

# COMBINATORIAL INEQUALITIES

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ABSTRACT. This is an expanded version of the *AMS Notices* column with the same title. The text is unchanged, but we added acknowledgements and a large number of endnotes which provide the context and the references.

*Combinatorics* has always been a battleground of tools and ideas. That's why it's so hard to do, or even define.<sup>1</sup> The inequalities are a particularly interesting case study as they seem to be both the most challenging and the least explored in Enumerative and Algebraic Combinatorics. Here are a few of my favorites, with some backstories.<sup>2</sup>

We start with *unimodality*<sup>3</sup> of binomial coefficients:

$$(1) \quad \binom{n}{k-1} \leq \binom{n}{k}, \quad \text{for all } 1 \leq k \leq n/2.$$

This is both elementary and well known – the proof is an easy calculation. But ask yourself the following natural question: does the difference  $B(n, k) := \binom{n}{k} - \binom{n}{k-1}$  count anything interesting? It should, of course, right? Imagine there is a natural injection

$$\psi : \binom{[n]}{k-1} \rightarrow \binom{[n]}{k}$$

from  $(k-1)$ -subsets to  $k$ -subsets of  $[n]$ , where  $[n] := \{1, \dots, n\}$ . Then  $B(n, k)$  can be described as the number of  $k$ -subsets of  $[n]$  that are not in the image of  $\psi$ , as good answer as any. But how do you construct the injection  $\psi$ ?<sup>4</sup>

Let us sketch the construction based on the classical *reflection principle* for the *ballot problem*, which goes back to the works of Bertrand and André in 1887. Start with a  $(k-1)$ -subset  $X$  of  $[n]$ , and let  $\ell$  be the smallest integer s.t.  $|X \cap [2\ell + 1]| = \ell$ . Such  $\ell$  exists since  $k \leq n/2$ . Define

$$\psi(X) := (X \setminus [2\ell + 1]) \cup ([2\ell + 1] \setminus X).$$

Observe that  $|\psi(X)| = k$  and check that  $\psi$  is the desired injection. This gives an answer to the original question:  $B(n, k)$  is the number of  $k$ -subsets  $Y \subset [n]$ , s.t.  $|Y \cap [m]| \leq m/2$  for all  $m$ .<sup>5</sup>

At this point you might be in disbelief in me dwelling on the easy inequality (1). Well, it only gets harder from here. Consider, e.g., the following question: Does there exist an injection  $\psi$  as above, s.t.  $X \subset \psi(X)$  for all  $X \in \binom{[n]}{k-1}$ ? We leave it to the reader as a challenge.<sup>6</sup>

There is also a curious connection to Algebraic Combinatorics:  $B(n, k) = f^{(n-k, k)}$ , the dimension of the irreducible  $S_n$ -module corresponding to the partition  $(n-k, k)$ . To understand how this could happen, think of both sides of (1) as dimensions of permutation representations of  $S_n$ . Turn both sides into vector spaces and modify  $\psi$  accordingly, to make it an  $S_n$ -invariant linear map. This would make it more natural and uniquely determined.<sup>7</sup> As a consequence, we obtain a combinatorial interpretation  $B(n, k) = |\text{SYT}(n-k, k)|$ , the number of standard Young tableaux of shape  $(n-k, k)$ , a happy outcome in every way.

Consider now *unimodality of Gaussian coefficients*:

$$(2) \quad p(n, k, \ell - 1) \leq p(n, k, \ell), \quad \text{for all } 1 \leq \ell \leq k(n-k)/2, \quad \text{where}$$

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$p(n, k, \ell)$  is the number of integer partitions  $\lambda \vdash \ell$  that fit into a  $k \times (n - k)$  rectangle, i.e.  $\lambda$  has parts of size at most  $(n - k)$ , and has at most  $k$  parts. To understand the context of this inequality, recall:

$$\sum_{\ell=0}^{k(n-k)} p(n, k, \ell) q^\ell = \binom{n}{k}_q := \frac{(n!)_q}{(k!)_q \cdot ((n-k)!)_q}, \quad \text{where } (n!)_q := \prod_{i=1}^n \frac{q^i - 1}{q - 1}.$$

To connect this to (1), note that  $\binom{n}{k}_1 = \binom{n}{k}$ , and that  $\binom{n}{k}_q$  is the number of  $k$ -subspaces of  $\mathbb{F}_q^n$ . In (2), we view  $\binom{n}{k}_q$  as a polynomial in  $q$  and compare its coefficients. Now, the Schubert cell decomposition of the Grassmannian over  $\mathbb{F}_q$ , or a simple induction can be used to give the partition interpretation.<sup>8</sup>

The inequality (2) is no longer easy to prove. Conjectured by Cayley in 1856, it was established by Sylvester in 1878; the original paper is worth reading even if just to see how pleased Sylvester was with his proof. In modern language, Sylvester defined the  $\mathfrak{sl}_2(\mathbb{C})$  action on certain homogeneous polynomials and the result follows from the highest weight theory (in its simplest form for  $\mathfrak{sl}_2$ ).<sup>9</sup>

Let's continue with the questions as we did above. Consider the difference  $C(n, k, \ell) := p(n, k, \ell) - p(n, k, \ell - 1)$ . Does  $C(n, k, \ell)$  count anything interesting? Following the pattern above, wouldn't it be natural to define some kind of nice injection from partitions of size  $(\ell - 1)$  to partition of size  $\ell$ , by simply adding a corner square according to some rule? That would be an explicit combinatorial (as opposed to algebraic) version of Sylvester's approach.

Unfortunately we don't know how to construct such a nice injection.<sup>10</sup> It's just the first of the many frustrations one encounters with algebraic proofs. Most of them are simply too rigid to be "combinatorialized". It doesn't mean that there is no combinatorial interpretation for  $C(n, k, \ell)$  at all. There is one very uninteresting interpretation due to Panova and myself, based on a very interesting (but cumbersome) identity by O'Hara.<sup>11</sup> Also, from the Computer Science point of view, it is easy to show that  $C(n, k, \ell)$  as a function is in  $\#\text{P}$ . We leave it to the reader to figure out why (or what does that even mean).<sup>12</sup>

To finish this story, we should mention Stanley's 1989 approach to (2) using finite group actions.<sup>13</sup> More recently, Panova and I introduced a different technique based on properties of the *Kronecker coefficients* of  $S_n$ , via the equality  $C(n, k, \ell) = g((n-k)^k, (n-k)^k, (n-\ell, \ell))$ .<sup>14</sup> Here the Kronecker coefficients  $g(\lambda, \mu, \nu)$  can be defined as structure constants for products of  $S_n$  characters:  $\chi^\mu \chi^\nu = \sum_\lambda g(\lambda, \mu, \nu) \chi^\lambda$ .<sup>15</sup> Both approaches imply stronger inequalities than (2), but neither gets us closer to a simple injective proof.

We turn now to *log-concavity of independent sets*:

$$(3) \quad a_{k-1}(M) \cdot a_{k+1}(M) \leq a_k(M)^2, \quad \text{where}$$

$a_k(M)$  is the number of independent  $k$ -subsets of a matroid  $M$ .<sup>16</sup> Note that the log-concavity implies unimodality, and in the special case of a *free matroid* (all elements are independent) this gives (1).<sup>17</sup>

The inequality (3) is a celebrated recent result by Adiprasito, Huh and Katz (2018), which showed that a certain "cohomology ring" associated with  $M$  satisfies the hard Lefschetz theorem and the Hodge–Riemann relations. This resolved conjectures by Welsh and Mason (1970s).<sup>18</sup>

It would be naïve for us to ask for a direct combinatorial proof via an injection, or by some other elementary means.<sup>19</sup> For example, Stanley in 1981 used the Aleksandrov–Fenchel inequalities in convex geometry to prove that the log-concavity is preserved under taking truncated sum with a free matroid<sup>20</sup>, already an interesting but difficult special case proved by inherently non-combinatorial means.<sup>21</sup>

There is also a Computational Complexity version of the problem which might be of interest. Let  $A(k, M) := a_k(M)^2 - a_{k-1}(M) \cdot a_{k+1}(M)$ . Does  $A(k, M)$  count any set of combinatorial objects?

For the sake of clarity, let  $G = (V, E)$  be a simple connected graph and  $M$  the corresponding matroid, i.e. bases in  $M$  are spanning trees in  $G$ . Then  $a_k(M)$  is the number of spanning forests

in  $G$  with  $k$  edges. Note that computing  $a_k(M)$  is  $\#\text{P}$ -complete in full generality.<sup>22</sup> Therefore, computing  $A(k, M)$  is  $\#\text{P}$ -hard.

Now,  $A(k, M)$  is in  $\text{GapP}$ , i.e. equal to the difference of two  $\#\text{P}$ -functions. Does  $A(k, M)$  lie in  $\#\text{P}$ ? This seems unlikely, but the current state of art of Computational Complexity doesn't seem to provide us with tools to even approach a negative solution.<sup>23</sup>

To fully appreciate the last example, consider the *log-concavity of matching numbers*:

$$(4) \quad m_{k-1}(G) \cdot m_{k+1}(G) \leq m_k(G)^2, \quad \text{where}$$

$m_k(G)$  is the number of  $k$ -matchings in a simple graph  $G = (V, E)$ , i.e.  $k$ -subsets of edges which are pairwise disjoint. For example,  $m_n(K_{2n}) = (2n-1) \cdots 3 \cdot 1$ . While perfect matchings don't necessarily define a matroid, they do have a similar flavor from a Combinatorial Optimization point of view.<sup>24</sup> The inequality (4) goes back to Heilmann and Lieb (1972) and is a rare case when the injection strategy works well.<sup>25</sup> The following argument is due to Krattenthaler (1996).<sup>26</sup>

Take a  $(k-1)$ -matching  $\beta$  whose edges we color *blue* and a  $(k+1)$ -matching  $\gamma$  whose edges we color *green*. The union  $\beta \cup \gamma$  of these two sets of edges splits into connected components, which are either paths or cycles, all alternately colored. Ignore for the time being all cycles and paths of even lengths. Denote by  $(r-1)$  the number of odd-length paths which have extra color blue. There are then  $(r+1)$  odd-length paths which have extra color green.

Now, allow switching colors in any of the  $2r$  odd-length paths. After recoloring, we want to have  $r$  odd-length paths extra color blue and the same with green. This amounts to a constructive injection from  $(r-1)$ -subsets of  $[2r]$  to  $r$ -subsets of  $[2r]$ , which we already know how to do as a special case of proving (1).

We leave to the reader the problem of finding an explicit combinatorial interpretation for  $M(k, G) := m_k(G)^2 - m_{k-1}(G) \cdot m_{k+1}(G)$ , proving that this function is in  $\#\text{P}$ . Note that computing  $m_k(G)$  is famously  $\#\text{P}$ -complete, which implies that so is  $M(k, G)$ . This makes the whole connection to Computational Complexity even more confusing. What exactly makes matchings special enough for this argument to work?<sup>27</sup>

If there is any pattern to the previous examples, it can be summarized as follows: the deeper one goes in an algebraic direction, the more involved are the inequalities and the less of a chance of a combinatorial proof. To underscore this point, consider the following three *Young tableaux inequalities*:

$$(5) \quad (f^\lambda)^2 \leq n!, \quad (c_{\mu\nu}^\lambda)^2 \leq \binom{n}{k}, \quad c_{\mu\nu}^\lambda \leq c_{\mu \vee \nu, \mu \wedge \nu}^\lambda, \quad \text{for all } \lambda \vdash n, \mu \vdash k, \nu \vdash n-k.$$

Here  $f^\lambda = |\text{SYT}(\lambda)|$  is the number of standard Young tableaux of shape  $\lambda$ , equal to the dimension of the corresponding irreducible  $S_n$ -module as above. Similarly,  $c_{\mu\nu}^\lambda = |\text{LR}(\lambda/\mu, \nu)|$  is the *Littlewood–Richardson coefficient*, equal to the number of Littlewood–Richardson tableaux of shape  $\lambda/\mu$  and weight  $\nu$ . It can be defined as a structure constant for products of Schur functions:  $s_\mu s_\nu = \sum_\lambda c_{\mu\nu}^\lambda s_\lambda$ . Finally,  $\mu \vee \nu$  and  $\mu \wedge \nu$  denote the union and intersection, respectively, of the corresponding Young diagrams.<sup>28</sup>

Now, the first inequality in (5) is trivial algebraically, but its combinatorial proof is highly nontrivial – it is a restriction of the *RSK correspondence*.<sup>29</sup> The second inequality is quite recent and follows easily from the definition and the Frobenius reciprocity. We believe it is unlikely that there is a combinatorial injection, even though there is a nice double counting argument.<sup>30</sup>

Finally, the third inequality in (5) is a corollary of the powerful inequality by Lam, Postnikov and Pylyavskyy (2007) using the curious Temperley–Lieb immanent machinery.<sup>31</sup> The key ingredient in the proof is Haiman's theorem which in turn uses the Kazhdan–Lusztig conjecture proven by Beilinson–Bernstein and Brylinski–Kashiwara.<sup>32</sup> While stranger things have happened, we would be very surprised if this inequality had a simple combinatorial proof.<sup>33</sup>

We conclude on a positive note, with a combinatorial inequality where everything works as well as it possibly could. Consider the following *majorization property of contingency tables*:

$$(6) \quad T(\mathbf{a}, \mathbf{b}) \leq T(\mathbf{a}', \mathbf{b}') \quad \text{for all } \mathbf{a}' \trianglelefteq \mathbf{a}, \mathbf{b}' \trianglelefteq \mathbf{b}.$$

Here  $\mathbf{a} = (a_1, \dots, a_m)$ ,  $a_1 \geq \dots \geq a_m > 0$ , and  $\mathbf{b} = (b_1, \dots, b_n)$ ,  $b_1 \geq \dots \geq b_n > 0$ , are two integer sequences with equal sum:

$$\sum_{i=1}^m a_i = \sum_{j=1}^n b_j = N.$$

A *contingency table* with margins  $(\mathbf{a}, \mathbf{b})$  is an  $m \times n$  matrix of non-negative integers whose  $i$ -th row sums to  $a_i$  and whose  $j$ -th column sums to  $b_j$ , for all  $i \in [m]$  and  $j \in [n]$ .  $T(\mathbf{a}, \mathbf{b})$  denotes the number of all such matrices. Finally, for sequences  $\mathbf{a}$  and  $\mathbf{a}'$  with the same sum, we write  $\mathbf{a} \trianglelefteq \mathbf{a}'$  if  $a_1 \leq a'_1$ ,  $a_1 + a_2 \leq a'_1 + a'_2$ ,  $a_1 + a_2 + a_3 \leq a'_1 + a'_2 + a'_3$ ,  $\dots$ . In other words, the inequality (6) says that there are more contingency tables when the margins are more evenly distributed.

Contingency tables can be viewed as adjacency matrices of bipartite multi-graphs with given degree distribution. They play an important role in Statistics and Network Theory.<sup>34</sup> We learned the inequality (6) from a paper by Barvinok (2007), but it feels like something that should have been known for decades.<sup>35</sup>

Now, we know two fundamentally different proofs of (6). The first is an algebraic proof using Schur functions which amounts to proving the following standard inequality for Kostka numbers:  $K_{\lambda\mu} \leq K_{\lambda\nu}$  for all  $\mu \trianglerighteq \nu$ , where  $K_{\lambda\mu}$  is the number of semistandard Young tableaux of shape  $\lambda$  and weight  $\mu$ . This inequality can also be proved directly, so combined with the RSK we obtain an injective proof of (6).<sup>36</sup>

Alternatively, one can prove the inequality directly for  $2 \times n$  rectangles and  $(+1, -1)$  changes in row (column) sums.<sup>37</sup> Combining these injections together gives a cumbersome, yet explicit injection. In principle, either of the two approaches can then be used to give a combinatorial interpretation for  $T(\mathbf{a}', \mathbf{b}') - T(\mathbf{a}, \mathbf{b})$ .<sup>38</sup>

In conclusion, let us note that we came full circle. Let  $m = 2$ ,  $a_1 = n - k + 1$ ,  $a_2 = k - 1$ ,  $a'_1 = n - k$ ,  $a'_2 = k$ , and  $b_1 = \dots = b_n = b'_1 = \dots = b'_n = 1$ .<sup>39</sup> Observe that  $T(\mathbf{a}, \mathbf{b}) = \binom{n}{k-1}$  and  $T(\mathbf{a}', \mathbf{b}') = \binom{n}{k}$ . The inequality (1) is a special case of (6) then.

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## NOTES

1. See the discussion in Ch. 6 of G.-C. Rota, *Discrete thoughts*, Birkhäuser, Boston, MA, 1992. See also our blog post “*What is Combinatorics?*” at <https://wp.me/p211iQ-bQ>.
2. The selection here is somewhat arbitrary, but we tried to include some of the greatest hits which fit a coherent narrative. Unfortunately, many interesting results are omitted. For example, we omitted a direct injective proof of the ad hoc inequality  $E_n \cdot F_n \geq n!$ , where  $E_n$  is the *Euler number* (the number of alternating permutations in  $S_n$ ), and  $F_n$  is the *Fibonacci number*. A more general inequality  $e(P)e(P^*) \geq n!$  for the number of linear extensions of a 2-dim poset and its plane dual remains open. See A. Morales, I. Pak, G. Panova, Why is  $\pi < 2\phi$ ?, *Amer. Math. Monthly* **125** (2018), 715–723, and our blog post “*Fibonacci times Euler*” at <https://wp.me/p211iQ-1J>.
3. For the general background on unimodal and log-concave sequences, see P. Brändén, Unimodality, log-concavity, real-rootedness and beyond, in *Handbook of Enumerative Combinatorics*, CRC Press, Boca Raton, FL, 2015, 437–483; [arXiv:1410.6601](https://arxiv.org/abs/1410.6601).
4. There are several proofs of this results, via injection and otherwise, some mentioned in the main body of the paper and in the notes below. A nice introductory overview is given in D. Zeilberger, [arXiv:1003.1273](https://arxiv.org/abs/1003.1273).
5. For details of this proof and various generalizations see B. E. Sagan, Unimodality and the Reflection Principle, *Ars Combin.* **48** (1998), 65–72. For the history of the ballot problem, the reflection principle, and for further references, see I. Pak, History of Catalan Numbers, §7, in R. P. Stanley, *Catalan Numbers*, Cambridge Univ. Press, 2015, App. B.

6. One can prove this from the *Hall marriage theorem*. In fact, a stronger result holds, that the Boolean algebra  $\mathcal{B}_n$  can be partitioned into symmetric saturated chains. Use induction. Suppose  $\mathcal{B}_{n-1} = \sqcup C_i$  and let  $C'_i := C_i \cup \{n\}$ . Now move the smallest element of  $C_i$  to  $C'_i$  and observe that  $\mathcal{B}_n = \sqcup C_i \sqcup C'_i$  as desired. In fact, this injection can be computed in polynomial time via the elegant “parenthesization” algorithm. For this construction and further references see C. Greene, D. J. Kleitman, Strong versions of Sperner’s theorem, *J. Combin. Theory Ser. A* **20** (1976), 80–88, and §3 in C. Greene, D. J. Kleitman, Proof techniques in the theory of finite sets, in *Studies in combinatorics*, MAA, 1978, 22–79.

7. This is equivalent to saying that

$$M^{(n-k,k)} := \text{Ind}_{S_k \times S_{n-k}}^{S_n} \mathbf{1} = \mathbb{S}^{(n-k,k)} \oplus \mathbb{S}^{(n-k+1,k-1)} \oplus \dots \oplus \mathbb{S}^{(n-1,1)} \oplus \mathbb{S}^{(n)},$$

where  $\mathbb{S}^\lambda$  is the irreducible  $S_n$ -module corresponding to  $\lambda$ . The uniqueness of the  $S_n$ -invariant injective map  $\psi : M^{(n-k+1,k-1)} \rightarrow M^{(n-k,k)}$  follows from the unique decomposition in the multiplicity free case.

8. For a full explanation of this argument see §1.10 in R. P. Stanley, *Enumerative Combinatorics*, Vol. 1 (Second ed.), Cambridge Univ. Press, 2012. See also I. Pak, When and how  $n$  choose  $k$ , in *AMS DIMACS Ser.*, vol. 43 (1998), 191–238.
9. For a friendly exposition of this approach, see R. A. Proctor, Solution of two difficult combinatorial problems with linear algebra, *Amer. Math. Monthly* **89** (1982), 721–734. A different elegant proof is given in §23 of J. Matoušek, *Thirty-three miniatures*, AMS, Providence, RI, 2010. For a more general setting, see R. P. Stanley, Unimodal sequences arising from Lie algebras, in *Lecture Notes Pure Appl. Math.* **57**, Dekker, New York, 1980, 127–136.
10. The existence of an injection is known and follows from the *Dilworth theorem* and the *Peck property*, see R. P. Stanley, Quotients of Peck posets, *Order* **1** (1984), 29–34. There is no known efficient algorithm to compute this injection. Finally, the existence of a symmetric saturated chain decomposition in this case is a well-known open problem, see R. P. Stanley, Weyl groups, the hard Lefschetz theorem, and the Sperner property, *SIAM J. Algebraic Discrete Methods* **1** (1980), 168–184.
11. A combinatorial interpretation for  $C(n, k, \ell)$  based on O’Hara’s identity due to I. Pak and G. Panova, is given on p. 9 in <https://tinyurl.com/ydemhyf5>. For O’Hara’s identity and its simple proof, see I. G. Macdonald, An elementary proof of a  $q$ -binomial identity, in *q-series and partitions*, IMA, Springer, New York, 1989, 73–75. For O’Hara’s combinatorial proof based on the identity, see D. Zeilberger, Kathy O’Hara’s constructive proof of the unimodality of the Gaussian polynomials, *Amer. Math. Monthly* **96** (1989), 590–602.
12. This means that there is a set of 0-1 sequences whose size is  $C(n, k, \ell)$ , s.t. the membership can be decided in polynomial time (in  $n$ ). This follows immediately from the fact that  $C(n, k, \ell)$  can be computed in polynomial time since Gaussian coefficients have an explicit recursive formula

$$\binom{n}{k}_q = \binom{n-1}{k}_q + q^{n-k} \binom{n-1}{k-1}_q$$

and have  $O(n^2)$  degree as a polynomial in  $q$ .

13. R. P. Stanley, Log-concave and unimodal sequences in algebra, combinatorics, and geometry, in *Ann. New York Acad. Sci.* **576** (1989), 500–535. See also White’s paper<sup>36</sup>.
14. I. Pak, G. Panova, Strict unimodality of  $q$ -binomial coefficients, *C.R. Acad. Sci. Paris, Ser. I. Math.* **351** (2013), 415–418, and I. Pak, G. Panova, Bounds on certain classes of Kronecker and  $q$ -binomial coefficients, *J. Combin. Theory Ser. A.* **147** (2017), 1–17.
15. The Kronecker coefficients are quite interesting, but very mysterious and nobody’s friend. They are hard to compute, have no known combinatorial interpretation, and tend to appear in different areas of science. For connections to Computational Complexity, see C. Ikenmeyer, K. Mulmuley, M. Walter, On vanishing of Kronecker coefficients, *Comput. Complexity* **26** (2017), 949–992, and I. Pak, G. Panova, On the complexity of computing Kronecker coefficients, *Comput. Complexity* **26** (2017), 1–36. For connections to Quantum Information Theory, see M. Christandl, G. Mitchison, The spectra of quantum states and the Kronecker coefficients of the symmetric group, *Comm. Math. Phys.* **261** (2006), 789–797.
16. For the background on matroids, see N. White, *Theory of Matroids*, Cambridge Univ. Press, 1986, and J. Oxley, *Matroid theory* (Second ed.), Oxford Univ. Press, 2011.
17. To be clear, for a free matroid  $M$  on  $n$  elements, we have  $a_k(M) = \binom{n}{k}$ .
18. K. Adiprasito, J. Huh, E. Katz, Hodge theory for combinatorial geometries, *Ann. Math.* **188** (2018), 381–452. See also a friendly introduction by the same authors in *Notices AMS* **64** (2017), no. 1, 26–30.
19. This does not mean that no simple non-combinatorial proof exists. In recent years, the Adiprasito–Huh–Katz approach has been substantially simplified and extended. See N. Anari, K. Liu, S. Oveis Gharan, C. Vinzant, Log-Concave Polynomials III: Mason’s Ultra-Log-Concavity Conjecture for Independent Sets of Matroids, [arXiv:1811.01600](https://arxiv.org/abs/1811.01600), and P. Brändén, J. Huh, Hodge–Riemann relations for Potts model partition functions, [arXiv:1811.01696](https://arxiv.org/abs/1811.01696).
20. To understand the abstract definition of the truncated sum of matroids, think of two sets of vectors  $S$  and  $S'$  in vector spaces  $V$  and  $V'$ , of dimensions  $d$  and  $d'$ , respectively. Suppose  $\max\{d, d'\} \leq k \leq d + d'$ . The truncated matroid is given by projection  $S \times S' \subset V \times V'$  onto a generic  $k$ -subspace.
21. R. P. Stanley, Two combinatorial applications of the Aleksandrov–Fenchel inequalities, *J. Combin. Theory Ser. A* **31** (1981), 56–65. For the extension of this proof to truncated sums of general matroids (using a different terminology and motivation), see L. Gurvits, A short proof, based on mixed volumes, of Liggett’s theorem on the convolution of ultra-logconcave sequences, *El. J. Combin.* **16** (2009), no. 1, Note 5, 5 pp.
22. For a comprehensive discussion of complexity of this and other values of the Tutte polynomial, see D. J. A. Welsh, *Complexity: knots, colourings and counting*, Cambridge Univ. Press, 1993.
23. For the recent overview of complexity of computing combinatorial sequences, see I. Pak, Complexity problems in enumerative combinatorics, in *Proc. ICM Rio de Janeiro*, Vol. 3, 2018, 3139–3166; revised and expanded version in [arXiv:1803.06636](https://arxiv.org/abs/1803.06636).
24. By that we mean that the linear programming works in polynomial time on both perfect matchings in graphs and bases of matroids. For an overview and a careful explanation of the connection, see W. H. Cunningham, Matching, matroids, and extensions, *Math. Program.* **91** (2002), Ser. B, 515–542.
25. O. J. Heilmann, E. H. Lieb, Theory of monomer-dimer systems, *Comm. Math. Phys.* **25** (1972), 190–243.
26. C. Krattenthaler, Combinatorial proof of the log-concavity of the sequence of matching numbers, *J. Combin. Theory Ser. A* **74** (1996), 351–354.
27. For several explicit combinatorial injections in the context of log-concavity, see B. E. Sagan, Inductive and injective proofs of log concavity results, *Discrete Math.* **68** (1988), 281–292. In the opposite direction, see an example of a log-convex combinatorial sequence in S. DeSalvo, I. Pak, Log-concavity of the partition function, *Ramanujan J.* **38** (2015), 61–73. Here the result is that partition function satisfies  $p(n)^2 - p(n-1)p(n+1) \geq 0$ , for all  $n > 25$ . It is unlikely this inequality has a direct combinatorial proof.
28. For a comprehensive introduction to combinatorics of Young tableaux and the Littlewood–Richardson coefficients, see B. E. Sagan, *The Symmetric Group*, Springer, New York, 2001, and R. P. Stanley, *Enumerative Combinatorics*, Vol. 2, Cambridge U. Press, Cambridge, 1999, Ch. 7.
29. The RSK is a fundamental bijection described in both Sagan and Stanley’s books, see above.
30. The inequality is given in I. Pak, G. Panova, D. Yeliussizov, On the largest Kronecker and Littlewood–Richardson coefficients, *J. Combin. Theory Ser. A* **165** (2019), 44–77, §4. The “double counting injection” is given by the

following one line argument:

$$\sum_{\lambda \vdash n} (c_{\mu, \nu}^{\lambda})^2 f^{\mu} f^{\nu} = \sum_{\lambda \vdash n} c_{\mu, \nu}^{\lambda} (c_{\mu, \nu}^{\lambda} f^{\mu} f^{\nu}) \leq \sum_{\lambda \vdash n} c_{\mu, \nu}^{\lambda} f^{\lambda} = \binom{n}{k} f^{\mu} f^{\nu}.$$

Here the inequality and the last equality can be made injective by using two different combinatorial interpretations of the Littlewood–Richardson coefficients.

31. T. Lam, A. Postnikov, P. Pylyavskyy, Schur positivity and Schur log-concavity, *Amer. J. Math.* **129** (2007), 1611–1622. The inequality there is more general and stated in a different but equivalent way; to see the connection, consider only straight shapes and write the inequalities in Schur functions basis. Also, this inequality is a refinement of an older, FKG-type inequality:

$$f^{\lambda} \cdot f^{\mu} \leq f^{\lambda \vee \mu} \cdot f^{\lambda \wedge \mu}.$$

The proof of the latter is elementary and based on the hook-length formula, see A. Björner, A  $q$ -analogue of the FKG inequality and some applications, *Combinatorica* **31** (2011), 151–164.

32. M. Haiman, Hecke algebra characters and immanant conjectures, *JAMS* **6** (1993), 569–595.

33. Another notable example is the inequality implied by the *Naruse hook-length formula* (NHLF):

$$(\diamond) \quad f^{\lambda/\mu} \geq n! \prod_{(i,j) \in \lambda/\mu} \frac{1}{h(i,j)},$$

where  $f^{\lambda/\mu} = |\text{SYT}(\lambda/\mu)|$  is the number of standard Young tableaux of skew shape and  $h(i,j) = \lambda_i + \lambda'_j - i - j + 1$  is the hook-length in  $\lambda$ . When  $\mu = \emptyset$ , the inequality becomes an equality; it is the celebrated *hook-length formula* with a plethora of bijective, probabilistic, and algebraic proofs. Unfortunately, no direct injective proof is known for  $(\diamond)$ . In a special case of  $\lambda/\mu = \tau^*$ , a Young diagram  $\tau$  rotated 180°, it gives the following *hook inequality*:

$$(\heartsuit) \quad \prod_{(i,j) \in \tau} h(i,j) \leq \prod_{(i,j) \in \tau} h^*(i,j),$$

where  $h^*(i,j) = i + j - 1$  are the hooks in  $\tau^*$ . Even this inequality does not seem to have a direct injective proof. For  $(\heartsuit)$  and further discussion and references to the NHLF, see A. Morales, I. Pak, G. Panova, Asymptotics of the number of standard Young tableaux of skew shape, *Eur. J. Combin.* **70** (2018), 26–49. For a recent combinatorial (but not injective!) proof of  $(\heartsuit)$ , see I. Pak, F. Petrov, V. Sokolov, Hook inequalities, [arXiv:1903.11828](https://arxiv.org/abs/1903.11828).

34. See e.g. P. Diaconis, A. Gangolli, Rectangular arrays with fixed margins, *Disc. Prob. Alg.* **72** (1995), 15–41, and B. S. Everitt, *The analysis of contingency tables* (Second ed.), Chapman & Hall, London, 1992.

35. See Eq. (4) on p. 111 in A. Barvinok, Brunn–Minkowski inequalities for contingency tables and integer flows, *Adv. Math.* **211** (2007), 105–122. Note that a more general inequality (3) in that paper remains open.

36. Barvinok’s proof is exactly the combinatorial identity given by the RSK written in terms of the Kostka numbers. For a direct injective proof of the inequality for Kostka numbers, see D. E. White, Monotonicity and unimodality of the pattern inventory, *Adv. Math.* **38** (1980), 101–108.

37. To see this, note that for  $\mathbf{a} = (a_1, \dots, a_n)$ ,  $\mathbf{b} = (N - k, k)$ ,  $N = a_1 + \dots + a_n$ , we have:

$$T(\mathbf{a}, \mathbf{b}) = [q^k] \prod_{i=1}^n (1 + q + \dots + q^{a_i - 1})$$

Each term in the product is unimodal, thus so is the product (see e.g. Stanley’s paper<sup>13</sup> above). This implies unimodality. To get a direct injection, see the Greene–Kleitman symmetric chain decomposition approach described above.<sup>6</sup> Curiously, White’s argument<sup>36</sup> can be also translated into this language.

38. I am being very vague here both in an effort to avoid technicalities and at least to some extent not to reveal the work in progress. If and when this work is completed, I will update this note.

39. This is the case in the note above<sup>37</sup>, of course.