

LIFTS, DERANDOMIZATION, AND DIAMETERS OF SCHREIER GRAPHS OF MEALY AUTOMATA

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ABSTRACT. It is known that random 2 -lifts of graphs give rise to expander graphs. We present a new conjectured derandomization of this construction based on certain *Mealy automata*. We verify that these graphs have polylogarithmic diameter, and present a class of automata for which the same is true. However, we also show that some automata in this class do not give rise to expander graphs.

1. INTRODUCTION

In [BL], Bilu and Linial showed that random 2 -lifts of expanding graphs remain expanding with high probability. This gives a probabilistic construction of expander families. Several ways to derandomize this procedure are also given in [BL], but none of them give a *strongly explicit* description of a family of expander graphs. That is, a description in which the actual graph is much larger than working memory, but a computer can list neighbors of a vertex in polylogarithmic (in the size of the graph) time.

We consider the following two families of 2 -lifts of graphs. The *Aleshin graphs* A_0, A_1, A_2, \dots are a sequence of 3 -regular edge-labeled directed graphs. The first graph A_0 is defined to be a single vertex with three self-loops labeled a, b , and c . Given the graph A_n , the next graph A_{n+1} is defined as a certain *graph lift* of A_n : Each vertex $v \in A_n$ lifts to two vertices $v_0, v_1 \in A_{n+1}$, and the edges transform as follows:

$$\begin{array}{lll}
 v \xrightarrow{a} w & \text{lifts to} & \begin{array}{l} v_0 \xrightarrow{c} w_0 \\ v_1 \xrightarrow{c} w_1 \end{array} \\
 v \xrightarrow{b} w & \text{lifts to} & \begin{array}{l} v_0 \xrightarrow{a} w_0 \\ v_1 \xrightarrow{b} w_1 \end{array} \\
 v \xrightarrow{c} w & \text{lifts to} & \begin{array}{l} v_0 \xrightarrow{b} w_0 \\ v_1 \xrightarrow{a} w_1 \end{array}
 \end{array}$$

That is, e.g., if A_n has an edge labeled c from v to w , then A_{n+1} has an edge labeled b from v_0 to w_1 , and an edge labeled a from v_1 to w_0 .

Another family, the *Bellaterra graphs* B_0, B_1, B_2, \dots is defined the same way, except with transformation rules

$$\begin{array}{lll}
 v \xrightarrow{a} w & \text{lifts to} & \begin{array}{l} v_0 \xrightarrow{c} w_0 \\ v_1 \xrightarrow{c} w_1 \end{array} \\
 v \xrightarrow{b} w & \text{lifts to} & \begin{array}{l} v_0 \xrightarrow{a} w_0 \\ v_1 \xrightarrow{b} w_1 \end{array} \\
 v \xrightarrow{c} w & \text{lifts to} & \begin{array}{l} v_0 \xrightarrow{b} w_0 \\ v_1 \xrightarrow{a} w_1 \end{array}
 \end{array}$$

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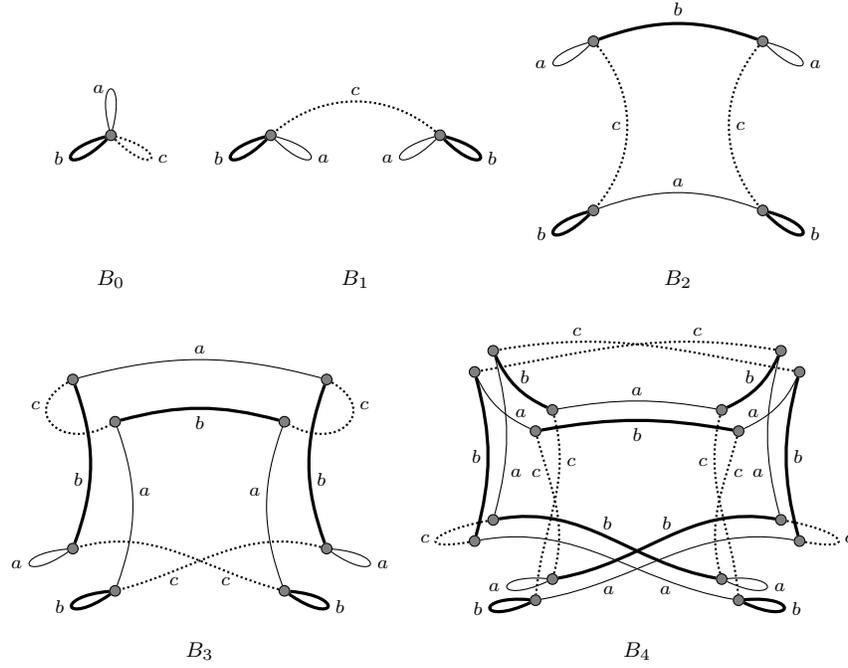


FIGURE 1. The Bellaterra graphs.

It is not hard to check that the reverse of every edge in B_n is also in B_n , so these can be thought of as undirected graphs. The first few graphs in this family are pictured in Figure 1.

The main result of this paper is the following theorem:

Theorem 1.1. *The diameter of the Aleshin graphs $\{A_i\}_{i=1}^{\infty}$ and Bellaterra graphs $\{B_i\}_{i=1}^{\infty}$ grows at most quadratically in n , i.e.,*

$$\text{diam}(A_n) = O(n^2) \quad \text{and} \quad \text{diam}(B_n) = O(n^2) \quad \text{as } n \rightarrow \infty.$$

Prior to this paper, there were no nontrivial bounds on the diameter of A_n ; even subexponential bounds remained out of reach. Note also that in principle we can start with *any* 3-labeled graph in place of $A_0 = B_0$, and proceed making lifts as above. We do not consider these in the paper, and our algebraic techniques do not apply.

Observe that both families of graphs are *very explicit* in the following sense: there is a polynomial time algorithm which, given a number n and $v \in \Gamma_n$, lists the neighbors of v . “Polynomial time” here refers to a runtime which is polynomial in the number of bits necessary to describe the input. It takes n bits to describe a vertex of B_n or A_n , so the algorithm should run in time $O(n^d)$, for some d .

In particular, it follows that they are *strongly explicit* in the sense of [BL]: There is a polynomial time (in the size of the inputs) algorithm which, given a number n , and vertices $v, w \in \Gamma_n$, decides whether v and w are adjacent in Γ_n .

As we will see below, these graphs can be described in terms of invertible *Mealy automata*. The associated automata are small: they act on binary strings and have only 3 states. A detailed study of all such small automata was performed in [B+]. The Bellaterra and Aleshin automata are numbered 846 and 2240 in that article. They are the only nontrivial *bireversible* ones. Spectra of the first few associated graphs are also computed in [B+], and the data suggest that the Aleshin graphs are a family of expanders with eigenvalue gap roughly 0.2.

Conjecture 1.2. *The Aleshin graphs $\{A_i\}_{i=1}^{\infty}$ are a family of two-sided expanders.*

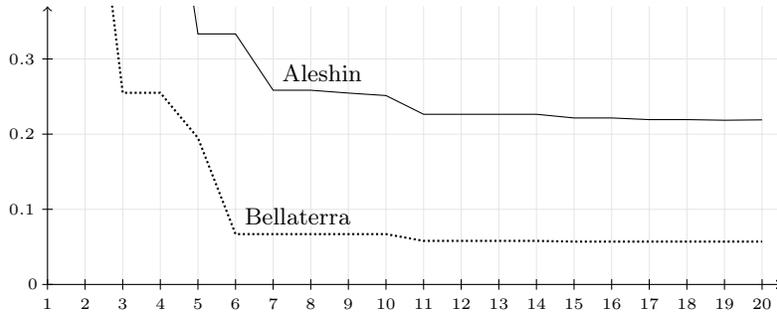


FIGURE 2. Eigenvalue gaps of the Bellaterra and Aleshin graphs.

Here by *two-sided* we mean that both the second largest and the smallest eigenvalues of 3-regular graphs A_n are bounded away as follows: $\lambda_2 < 3 - \varepsilon$ and $\lambda_n > -3 + \varepsilon$ (see e.g. [Tao]). This conjecture and the following one appear as Problem 10.1 in [G2].

Though it is less clear from the data in [B+], our own computations (see Figure 2) suggest that the Bellaterra graphs are also expanders, with eigenvalue gap roughly 0.05, so we make the stronger conjecture:¹

Conjecture 1.3. *The Bellaterra graphs $\{B_i\}_{i=1}^{\infty}$ are a family of two-sided expanders.*

If so, they are a strongly explicit derandomization of the probabilistic construction in [BL]. One consequence of being an expander family is logarithmic diameter growth with respect to the size of the graph, so if Conjecture 1.2 holds then $\text{diam}(A_n)$ grows linearly in n , a stronger claim than in the theorem.

Unfortunately, we are not near proving either conjectures and in fact our tools are too weak to prove them. Later in the paper, we state and prove general conditions on an automaton which guarantee polynomial diameter growth in the associated graphs (Section 8). We then prove that for some automata which satisfy those conditions, we do not get expanders (see Section 10). In other words, a different, perhaps combinatorial technique is needed to prove the expansion.

2. MEALY AUTOMATA

The Bellaterra graphs $\{B_n\}_{n=1}^{\infty}$ are *very explicit* in the sense of [HLW].² That is, there is a polynomial time algorithm which, given a number n and a vertex $v \in B_n$, lists the neighbors of v in B_n . It takes n bits to describe a vertex in B_n , so the runtime of the algorithm should be polynomial in n .

In fact, there is a linear time algorithm. Even more strongly, the computation can be implemented with a Mealy automaton, i.e., a finite state automaton which outputs a letter each time it reads a letter.

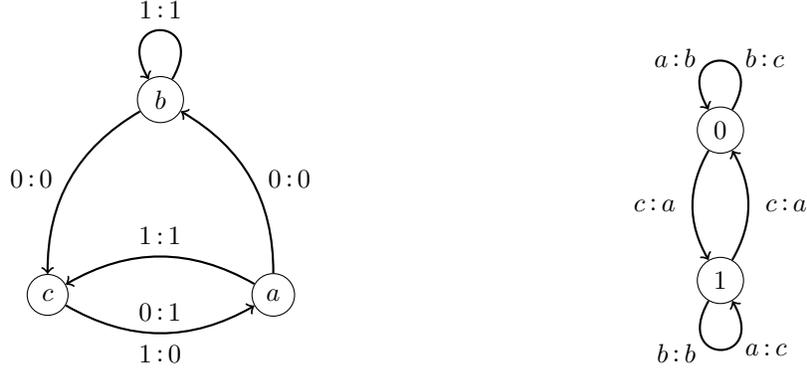
Definition 2.1. A *Mealy automaton* $\mathcal{M} = (Q, A, \tau, \sigma)$ is a pair of finite sets Q, A , together with functions $\sigma : Q \times A \rightarrow A$, and $\tau : Q \times A \rightarrow Q$.

The sets Q and A are called the *states* and *alphabet*, respectively. The functions σ and τ are called the *output* and *transition* functions, respectively. When $|Q| = q$ and $|A| = a$, we call \mathcal{M} a (q, a) -automaton. We adopt the following notation:

$$\begin{aligned} {}^q x &= \sigma_q(x) = \sigma(q, x) \\ q^x &= \tau_x(q) = \tau(q, x). \end{aligned}$$

¹See Remark 11.3.

²Sometimes, these are called *fully explicit*, see e.g. [Vad].

FIGURE 3. The Bellaterra automaton \mathcal{B} and its dual $\bar{\mathcal{B}}$.

Let A^* and A^∞ denote the set of finite and right-infinite words in the alphabet A , respectively, and let $A^{*,\infty} = A^* \cup A^\infty$ denote the set of all words in A . A Mealy automaton in the state $q \in Q$ acts in a length-preserving way on words in $A^{*,\infty}$ by reading the first letter x , outputting the letter $\sigma(q, x)$, and acting on the rest of the word from the state $\tau(q, x)$. That is, each $q \in Q$ has a corresponding length-preserving map $A^{*,\infty} \rightarrow A^{*,\infty}$ defined recursively by

$${}^q(x_0x_1 \dots x_n) = y_0 {}^r(x_1 \dots x_n),$$

$$\text{and } {}^q(x_0x_1x_2 \dots) = y_0 {}^r(x_1x_2 \dots),$$

where $y_0 = \sigma(q, x_0)$ and $r = \tau(q, x_0)$. This extends to a left action of finite words Q^* on words in $A^{*,\infty}$ via, e.g.,

$${}^{qr}s = {}^q({}^rs).$$

So we have defined an extension of $\sigma : Q \times A \rightarrow A$ to a map $\sigma : Q^* \times A^{*,\infty} \rightarrow A^{*,\infty}$ given by

$$\sigma(w, s) = \sigma_w(s) = {}^ws.$$

A Mealy automaton can be depicted with a Moore diagram: a directed graph with a vertex for each state $q \in Q$ and a labeled edge

$$q \xrightarrow{x:y} r$$

for every $q \in Q$ and every $x \in A$, where $y = \sigma(q, x)$ and $r = \tau(q, x)$. That is, an edge $q \xrightarrow{x:y} r$ denotes that if the Mealy automaton is in state q and reads the letter x , then it outputs the letter y and transitions to the state r . We will sometimes simply write $q \xrightarrow{x:y} r$ to denote that $y = \sigma(q, x)$ and $r = \tau(q, x)$.

Example 2.2. Consider the *Bellaterra automaton* \mathcal{B} pictured in Figure 3. More formally, $\mathcal{B} = (Q, A, \tau, \sigma)$ is defined by

$$A = \{0, 1\}, \quad Q = \{a, b, c\}$$

$$\sigma_a = \sigma_b = \text{id}, \quad \sigma_c = (0 \ 1),$$

$$\text{and } \tau_0 = (a \ b \ c), \quad \tau_1 = (a \ c)(b),$$

where we use the usual cycle notation for permutations, so e.g., $\tau_0(a) = b$, $\tau_0(b) = c$, $\tau_0(c) = a$. Note that σ_q and τ_a are permutations in this case, which need not be true for general automata. In the terminology defined below, this means \mathcal{B} is *invertible* and *reversible*, respectively.

Given a number n , the Bellaterra graph B_n can be described as the graph whose vertices are length n binary strings, with an edge

$$s \xrightarrow{q} ({}^qs)$$

for each vertex $s \in A^n$ and each state $q \in Q$. Where we give different colors to the edges corresponding to states a , b and c , as in the introduction and Figure 1. For example, we have

$${}^c(0000) = 1 {}^a(000) = 10 {}^b(00) = 100 {}^c(0) = 1001,$$

so $0000 \xrightarrow{c} 1001$.

Some symmetry between states and letters of a Mealy automaton is already apparent in the definition. The nature of this symmetry becomes more clear if we consider computing compositions of maps associated to the states of an automaton, we have, e.g.,

$${}^{q_1 q_0}(x_0 x_1 \dots x_n) = {}^{q_1}(y_0 {}^{r_0}(x_1 \dots x_n)) = z_0 {}^{r_1 r_0}(x_1 \dots x_n) = \dots,$$

where $q_0 \xrightarrow{x_0 : y_0} r_0$, and $q_1 \xrightarrow{y_0 : z_0} r_1$. The computation proceeds by taking any instance of ${}^q(x \dots)$ in the expression, and replacing it with $y^r(\dots)$, where $q \xrightarrow{x : y} r$.

If we ignore parentheses, states in Q and letters in A play a symmetric role in this process, except that letters in Q are written higher and disappear when they are at the right side of the expression. Taking this symmetry into account, the automaton also naturally defines an action of the letters in A on finite words in Q^* :

$$(q_n \dots q_1 q_0)^x = (q_n \dots q_1)^y r_0,$$

where $q_0 \xrightarrow{x : y} r_0$. Letters in A also act on the set of *left*-infinite words in the alphabet Q :

$$(\dots q_2 q_1 q_0)^x = (\dots q_2 q_1)^y r_0.$$

We let $Q^{-\infty}$ denote this set of left-infinite words, and let $Q^{-\infty,*}$ denote $Q^* \cup Q^{-\infty}$, so we have an action of A on $Q^{-\infty,*}$. This naturally extends to a *right* action of A^* on $Q^{-\infty,*}$, via, e.g.

$$w^{xy} = (w^x)^y.$$

So we have defined a map $\tau : Q^{-\infty,*} \times A^* \rightarrow Q^{-\infty,*}$, given by

$$\tau(w, s) = \tau_s(w) = w^s.$$

It is straightforward to check that for any $s \in A^*$, $t \in A^{*,\infty}$, $w \in Q^*$, $v \in Q^{-\infty,*}$, the actions we have defined satisfy the following relations:

$$\begin{aligned} w(st) &= \tilde{s} \tilde{w}(t), \quad \text{and} \\ (vw)^s &= (v)^{\tilde{s}} \tilde{w}, \end{aligned}$$

where $\tilde{s} = {}^w s$ and $\tilde{w} = w^s$.

If we need to specify that we are referring to a particular automaton \mathcal{M} , we will write $\sigma_{\mathcal{M}}$ for σ and $\sigma_{\mathcal{M},w}$ for σ_w , and similarly for τ .

With this symmetry in mind, it is sensible to define the *dual* of an automaton $\mathcal{M} = (Q, A, \tau, \sigma)$ to be the automaton $\overline{\mathcal{M}} = (\widehat{Q}, \widehat{A}, \widehat{\tau}, \widehat{\sigma})$ given by interchanging the roles of the states and alphabet. That is, we take

$$\widehat{A} = Q, \quad \widehat{Q} = A, \quad \widehat{\sigma}(a, q) = \tau(q, a), \quad \text{and} \quad \widehat{\tau}(a, q) = \sigma(q, a).$$

In other words, for $q, r \in Q$ and $x, y \in A$, we have $x \xrightarrow{q:r} y$ in $\overline{\mathcal{M}}$ if and only if $q \xrightarrow{x:y} r$ in \mathcal{M} .

Computations in the dual automaton are computations in the original automaton, with each step written backwards. It follows that, e.g., for every $s \in A^*$ and $w \in Q^*$ we have

$$\sigma_{\overline{\mathcal{M}}}(s, w) = \overline{\tau_{\mathcal{M}}(\bar{w}, \bar{s})},$$

where \bar{u} denotes the reversal of u , i.e. the word u written backwards.

Example 2.3. The dual of the Bellaterra automaton is also pictured in Figure 3.

Example 2.4. Let $A = \{0, 1\}$. Consider the Mealy automaton pictured in Figure 4. This automaton is called the *adding machine*, because the map $\sigma_r : \{0, 1\}^* \rightarrow \{0, 1\}^*$ is simply addition of 1, where length n words in A^* are interpreted as binary representations of numbers modulo 2^n , with the least significant digit on the left.

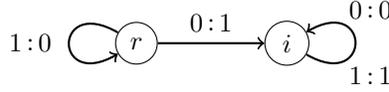


FIGURE 4. The binary adding machine

We say a Mealy automaton is *invertible* if σ_q is invertible for every $q \in Q$. This occurs if and only if the endomorphism $\sigma_w : A^* \rightarrow A^*$ is invertible for every $w \in Q^*$. We are primarily interested in invertible automata, though our results can be generalized to the non-invertible case.

The *inverse* of an invertible automaton $\mathcal{M} = (Q, A, \tau, \sigma)$ is the automaton $\mathcal{M}^{-1} = (Q', a, \tau', \sigma')$ given by

$$Q = \{q' \mid q \in Q\}, \quad \sigma'_{q'} = \sigma_q^{-1}, \quad \tau'(q', a) = \tau(q, \sigma_q^{-1}(a)).$$

In other words, the Moore diagram of \mathcal{M}^{-1} is the Moore diagram of \mathcal{M} , with each edge

$$q \xrightarrow{x:y} r$$

replaced by

$$q' \xrightarrow{y:x} r'.$$

It is straightforward to check that $\sigma_{\mathcal{M}^{-1}, q'} = \sigma_{\mathcal{M}, q}^{-1}$ for every $q \in Q$.

Consider two automata $\mathcal{M} = (Q, A, \tau, \sigma)$, $\mathcal{M}' = (Q', A, \tau', \sigma')$ acting on the same alphabet, with Q, Q' disjoint. Their *union* is the automaton $\mathcal{M} \cup \mathcal{M}' = (Q \cup Q', A, \tau'', \sigma'')$, where

$$\tau''(q, a) = \begin{cases} \tau(q, a) & q \in Q \\ \tau'(q, a) & q \in Q' \end{cases} \quad \text{and} \quad \sigma''(q, a) = \begin{cases} \sigma(q, a) & q \in Q \\ \sigma'(q, a) & q \in Q' \end{cases}$$

For example, $\mathcal{M} \cup \mathcal{M}^{-1}$ is an automaton with twice as many states as Q , in which every state q has an inverse state q' with $\sigma_{q'} = \sigma_q^{-1}$.

We say an automaton is *reversible* if its dual is invertible.

We say an automaton is *bireversible* if it is invertible, reversible, and its inverse is reversible. Note that the last condition does not follow from the other two.

Example 2.5. The three-state automaton in Figure 9 is reversible and invertible, but not bireversible.

Example 2.6. The adding machine pictured in Figure 4 and described in Example 2.4 is invertible, but not reversible.

Example 2.7. The Bellaterra automaton and the Aleshin automaton, pictured in Figure 3 and Figure 5 respectively, are both bireversible. The fact that the Bellaterra automaton is reversible is why we can specify τ_0 and τ_1 as permutations in Example 2.2.

3. SCHREIER GRAPHS

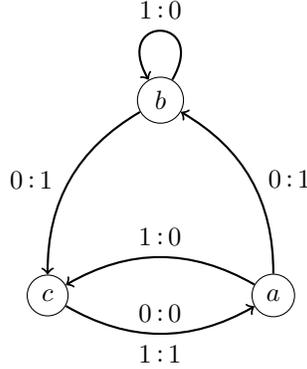
For our purposes, graphs are locally finite, directed, and may have self-loops and repeated edges. A graph is *regular* if the indegree and outdegree are the same across all vertices.

Let Γ be a graph. Given vertices $v, w \in \Gamma$, we write $v \rightarrow_{\Gamma} w$ if there is an edge in Γ from v to w . We write $d_{\Gamma}(v, w)$ for the distance between v and w , i.e. the length of the shortest undirected path between v and w . When there is no such path, we take $d_{\Gamma}(v, w) = \infty$. Given a nonnegative integer r , the *ball* of radius r centered at v is the set

$$B_{\Gamma}(v, r) = \{w \in \Gamma : d(v, w) \leq r\}.$$

The *diameter* of Γ is defined to be

$$\text{diam}(\Gamma) = \max_{v, w \in \Gamma} d_{\Gamma}(v, w).$$


 FIGURE 5. The Aleshin automaton \mathcal{A} .

When it is clear from context what graph we are discussing, we will drop the subscripts and simply write $v \rightarrow w$, $d(v, w)$, and $B(v, r)$.

In Example 2.2 we described the Bellaterra graphs in terms of a Mealy automaton. In the same way, we can associate a sequence of graphs to any Mealy automaton. Since we are primarily concerned with regular graphs, we require the automaton to be invertible.

Definition 3.1. Let $\mathcal{M} = (Q, A, \sigma, \tau)$ be an invertible Mealy automaton. Given $n \in \{1, 2, \dots\} \cup \{\infty\}$, the *Schreier graph* $\Gamma_{\mathcal{M}, n}$ is a directed graph, defined as follows: The vertices of $\Gamma_{\mathcal{M}, n}$ are length n words in $A^{*, \infty}$, i.e. elements of A^n . For each vertex $s \in \Gamma_{\mathcal{M}, n}$ and each state $q \in Q$, the Schreier graph $\Gamma_{\mathcal{M}, n}$ has an edge

$$s \rightarrow {}^q s.$$

Clearly, the number of edges leaving a vertex is $|Q|$. The Schreier graph of the inverse automaton, $\Gamma_{\mathcal{M}^{-1}, n}$, is simply $\Gamma_{\mathcal{M}, n}$ with the edges reversed. So, the number of edges entering a given vertex in $\Gamma_{\mathcal{M}^{-1}, n}$ is also $|Q|$, and $\Gamma_{\mathcal{M}, n}$ is regular.

There is a natural interpretation of the n -th Schreier graph of \mathcal{M} in terms of *products* of automata, which is slightly beyond the scope of this paper: the Schreier graph $\Gamma_{\mathcal{M}, n}$ is the Moore diagram of the n -fold *product* of the dual $\overline{\mathcal{M}}$ with itself (see e.g. [DR]).

Example 3.2. The n -th Bellaterra graph B_n is the Schreier graph $\Gamma_{\mathcal{B}, n}$, where \mathcal{B} is the Bellaterra automaton, pictured in Figure 3.

For example, the inductive definition of B_n in the introduction has the rule that

$$v \xrightarrow{b} w \quad \text{lifts to} \quad \begin{array}{c} v_0 \xrightarrow{a} w_0 \\ v_1 \xrightarrow{b} w_1 \end{array}.$$

In the definition in terms of automata, this corresponds to the fact that, if $v, w \in \{0, 1\}^n$ are binary words with ${}^b v = w$, then

$${}^a(0v) = 0{}^b v = 0w \quad \text{and} \quad {}^b(1v) = 0{}^b v = 1w.$$

Example 3.3. Similarly, the n -th Aleshin graph A_n is the Schreier graph $\Gamma_{\mathcal{A}, n}$, where \mathcal{A} is the Aleshin automaton, first considered in [A], pictured in Figure 5.

4. AUTOMATON GROUPS

Let $\mathcal{M} = (Q, A, \tau, \sigma)$ be an invertible Mealy automaton. As seen in Section 2, we have invertible maps $\sigma_q : A^{*, \infty} \rightarrow A^{*, \infty}$ for each $q \in Q$. This gives an action of the free group F_Q on $A^{*, \infty}$. We can extend the definition of σ as follows: For $w \in F_Q$, we can define σ_w in the natural way, e.g.,

$$\sigma_{qr^{-1}} = \sigma_q \sigma_r^{-1}.$$

As usual, we will adopt the notational convention

$$\sigma_w(s) = \sigma(w, s) = {}^w s.$$

The *automaton group* associated to \mathcal{M} is the group $G_{\mathcal{M}}$ generated by the automorphisms σ_q . For example, letting \mathcal{A} denote the Aleshin automaton, it was shown in [VV1] that $\sigma_{\mathcal{A},a}$, $\sigma_{\mathcal{A},b}$, and $\sigma_{\mathcal{A},c}$ satisfy no nontrivial relation, so $G_{\mathcal{A}}$ is the free group F_3 . However, it is straightforward to check that $\sigma_{\mathcal{B},a}^2 = \sigma_{\mathcal{B},b}^2 = \sigma_{\mathcal{B},c}^2 = \text{id}$, where \mathcal{B} is the Bellaterra automaton. It is known that $G_{\mathcal{B}} \cong \langle a, b, c \mid a^2, b^2, c^2 \rangle = C_2 * C_2 * C_2$, i.e. the words above satisfy no other relations. This was shown by E. Muntz and D. Savchuk, and the proof is included in [Nek] (see also [B+]).

Information about $G_{\mathcal{M}}$ as an abstract group can be used to obtain information about the Schreier graphs $\Gamma_{\mathcal{M},n}$. See Remarks 11.8 and 11.7.

5. TREES AND AUTOMORPHISMS

In this context it is natural to think of the set of finite words A^* as vertices in a regular rooted tree, where the empty word is the root and the children of the word s are the words sx for $x \in A$. We will need to talk about rooted trees more generally, so we make the following definitions.

A *rooted tree* (or simply *tree*) is a graph \mathbf{T} with a distinguished root vertex $r \in \mathbf{T}$ such that for each $v \in \mathbf{T}$ there is exactly one directed path from r to v . The *level* of v , denoted $\ell(v)$ is the length of this path. The *n -th level* of \mathbf{T} , denoted \mathbf{T}_n is the set of all vertices $v \in \mathbf{T}$ such that $\ell(v) = n$. A *subtree* of \mathbf{T} is a subgraph containing r which is itself a rooted tree. A *tree isomorphism* between two trees \mathbf{S} and \mathbf{T} is a graph isomorphism which sends the root of \mathbf{S} to the root of \mathbf{T} . An *automorphism* of \mathbf{T} is an isomorphism from \mathbf{T} to \mathbf{T} . Let $\text{Aut}(\mathbf{T})$ be the group of automorphisms of \mathbf{T} .

Then, given a Mealy automaton $\mathcal{M} = (Q, A, \sigma, \tau)$, for any $q \in Q$ the map $\sigma_q : A^* \rightarrow A^*$ is a tree automorphism. That is, σ_q is a bijection which fixes the empty word, and sends children of x to children of $\sigma_q(x)$. In other words, for every $s \in A^*$ and $x \in A$ there is some $y \in A$ such that

$$\sigma_q(sx) = \sigma_q(s)y.$$

Of course, it follows that for any $w \in Q^*$, the map $\sigma_w : A^* \rightarrow A^*$ is a composition of tree automorphisms and is itself a tree automorphism. Infinite words, i.e., elements of A^∞ , can be thought of as rays from the root of the tree A^* , and σ_w acts on them in the natural way.

Note that in order to think of $\tau_a : Q^* \rightarrow Q^*$ as a tree automorphism, we must think of Q^* as a tree in the reverse way, i.e. the children of w are of the form qw for $q \in Q$, rather than of the form wq . Given a tree automorphism $g : A^* \rightarrow A^*$ and a word $s \in A^*$, the *section* of g at s , is the tree automorphism $g|_s : A^* \rightarrow A^*$ defined by $g(st) = g(s)g|_s(t)$. Note that we are using a canonical identification between branches of the tree A^* . There need not be such an identification in a general tree, so this definition of sections is specific to trees of words.

We call a tree automorphism $\alpha : A^* \rightarrow A^*$ *automatic* if it $\alpha = \sigma_{\mathcal{M},q}$ for state q of some Mealy automaton \mathcal{M} . Equivalently, α is automatic if and only if it has finitely many distinct sections. The set of automatic automorphisms forms a subgroup $\text{FAut}(A^*) < \text{Aut}(A^*)$.

An automorphism $g : A^* \rightarrow A^*$ is determined by its action on the first level, $(A^*)_1 = A^1 = A$, and its sections $g|_x$ at all $x \in A$. If $\rho : A \rightarrow A$ is a permutation, then for notational convenience we can extend ρ as an automorphism $A^* \rightarrow A^*$ via $\rho(xs) = \rho(x)s$, for all $x \in A$, $s \in A^*$. In other words, automorphism ρ acts only on the first letter of a word in A^* . If A is equipped with an ordering of its elements, say, $A = \{x_1, \dots, x_k\}$, then we write (g_1, \dots, g_k) for the automorphism $g : A^* \rightarrow A^*$ which acts trivially on A , and whose sections are $g|_{x_i} = g_i$ for all i . Then every automorphism can be uniquely decomposed into $g = \rho(g_1, \dots, g_k)$, for some permutation $\rho : A \rightarrow A$ and some automorphisms $g_i : A^* \rightarrow A^*$. Specifically, ρ is the restriction of g to A , and $g_i = g|_{x_i}$. Then, given an invertible Mealy automaton $\mathcal{M} = (Q, A, \tau, \sigma)$, the definition of the automorphisms σ_q can be phrased recursively as

$$\sigma_q = \sigma_q(\sigma_{q_1}, \dots, \sigma_{q_k}),$$

where $q_i = q^{x_i}$. Such a recursive definition is called a *wreath recursion*. For example, if $\mathcal{B} = (Q, A, \tau, \sigma)$ is the Bellaterra automaton, then we have the wreath recursion

$$\sigma_a = (\sigma_b, \sigma_c), \quad \sigma_b = (\sigma_c, \sigma_b), \quad \sigma_c = \rho(\sigma_a, \sigma_a),$$

where $\rho : \{0, 1\} \rightarrow \{0, 1\}$ swaps 0 and 1.

Finally, we say a tree automorphism $g : \mathbf{T} \rightarrow \mathbf{T}$ is *spherically transitive* (or just *transitive*) if its restriction to every level of \mathbf{T} is a transitive map. For example, if \mathcal{M} is the adding automaton pictured in Figure 4, then $\sigma_r : \{0, 1\}^* \rightarrow \{0, 1\}^*$ is spherically transitive, because its action on the n -th level is addition of 1 modulo 2^n .

6. THE BELLATERRA AUTOMATON

Consider the Bellaterra automaton $\mathcal{B} = (Q, A, \sigma, \tau)$, pictured in Figure 3. We want to show that the graphs $B_n = \Gamma_{\mathcal{B}, n}$ have small diameter. Our approach is to construct a short path from every binary string in $\{0, 1\}^n$ to the string $1^n = 111 \dots 1$.

We do this by repeatedly increasing how many times the digit 1 occurs at the beginning of the string. To achieve this we must find words in Q^* which send $1^{k-1}0$ to 1^k , that is, words whose action on 1^k changes only the last digit. So, we are looking for words in Q^* which do not fix the infinite word $1^\infty = 111 \dots$, but do preserve the first k of its letters. There are enough words which do not fix 1^∞ because τ_1 acts “transitively enough” on Q^* , so that almost every orbit under its action contains some word $w \in Q^*$ which swaps 0 and 1. Then we can select words from this set which fix 1^k using the pigeonhole principle.

The following lemma established the relationship between words which fix 1^∞ and orbits of τ_1 . This result is known (see e.g. [AR, Thm. 2]); we include a short proof for completeness.

Lemma 6.1. *Let $\mathcal{M} = (Q, A, \tau, \sigma)$ be a Mealy automaton, let x be a letter in A , and let w be a word in Q^* . Then w fixes the infinite word $xxx \dots = x^\infty$ if and only if every element of the orbit of w under τ_x fixes x . That is,*

$${}^w(xxx \dots) = xxx \dots \quad \text{if and only if} \quad \sigma(\tau_x^n(w), x) = x \text{ for all } n \geq 0.$$

Proof. Say

$${}^w(xxx \dots) = y_0 y_1 y_2 \dots$$

Then y_n is the last letter of ${}^w(x^{n+1})$. Letting $X = x^n$, we have

$${}^w(x^{n+1}) = {}^w(Xx) = {}^w(X) \tau^{(w, X)}x,$$

So,

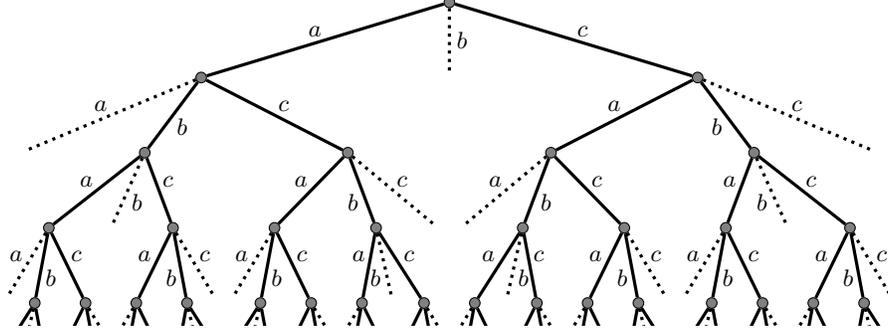
$$y_n = \tau^{(w, X)}x = \sigma(\tau(w, X), x) = \sigma(\tau_X(w), x) = \sigma(\tau_x^n(w), x),$$

and therefore ${}^w xxx \dots = xxx \dots$ if and only if $\sigma(\tau_x^n(w), x) = x$ for every $n \geq 0$, as desired. \square

Since we wish to find many words in Q^* which do not fix 1^∞ , we want to find large orbits of τ_1 . It is straightforward to check that a^2 , b^2 , and c^2 act trivially on A^* , (i.e. $\sigma_{aa} = \sigma_{bb} = \sigma_{cc} = \text{id}$) so we are primarily interested in *reduced words* in $\{a, b, c\}$, i.e. those which do not repeat the same letter twice in a row. Note that these words form a subtree of Q^* , which is nearly a binary tree: every vertex has two children, except the root. In fact, we will work with the binary subtree pictured in Figure 6, consisting of those reduced words in Q^* which end with a or c .

We will need a simple result on the transitivity of automorphisms of a binary tree. Let $A = \{0, 1\}$. Define a group homomorphism $\chi : \text{Aut}(A^*) \rightarrow \mathbb{Z}_2[[t]]$, by

$$\chi(g) = \sum_{n=1}^{\infty} c_n t^{n-1},$$

FIGURE 6. The subtree \mathbf{T} of Q^* .

where $(-1)^{c_n}$ is the sign of the permutation given by the action of g the n -th level of A^* . Values of this homomorphism can be computed recursively via

$$\chi(g) = c_1 + t(\chi(g|_x) + \chi(g|_y)),$$

where c_1 is 0 if g fixes the two elements of A , and $c_1 = 1$ if g swaps them. We call $\chi(g)$ the *characteristic function of g* . Of course, this definition makes sense when A is any two-element set, so we will state the lemma more generally:

Lemma 6.2. *Let $A = \{x, y\}$. An automorphism $g \in \text{Aut}(A^*)$ is spherically transitive if and only if $\chi(g) = 1/(1-t)$.*

Proof. If g is spherically transitive, then its action on the n -th level of A^n is a (2^n) -cycle, which is an odd permutation for all $n \geq 1$. Hence, $c_n = 1$ for all $n \geq 1$, and

$$\chi(g) = \sum_{n=1}^{\infty} t^{n-1} = \frac{1}{1-t}.$$

In the other direction, suppose $\chi(g) = 1/(1-t)$, i.e. g acts as an odd permutation on A^n for every $n \geq 1$. We will show by induction on n that the action of g on A^n is a (2^n) -cycle for all $n \geq 0$. This is trivial for $n = 0$.

For the inductive step, suppose g acts as a (2^n) -cycle on A^n . Given a word $s \in A^n$, we either have $g^{2^n}(sx) = sx$ or $g^{2^n}(sx) = sy$. In the first case sx belongs to a (2^n) -cycle of g , in the second case, sx belongs to a (2^{n+1}) -cycle. So, any word in A^{n+1} ending in x belongs to either a (2^n) -cycle or a (2^{n+1}) -cycle, and similarly for words ending in y . So, the action of g on A^{n+1} decomposes into either two (2^n) -cycles or a single (2^{n+1}) -cycle. But the former is an even permutation, so g must act as a (2^{n+1}) -cycle on A^{n+1} , as desired. \square

Lemma 6.3. *Let $\mathcal{B} = (Q, A, \tau, \sigma)$ denote the Bellaterra automaton. Then for every natural number n , the map τ_1 acts transitively on the set of reduced words of length n ending with a or c .*

Proof. We will write the argument down in terms of the dual automaton $\overline{\mathcal{B}}$, pictured in Figure 3. This switches the role of τ and σ . Taking the dual also reverses words, so we want to show that $\sigma_1 = \sigma_{\overline{\mathcal{B}}, 1}$ acts transitively on the binary subtree $\mathbf{T} \subset Q^*$ of reduced words which *begin* with a or c . This tree is pictured in Figure 6.

The binary tree \mathbf{T} is not a tree of words over a two-letter alphabet, so it is convenient to put it into bijection with one. Let \mathbf{R} be the binary tree $\{\uparrow, \downarrow\}^*$. We define the maps $\phi_a, \phi_b, \phi_c : \mathbf{R} \rightarrow Q^*$ recursively by

$$\begin{aligned} \phi_a(\uparrow w) &= b\phi_b(w) & \phi_a(\downarrow w) &= c\phi_c(w) \\ \phi_b(\uparrow w) &= c\phi_c(w) & \phi_b(\downarrow w) &= a\phi_a(w) \\ \phi_c(\uparrow w) &= a\phi_a(w) & \phi_c(\downarrow w) &= b\phi_b(w) \end{aligned}$$

It is straightforward to check by induction on word length that for each $x \in Q$, the map ϕ_x defines a tree isomorphism between \mathbf{R} and the tree of reduced words in Q^* which do not begin with x . In particular, ϕ_b is a bijection between \mathbf{R} and \mathbf{T} .

Now consider the dual $\bar{\mathbf{B}}$ of the Bellaterra automaton, and in particular the corresponding automorphisms $\sigma_0, \sigma_1 \in \text{Aut}(Q^*)$. Given $x, y \in Q$, $d \in A$, define

$$\sigma_{x,d,y} = \phi_x^{-1} \sigma_d \phi_y \in \text{Aut}(\mathbf{R}).$$

Note that, a priori, the domain of ϕ_x^{-1} may not coincide with the image of $\sigma_d \phi_y$, so $\sigma_{x,d,y}$ may be ill-defined for some values of x, d, y . However, the computations below give an explicit recursion for computing $\sigma_{1,b,1}$, which also demonstrates that it is well-defined.

We can compute that, e.g.,

$$\begin{aligned} \sigma_{b,1,b}(\uparrow w) &= \phi_b^{-1}(\sigma_1(\phi_b(\uparrow w))) \\ &= \phi_b^{-1}(\sigma_1(c \phi_c(w))) \\ &= \phi_b^{-1}(a \sigma_0(\phi_c(w))) \\ &= \downarrow \phi_a^{-1}(\sigma_0(\phi_c(w))) \\ &= \downarrow \sigma_{a,0,c} \end{aligned}$$

In particular, $\sigma_{b,1,b}|_{\uparrow} = \sigma_{a,0,c}$.

Similar computations give the complete recursive description of $\sigma_{b,1,b}$, which we write down using the usual wreath recursion notation $g = \rho^\varepsilon(g|_{\uparrow}, g|_{\downarrow})$, where ρ swaps \uparrow and \downarrow :

$$\begin{aligned} \sigma_{b,1,b} &= \rho(\sigma_{a,0,c}, \sigma_{c,1,a}) \\ \sigma_{a,0,c} &= (\sigma_{b,0,a}, \sigma_{c,0,b}) \\ \sigma_{c,1,a} &= \rho(\sigma_{b,1,b}, \sigma_{a,0,c}) \\ \sigma_{b,0,a} &= (\sigma_{c,0,b}, \sigma_{a,1,c}) \\ \sigma_{c,0,b} &= (\sigma_{a,1,c}, \sigma_{b,0,a}) \\ \sigma_{a,1,c} &= \rho(\sigma_{c,1,a}, \sigma_{b,1,b}) \end{aligned}$$

Defining $F_{x,d,y} = \chi(\sigma_{x,d,y})$, this gives us the following linear equations in the ring $\mathbb{Z}_2[[t]]$:

$$\begin{aligned} F_{b,1,b} &= 1 + t(F_{a,0,c} + F_{c,1,a}) \\ F_{a,0,c} &= t(F_{b,0,a} + F_{c,0,b}) \\ F_{c,1,a} &= 1 + t(F_{b,1,b} + F_{a,0,c}) \\ F_{b,0,a} &= t(F_{c,0,b} + F_{a,1,c}) \\ F_{c,0,b} &= t(F_{a,1,c} + F_{b,0,a}) \\ F_{a,1,c} &= 1 + t(F_{c,1,a} + F_{b,1,b}) \end{aligned}$$

Solving this system of equations yields

$$\begin{aligned} F_{b,1,b} &= 1/(1-t) \\ F_{a,0,c} &= 0 \\ F_{c,1,a} &= 1/(1-t) \\ F_{b,0,a} &= t/(1-t) \\ F_{c,0,b} &= t/(1-t) \\ F_{a,1,c} &= 1, \end{aligned}$$

So we have

$$\chi(\phi_b^{-1} \sigma_{\bar{\mathbf{B}},1} \phi_b) = 1/(1-t).$$

By Lemma 6.2, the automorphism $\phi_b^{-1} \sigma_1 \phi_b$ acts transitively on \mathbf{R} . Hence σ_1 acts transitively on \mathbf{T} . That is, for each n it acts transitively on the set of length n reduced words in $\{a, b, c\}$

which begin with a or c . Hence, in the unreversed Bellaterra automaton \mathcal{B} , we have that $\tau_{\mathcal{B},1}$ acts transitively on the words of a given length which end with a or c . \square

We can now put the argument together to show that B_n has at most quadratic diameter.

Theorem 6.4. *Let B_n denote the n -th Bellaterra graph. Then $\text{diam}(B_n) = O(n^2)$.*

Proof. Let $\mathcal{B} = (Q, A, \tau, \sigma)$ denote the Bellaterra automaton, so that $B_n = \Gamma_{\mathcal{B},n}$. It is enough to show that for some C the ball of radius Cn^2 around the vertex $1^n = 11\dots 1$ covers all of B_n . That is, we will show that for every number n , and every $v \in \mathcal{B}_n$,

$$d(1^n, v) \leq Cn^2.$$

The only letter in $Q = \{a, b, c\}$ which swaps the elements of A is c . The other two letters fix 0 and 1. Hence, a word $w \in Q^*$ fixes 1 if and only if it has an even number of c 's.

For each $n > 0$, there is a reduced word ending in a or c which contains an odd number of c 's. We can take, e.g. $abab\dots abc$ or $baba\dots abc$. By Lemma 6.3, if w is any reduced word of length n ending in a or c , then its orbit under τ_1 contains some word which does not fix 1. Hence, by Lemma 6.1, w does not fix the infinite word $111\dots = 1^\infty$.

Given a number $n \geq 1$, we have $|A^n| = 2^n$, and there are $2^{n+1} - 1$ reduced words of length n or less which end in a or c . By the pigeonhole principle, there must be two such words, v, w with $v(1^n) = w(1^n)$. We may assume $\ell(v) \leq \ell(w)$. Since a^2, b^2 , and c^2 all act trivially on A^* , reversing a word inverts its action on A^* . Let u be the reduced word formed by canceling pairs of repeated letters in $\bar{v}w$. Then,

$$\begin{aligned} u(1^n) &= \bar{v}w(1^n) = \bar{v}v(1^n) = 1^n, \\ \text{and } \ell(u) &\leq \ell(\bar{v}w) = \ell(v) + \ell(w) \leq 2n. \end{aligned}$$

Since $v \neq w$, u is not the empty word. We assumed that $\ell(v) \leq \ell(w)$, so the last letter of w is not canceled. Hence u also ends in a or c , and therefore

$$u(111\dots) \neq 111\dots$$

Let k be the maximal integer such that $u(1^k) = 1^k$. We know $k \geq n$ and $u(1^{k+1}) = 1^k 0$. So, letting $s = 1^{k-n}$, $t = 1^{n+1}$, and $t' = 1^n 0$, we have

$$st' = u(st) = u_s u'(t) = s u'(t),$$

where $u' = u^s$. So we have

$$\begin{aligned} u'(1^{n+1}) &= 1^n 0, \\ \text{and } \ell(u') &= \ell(u) \leq 2n. \end{aligned}$$

This construction works for all $n \geq 1$. That is, for every $n \geq 1$, there exists a $u_n \in Q^*$ with $\ell(u_n) \leq 2n$ and $u_n(1^n 0) = 1^{n+1}$.

We now prove by induction on n that for every $s \in A^n$, there is a $w \in Q^*$ with $\ell(w) \leq n^2$ such that $w_s = 1^n$. The base cases $n = 0$ and $n = 1$ are trivial. For the inductive step, consider any $n \geq 1$. Given $s \in A^{n+1}$, let s' be s with the last digit removed. By the induction hypothesis know there is a word w with $\ell(w) \leq n^2$ such that $w_{s'} = 1^n$. Then either $w_s = 1^{n+1}$ or $w_s = 1^n 0$. In the first case, we are done. In the second case, $u_n w_s = 1^{n+1}$, and $\ell(u_n w) \leq 2n + n^2 \leq (n+1)^2$, so we are done.

So, we have shown that in the graph $B_n = \Gamma_{\mathcal{B},n}$, we have $d(1^n, s) \leq n^2$ for every $s \in A^n$. It follows that for any $s, t \in B_n$,

$$d(s, t) \leq d(s, 1^n) + d(1^n, t) \leq 2n^2,$$

i.e. $\text{diam}(B_n) \leq 2n^2$. \square

7. THE ALESHIN AUTOMATON

The Aleshin automaton \mathcal{A} and the Bellaterra automaton \mathcal{B} are closely related. Indeed, let $\tau_d : \{0, 1\}^* \rightarrow \{0, 1\}^*$ denote the map which swaps every digit of a binary word. Then it is straightforward to check by induction that

$$\begin{aligned}\tau_{\mathcal{A},a} &= \tau_d \tau_{\mathcal{B},a}, \\ \tau_{\mathcal{A},b} &= \tau_d \tau_{\mathcal{B},c}, \\ \text{and } \tau_{\mathcal{A},c} &= \tau_d \tau_{\mathcal{B},b}.\end{aligned}$$

With this observation, Theorem 6.4 has the following corollary.

Corollary 7.1. *Let A_n denote the n -th Aleshin graph. Then $\text{diam}(A_n) = O(n^2)$.*

Proof. For every pair $q, r \in \{a, b, c\}$, we have

$$\tau_{\mathcal{A},q}^{-1} \tau_{\mathcal{A},r} = \tau_{\mathcal{B},q}^{-1} \tau_d^{-1} \tau_d \tau_{\mathcal{B},r} = \tau_{\mathcal{B},q} \tau_{\mathcal{B},r}.$$

So, if two words in $\{0, 1\}^n$ are separated by a path of length 2 in the Bellaterra graph B_n , they are also separated by a path of length 2 in the Aleshin graph A_n . It follows that two endpoints of an even-length path in B_n are endpoints of a path in A_n of the same length.

For any word $s \in \{0, 1\}^n$ there is a path in B_n of length $O(n^2)$ from 1^n to s . We may assume that this path has even length since 1^n has an edge in B_n from itself to itself. This corresponds to a path in A_n of the same length, so for any $s \in \{0, 1\}^n$, there is a path in A_n of length $O(n^2)$ from 1^n to s . Therefore, $\text{diam}(\Gamma_{\mathcal{A},n}) = O(n^2)$, as desired. \square

8. GENERALIZATIONS

The proof of Theorem 6.4 can be adapted to prove a more general result. In order to generalize to automata with larger alphabets, we need to consider a restricted type of automaton. We say an Mealy automaton $\mathcal{M} = (Q, A, \tau, \sigma)$ is *cyclic* if it is invertible and the group generated by the output maps σ_q is a cyclic group of order $|A|$. In other words we can write $A = \{x_1, x_2, \dots, x_n\}$, so that $\langle \sigma_q \mid q \in Q \rangle = \langle (x_1 x_2 \dots x_n) \rangle$. In particular, any invertible automaton with $|A| = 2$ is cyclic, unless all the maps σ_q are trivial. We will also restrict attention to those cyclic automata with $|A|$ prime. This is necessary to guarantee that a word in Q^* which acts on A nontrivially cycles through all elements of A . This will enable us to reach any word of the form $x^n y$ from x^{n+1} in a short time, as long as we can reach some such word with $y \neq x$.

We first state and prove the general result with the weakest assumptions under which our argument guarantees polynomial growth of $\text{diam}(\Gamma_{\mathcal{M},n})$.

Theorem 8.1. *Let $\mathcal{M} = (Q, A, \sigma, \tau)$ be a cyclic Mealy automaton with $|A|$ prime, and let $\Gamma = \Gamma_{\mathcal{M},\infty}$. Suppose there is a letter $x \in A$ and constants $\alpha > 0$, $K > 1$ such that, for sufficiently large r ,*

$$|B_\Gamma(xxx \dots, r)| \geq K^{r^\alpha}.$$

Then there is a constant $C > 0$, such that for all n ,

$$\text{diam}(\Gamma_{\mathcal{M},n}) \leq Cn^{1+1/\alpha}.$$

Proof. Let $p = |A|$. By replacing \mathcal{M} with $\mathcal{M} \cup \mathcal{M}^{-1}$ if necessary, we may assume that for every $q \in Q$, there is a $q' \in Q$ with $\sigma_{q'} = \sigma_q^{-1}$. This replacement adds edges to the Schreier graphs $\Gamma_{\mathcal{M},n}$, but only reverses of edges which were already there, so $\text{diam}(\Gamma_{\mathcal{M},n})$ and $B_{\Gamma_{\mathcal{M},n}}(s, r)$ are unaffected. Then for a word $w \in Q^*$, we define w^{-1} to be w , reversed, with each letter q replaced by q' , so that $\sigma_{w^{-1}} = \sigma_w^{-1}$.

Given sufficiently large n , pick r such that

$$((\log_K p)n)^{1/\alpha} < r < 2((\log_K p)n)^{1/\alpha}.$$

Then $|B_\Gamma(xxx\dots, r)| > p^n$. By the pigeonhole principle, some two elements of $B_\Gamma(xxx\dots, r)$ have the same first n digits. That is, there are $v, w \in (Q \cup Q^{-1})^*$ with

$$\ell(v), \ell(w) \leq r, \quad v(x^n) = w(x^n), \quad \text{and} \quad v(xxx\dots) \neq w(xxx\dots).$$

So, there is a $u_0 = v^{-1}w \in (Q \cup Q^{-1})^*$ with

$$\begin{aligned} \ell(u_0) &\leq 2r, \\ u_0(x^n) &= x^n, \\ \text{and} \quad u_0(xxx\dots) &\neq xxx\dots \end{aligned}$$

There is some smallest value of $k \geq n+1$ such that $u_0(x^k) \neq x^k$. Let $X_0 = x^{k-n-1}$ and $X = x^n$, so that $x^k = X_0Xx$ and $u_0(x^k) = X_0Xy$ for some $y \in A$ with $y \neq x$. Let $u = u_0^{X_0}$ so in particular, $\ell(u) = \ell(u_0)$. Then,

$$X_0Xy = u_0(X_0Xx) = u_0X_0 \ u(Xx) = X_0 \ u(Xx),$$

so

$$u(Xx) = Xy.$$

Similarly, if $u' = u^X$, we have

$$u(Xz) = X \ u'z$$

for any $z \in A$. Since $u'x = y \neq x$ and \mathcal{M} is cyclic, the action of u' on A is a nontrivial cyclic permutation. Since $p = |A|$ is prime, u' acts transitively on A , and therefore u acts transitively on $\{Xz \mid z \in A\}$. It follows that for any $z, z' \in A$,

$$d(Xz, Xz') \leq p \ell(u) \leq 2pr \leq 4p((\log_K p)n)^{1/\alpha}.$$

Thus, there is a constant C such that for sufficiently large n , we have

$$d(x^n z, x^{n+1}) \leq Cn^{1/\alpha}, \quad \text{for all } z \in A.$$

By increasing the constant if necessary, we can make this true for all n .

Now let us show by induction on n that for all $s \in A^n$, we have $d(s, x^n) < Cn^{1+1/\alpha}$. The base case $n = 0$ is trivial. For the inductive step, take any $s \in A^{n+1}$, and let s' be its first n letters. We know $d(s', x^n) < Cn^{1+1/\alpha}$. There is some word $w \in (Q \cup Q^{-1})^*$ with $\ell(w) = d(s', x^n)$ and $w(s') = x^n$. Then $w_s = x^n z$ for some $z \in B$. Thus,

$$\begin{aligned} d(s, x^{n+1}) &\leq d(s, x^n z) + d(x^n z, x^{n+1}) \\ &\leq Cn^{1+1/\alpha} + Cn^{1/\alpha} \\ &\leq C(n+1)^{1+1/\alpha}, \end{aligned}$$

which completes the induction.

It follows that for any $s, t \in A^n$, $d(s, t) \leq d(s, x^n) + d(x^n, t) \leq 2Cn^{1+1/\alpha}$, i.e.,

$$\text{diam}(\Gamma_{\mathcal{M}, n}) \leq 2Cn^{1+1/\alpha}.$$

□

In all the cases where we apply this, $|B_\Gamma(xxx\dots, r)|$ will have exponential growth, so we state that case separately.

Corollary 8.2. *Let $\mathcal{M} = (Q, A, \sigma, \tau)$ be a cyclic Mealy automaton with $|A|$ prime. Let $\Gamma = \Gamma_{\mathcal{M}, \infty}$. If there is an $x \in A$ and a constant $K > 1$ such that*

$$|B_\Gamma(xxx\dots, r)| \geq K^r$$

for sufficiently large r , then there is a constant $C > 0$, such that

$$\text{diam}(\Gamma_{\mathcal{M}, n}) \leq Cn^2.$$

It is not always easy to guarantee that $|B_\Gamma(xxx\dots, r)| \geq K^r$ grows quickly, so we prove an additional result based on the size of orbits of τ_x in Q^n . Loosely, if the orbits grow quickly enough, it must be because there are enough distinct images of words of the form x^m .

Theorem 8.3. *Let $\mathcal{M} = (Q, A, \sigma, \tau)$ be a reversible³ cyclic Mealy automaton with $|A|$ prime. Suppose there is a letter $x \in A$ and constants $K > 1$, $\alpha > 0$ such that for sufficiently large n , there is a $w \in Q^n$ with*

$$|\{\tau_x^k(w) \mid k \in \mathbb{Z}\}| \geq K^{n^\alpha}.$$

Then there is a constant $C > 0$, such that for all n , we have

$$\text{diam}(\Gamma_{\mathcal{M},n}) \leq C n^{1+1/\alpha}.$$

Proof. Let $P = \{p \in \mathbb{N} \mid p \text{ prime}, p \leq |Q|\}$. It is easy to see by induction on length that for each $w \in Q^*$, the sequence $w, \tau_x(w), \tau_x^2(w), \dots$ is periodic with period m , where m is a product of some powers of primes in P . Define letters $q_{i,k}$ via

$$\tau_x^k(w) = q_{0,k}q_{1,k} \dots q_{n,k}.$$

That is, $q_{i,k}$ is the i -th letter of $\tau_x^k(w)$. Let m_i be the period of the sequence $q_{i,0}, q_{i,1}, \dots$. Then the period of $\tau_x^0(w), \tau_x^1(w), \dots$ is $m = \gcd(m_0, m_1, \dots, m_n)$. Let $M = \max_i m_i$. Each prime power in the prime factorization of m is a factor of some m_i , so it is at most M . The period m is the product of these prime powers, so $m \leq M^{|P|}$. That is, there is some i such that $m_i \geq m^{1/|P|}$.

Fix that i for the rest of the proof, and let v be the first i letters of w . Consider the infinite word $s = {}^v(xxx\dots)$, and let x_k be its k -th letter. Note that s is periodic, and let l be the period of the word s , i.e. the smallest number for which $x_{k+l} = x_k$ for all k . Note that $q_{i,k+1} = q_{i,k}^{x_k}$ and therefore

$$q_{i,k+l} = q_{i,k}^X,$$

where

$$X = x_k x_{k+1} \dots x_{k+l-1}.$$

Since the x_k repeat every l letters, we have $q_{i,k+l} = q_{i,k}^X$, and $q_{i,k+2l} = q_{i,k+l}^X$, and so on. Let $F = |Q|^l$. Then X^F acts trivially on Q , and hence $q_{i,k+lF} = q_{i,k}$. This is true for each k , so the $q_{i,k}$ have period $m_i \leq Fl$. Thus, the word $s = {}^v(xxx\dots)$ has period $l \geq m^{1/|P|}/F$.

Now let n be sufficiently large, so that there is a word $w \in Q^n$ whose orbit under τ_x has size $m \geq K^{n^\alpha}$. Then, from the above, for some $v \in Q^*$ with $\ell(v) \leq n$, the word $s = {}^v(xxx\dots)$ has period

$$l \geq \frac{1}{F} K^{n^\alpha}/|P| \geq \tilde{K}^{n^\alpha},$$

where we fix some $1 < \tilde{K} < K^{1/|P|}$, and the last inequality holds for sufficiently large n .

Let $v_k = \tau_x^k(v)$. Then ${}^{v_k}(xxx\dots)$ is a shift of s , and since s has period l there are l distinct such shifts. So, since each v_k satisfies $\ell(v_k) \leq n$, the set $\{{}^w(xxx\dots) \mid w \in Q^*, \ell(w) \leq n\}$ has at least $l \geq \tilde{K}^{n^\alpha}$ elements. It follows that $|B_\Gamma(xxx\dots, n)| \geq \tilde{K}^{n^\alpha}$, where $\Gamma = \Gamma_{\mathcal{M},\infty}$.

So Theorem 8.1 applies, and there is a constant C such that $\text{diam}(\Gamma_{\mathcal{M},n}) \leq C n^{1+1/\alpha}$. \square

We also state the following special case, which is a simple way to apply the theorem.

Corollary 8.4. *Let $\mathcal{M} = (Q, A, \sigma, \tau)$ be a reversible cyclic Mealy automaton with $|A|$ prime. Suppose there is some $a \in A$, and some $d \geq 2$, such that τ_a acts spherically transitively on some d -regular subtree $\mathbf{T} \subseteq Q^*$. Then there is a constant $C > 0$, such that for all n , we have*

$$\text{diam}(\Gamma_{\mathcal{M},n}) \leq C n^2.$$

³The assumption that \mathcal{M} is reversible may be lifted, if we replace $|\{\tau_x^k(w) \mid k \in \mathbb{Z}\}|$ with the length of the (eventual) period of $w, \tau_x(w), \tau_x^2(w), \dots$.

9. ELEMENT-COTRANSITIVE CYCLIC AUTOMATA

The simplest way for the conditions in Corollary 8.4 to be satisfied is when some τ_a acts spherically transitively on the entire tree Q^* . With that in mind, we make the following definitions.

We say an invertible Mealy automaton $\mathcal{M} = (Q, A, \sigma, \tau)$ is *q-transitive* if the tree automorphism $\sigma_q : A^* \rightarrow A^*$ is spherically transitive. We say \mathcal{M} is *element-transitive* if it is *q-transitive* for some $q \in Q$. We say \mathcal{M} is *element-cotransitive* if its dual is element-transitive. In particular, we require an element-cotransitive automaton to be reversible, but it need not be bireversible.

Then, according to Corollary 8.4, we have

Corollary 9.1. *Let $\mathcal{M} = (Q, A, \sigma, \tau)$ be a cyclic element-cotransitive Mealy automaton with $|A|$ prime. Then $\text{diam}(\Gamma_{\mathcal{M}, n}) = O(n^2)$.*

We do not know a general method for determining whether a tree automorphism given by an automaton is element-transitive, but there are special cases where checking it is easier. For example, [St] gives a generalization of Lemma 6.2 to all cyclic automata:

Lemma 9.2. *Let $\mathcal{M} = (Q, A, \tau, \sigma)$ be a cyclic automaton, with $|A| = m$. Then there is a cyclic permutation ρ of A , such that for each $q \in Q$ there is a k_q s.t. $\sigma_q = \rho^{k_q}$. Recursively define*

$$\chi(q) = k_q + t \sum_{x \in A} \chi(\tau_x(q)) \in \mathbb{Z}_m[[t]]$$

Then σ_q acts transitively on A^ if and only if each coefficient of $\chi(q)$ is a generator of \mathbb{Z}_m .*

An automaton is called *cocyclic* if its dual is cyclic. Now observe that the power series $\chi(q)$ for $q \in Q$ satisfy a recursive linear relation, which can be solved to write each $\chi(q)$ as a rational function. This implies:

Corollary 9.3. *Given a (co)cyclic Mealy automaton $\mathcal{M} = (Q, A, \tau, \sigma)$, there is an algorithm to determine whether it is element-(co)transitive.*

For example, it is straightforward to check that, there are 16 cocyclic invertible (3, 2)-automata, and only four are element-cotransitive. These four are the automata pictured in Figures 7b–7e, i.e., automata number 956, 2396, 870, and 2294 in [B+].⁴

10. FURTHER EXAMPLES

Below we give a brief outline of a few more example of interesting Mealy automata. We intentionally suppress some of the proofs which are often straightforward and occasionally involve simple computer checking.

Example 10.1. The automata in Figure 7 are precisely the invertible (3, 2)-automata which are element-cotransitive. They are labeled by their numbers in [B+].

The fact that no other (3, 2) automata are element-cotransitive can be verified by a computation. For each automaton, it is enough to look at Q^4 . That is, each invertible (3, 2)-automaton $\mathcal{M} = (Q, A, \sigma, \tau)$, except the ones in Figure 7, has no $a \in A$ for which τ_a does acts transitively on Q^4 .

The automata in Figures 7b–7e are cocyclic, so we can use Lemma 9.2 to verify that they are indeed element-cotransitive. The remaining automaton \mathcal{M} , pictured in Figure 7a, is not cocyclic, so we cannot apply Lemma 9.2 directly. However, it turns out that there is an automatic automorphism $\kappa : \{a, b, c\}^* \rightarrow \{a, b, c\}^*$ such that $\kappa^{-1}\tau_{\mathcal{M}, 1}\kappa$ can be computed by a cyclic automaton. Indeed, one can take $\kappa = \tau_{\mathcal{C}, x}$, where \mathcal{C} is the automaton in Figure 8. Then we can compute the power series $\chi(\kappa^{-1}\tau_{\mathcal{M}, 1}\kappa)$, and see directly that its coefficients are nonzero. At that point, Lemma 9.2 implies that $\kappa^{-1}\tau_{\mathcal{M}, 1}\kappa$ acts transitively on Q^* , and hence so does $\tau_{\mathcal{M}, 1}$.

⁴Note that [B+] does not distinguish between an automaton and its inverse. We do, so some of our automata are actually inverses of the automata described there.

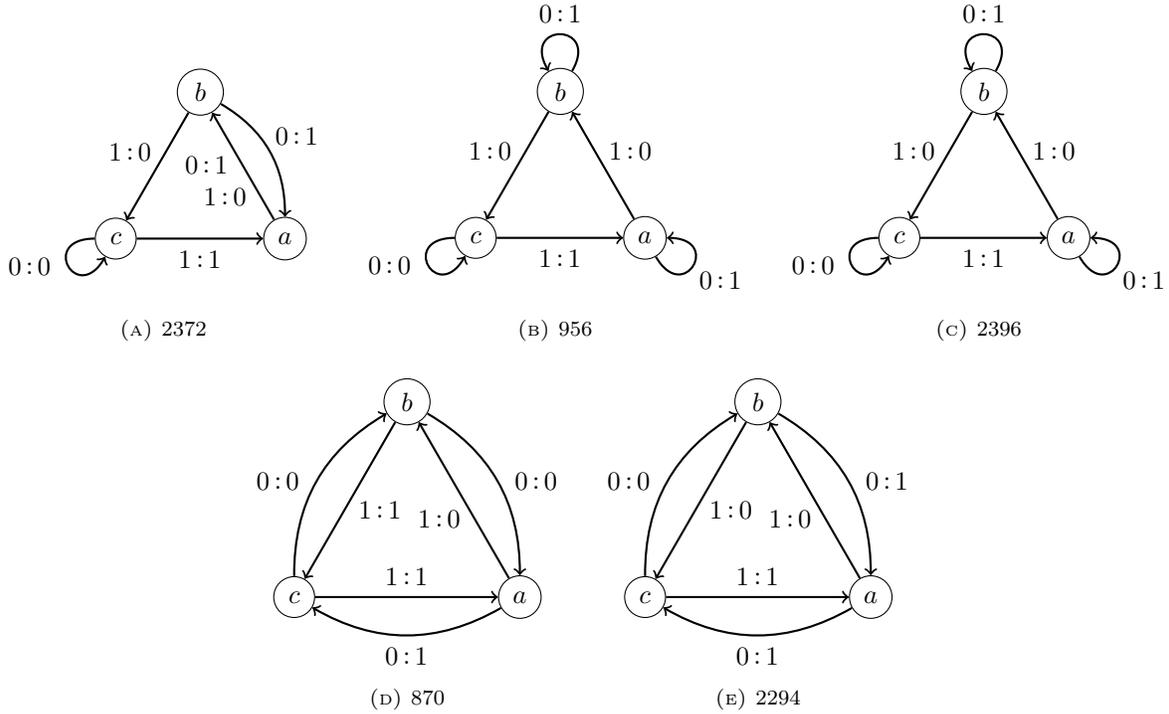


FIGURE 7. The five element-cotransitive 3-state automata on a binary alphabet.

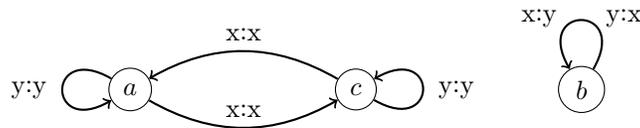


FIGURE 8. An automaton to conjugate \mathcal{M}_{2372} into a cocyclic automaton.

Proposition 10.2. *The element-cotransitive invertible $(3, 2)$ -automata are precisely the five automata pictured in Figure 7, up to relabeling of A and Q .*

Example 10.3. Of course, there are automata which are not element-cotransitive, but still satisfy the conditions of Corollary 8.4. As we saw, one example is the Bellaterra automaton. A natural and easy to analyze example is the automaton $\mathcal{M} = (Q, A, \tau, \sigma)$ that implements division by 3 modulo 2^n . (This is automaton number 924 in [B+]. See [BŠ] for more on this construction and related ones.) We will also see that its Schreier graphs do not form a family of expanders.

A quick way to define this automaton is that for $a, b \in Q = \{0, 1, 2\}$ and $x, y \in A = \{0, 1\}$, we have $a \xrightarrow{x:y} b$ if and only if

$$a + 3y = x + 2b.$$

This automaton is pictured in Figure 9. Note that for convenience we abuse notation slightly and call two of the states, 0 and 1, by the same name as the letters in the alphabet.

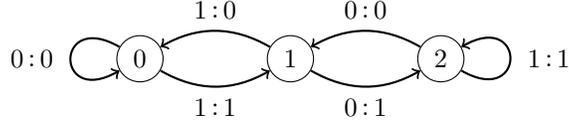


FIGURE 9

By assumption, if $a \xrightarrow{x:y} b$, then for any $x' \in \mathbb{Z}/2^{n-1}\mathbb{Z}$, we have the following equalities in $\mathbb{Z}/2^n\mathbb{Z}$:

$$\begin{aligned} x + 2b &= a + 3y \\ x + 2x' - a &= 3y + 2x' - 2b \\ \frac{(x + 2x') - a}{3} &= y + 2\frac{x' - b}{3}. \end{aligned}$$

That is, if x is the least significant binary digit of a number $X \in \mathbb{Z}/2^n\mathbb{Z}$, and $x' \in \mathbb{Z}/2^{n-1}\mathbb{Z}$ is the number corresponding to the rest of its digits, then the least significant digit of $(X - a)/3$ is y , and the rest of the digits are given by $(x' - b)/3$. It follows that if we identify a number $x \in \mathbb{Z}/2^n\mathbb{Z}$ with its binary representation in $\{0, 1\}^n$ (with the least significant digit on the left), then we have, for each $a \in \{0, 1, 2\}$,

$$\sigma_a(x) = \frac{x - a}{3}.$$

By a symmetric argument, the dual of this automaton implements division by 2 modulo 3. Phrasing this in terms of the original automaton \mathcal{M} , we interpret a length- m word in $\{0, 1, 2\}$ as the representation of a number modulo 3^m written in ternary with the least significant digit on the right. Then for each $x \in \{0, 1\}$,

$$\tau_x(a) = \frac{a - x}{2}.$$

In particular, τ_0 divides a number by 2. Since 2 generates the multiplicative group $(\mathbb{Z}/3^m\mathbb{Z})^*$, that group is an orbit of τ_0 . So for every m , there is an orbit of τ_0 in Q^m of size $2 \cdot 3^{m-1}$. By Theorem 8.3, it follows that $\text{diam}(\Gamma_{\mathcal{M},n}) = O(n^2)$. In fact, it can be checked explicitly that $\text{diam}(\Gamma_{\mathcal{M},n}) = O(n)$. This can be seen from the observation that the sequence of applications of σ_1 , σ_2 , and σ_3 necessary to send the binary number x to $00\dots 0$ is essentially the representation of x in base 3.

However, the group $G_{\mathcal{M}} = \langle \sigma_0, \sigma_1, \sigma_2 \rangle$ is generated by $\mu = \sigma_0^{-1}$ and $\alpha = \sigma_1^{-1}\sigma_0$, which are multiplication by 3 and addition of 1, respectively. E.g., $\sigma_2 = \mu^{-1}\alpha^{-2}$. It follows that the group action factors through the group of upper-triangular 2 by 2 matrices via

$$\mu \mapsto \begin{pmatrix} 3 & 0 \\ 0 & 1 \end{pmatrix} \quad \alpha \mapsto \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$$

This group is solvable, and therefore amenable. It follows that its Schreier graphs with respect to a fixed set of generators cannot be expanders [Lub, 3.3.7]. So, the family $\{\Gamma_{\mathcal{M},n}\}_{n=1}^{\infty}$ is not a family of expanders.

So, there are automata to which our general results apply, but whose Schreier graphs do not form a family of expanders. More work is necessary to find sufficient conditions for when an automaton gives rise to a family of expanders.

Example 10.4. It is possible to check by a computation that there are no element-cotransitive invertible $(4, 2)$ -automata. In fact, for any such automaton, any τ_x does not act transitively on Q^4 .

Example 10.5. There are seven $(5, 2)$ -automata which are not cocyclic, whose duals act element-transitively on Q^{10} . Of these, just one is bireversible, as the Aleshin and Bellaterra automata are. It is pictured in Figure 10. However, it is unlikely that there is an automatic automorphism $\kappa : Q^* \rightarrow$

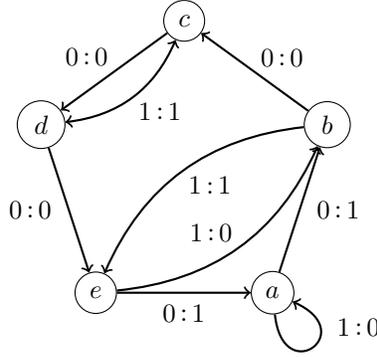


FIGURE 10. The only candidate to be a element-cotransitive bireversible $(5, 2)$ -automaton.

Q^* such that $\kappa^{-1}\tau_0\kappa$ is implemented by a cocyclic automaton. We have checked that if there is such a κ , the automaton implementing it would need to have at least 48668 states. So we cannot repeat the analysis in Example 10.1 of the automaton pictured in Figure 7a.

11. REMARKS AND FURTHER WORK

11.1. This paper is concerned with automata which mix the words in A^* , so we are interested in automata whose action on A^* changes many letters of a typical word. At the other end of the spectrum, there are *bounded* automata, which typically only change a few letters at the beginning of a word. Such automata were investigated in [B], where it was shown that their Schreier graphs typically have exponential growth.

11.2. The results in Section 8 can be extended to non-invertible Mealy automata as well. Since we are primarily interested only in regular graphs, we prove only the simpler case. For example, to state Theorem 8.1 more generally, one needs to consider the size of balls in $\Gamma_{\mathcal{M},\infty}$ defined in terms of directed paths, but the result about the diameter of $\Gamma_{\mathcal{M},n}$ still needs diameter to be defined in terms of undirected paths.

11.3. As noted in the proof of Corollary 7.1, any product of two generators of the Bellaterra group $G_{\mathcal{B}} = \langle \sigma_{\mathcal{B},a}, \sigma_{\mathcal{B},b}, \sigma_{\mathcal{B},c} \rangle$ belongs to the Aleshin group $G_{\mathcal{A}} = \langle \sigma_{\mathcal{A},a}, \sigma_{\mathcal{A},b}, \sigma_{\mathcal{A},c} \rangle$. We used this fact to show that, since the Bellaterra graphs have small diameter, so do the Aleshin graphs. In fact, it can also be used to show that if the Bellaterra graphs form a (two-sided) expander family, so do the Aleshin graphs. In other words, Conjecture 1.3 implies Conjecture 1.2.

11.4. A tree automorphism $\alpha : A^* \rightarrow A^*$ is spherically transitive if and only if it is conjugate in $\text{Aut}(A^*)$ to the adding machine ρ , i.e., the automorphism which interprets a word in A^n as the base- $|A|$ representation of a number modulo $|A|^n$, and adds one to that number. The adding machine is an automatic automorphism, e.g., the binary adding machine is pictured in Figure 4.

One might hope that whenever an automatic automorphism $\alpha \in \text{FAut}(A^*)$ is conjugate to ρ in $\text{Aut}(A^*)$, it is also conjugate to ρ in $\text{FAut}(A^*)$. If so, we would have an algorithm for determining whether a given automatic automorphism is transitive. In fact, since we can enumerate the element-transitive cyclic automata, it would be enough if every transitive $\alpha \in \text{FAut}(A^*)$ were conjugate in $\text{FAut}(A^*)$ to some cyclic automorphism.

However, Example 10.5 suggests that, in the dual of the automaton in Figure 10, σ_0 is transitive but not conjugate in $\text{FAut}(A^*)$ to any cyclic automaton, in particular to ρ . However, we do not have a proof of either of these facts.

Problem 11.1. *Exhibit a transitive $\alpha \in \text{FAut}(A^*)$ which is not conjugate (in $\text{FAut}(A^*)$) to a cyclic $\beta \in \text{FAut}(A^*)$. (Or prove that there is no such α .)*

Problem 11.2. *Characterize the automorphisms in $\text{FAut}(A^*)$ which are conjugate in $\text{FAut}(A^*)$ to a cyclic automorphism.*

We can, however, exhibit a cyclic $\alpha \in \text{FAut}(A^*)$ which is not conjugate in $\text{FAut}(A^*)$ to the adding machine ρ :

Proposition 11.3. *Let $\mathcal{M} = (Q, A, \tau, \sigma)$ be the dual of the automaton in Figure 7b, where $Q = \{0, 1\}$ and $A = \{a, b, c\}$. Then σ_1 acts transitively on A^* , but there is no $\kappa \in \text{FAut}(A^*)$ such that $\kappa^{-1}\sigma_1\kappa = \rho$*

Sketch of proof: Given an eventually periodic word $w \in A^*$, we let $h(w)$ denote the smallest number n such that w is periodic after the first n letters.

Note that if ρ is the adding machine, then for any eventually periodic word $v \in A^\infty$, we have

$$h(\rho^n(v)) = O(\log n).$$

Moreover, after a finite number of steps, the periodic part of $\rho^n(v)$ stabilizes. It follows that for any $\kappa \in \text{FAut}(A^*)$, we have

$$h(\kappa\rho^n(v)) = O(\log n)$$

and since this applies to any v ,

$$h(\kappa\rho^n\kappa^{-1}(v)) = O(\log n)$$

On the other hand, taking $\alpha = \sigma_1$, we can check that if we read $w \in A^n$ as a ternary number modulo 3^n (with $c = 0, a = 1, b = 2$), we have

$$\alpha(w) = \frac{w+1}{-2}.$$

It follows that for $w = ccc\dots$,

$$h(\alpha^{-n}(w)) \sim (\log_3 2)n$$

Thus α^{-1} and ρ are not conjugate in $\text{FAut}(A^*)$. It is easy to check that ρ and ρ^{-1} are conjugate, so α and ρ are not conjugate in $\text{FAut}(A^*)$. \square

11.5. More generally, an open problem is the classification of conjugacy classes in $\text{FAut}(A^*)$. The conjugacy classes of $\text{Aut}(A^*)$ can be described in terms of orbit trees [GNS]. This tree captures the information about the orbits of an automorphism $\alpha \in \text{Aut}(A^*)$, e.g. a ray with few branches in the orbit tree corresponds to a sequence of quickly growing orbits. Information about this tree can tell us whether we can apply, e.g., Theorem 8.3.

Of course, not all orbit trees arise from elements of $\text{FAut}(A^*)$, since there are uncountably many. Moreover, not all automatic automorphisms with the same orbit tree are conjugate in $\text{FAut}(A^*)$, as seen in Proposition 11.3.

In [BBSZ], the problem is solved for *bounded automorphisms*, and more generally automorphisms with *finite orbit-signalizer*. Such “small” automorphisms are unlikely to give expanders, so we are interested in the other end of the spectrum, automorphisms with many nontrivial sections on every level.

11.6. *Automaton groups*, i.e., groups of the form $G_{\mathcal{M}}$ for some Mealy automaton \mathcal{M} , are of independent interest in group theory.⁵ A famous example is the *Grigorchuk group*, which is the first known group whose growth function is intermediate between polynomial and exponential (see [GP, G1]). For more on automaton groups, see [BGŠ, GNS, Nek].

⁵Note that the term *automatic group* has a different meaning in the literature, one we do not use in this paper.

11.7. The structure of $G_{\mathcal{M}}$ as an abstract group can give us information on whether or not the graphs $\Gamma_{\mathcal{M},n}$ form a family of expanders. For example, if $G_{\mathcal{M}}$ is amenable then $\{\Gamma_{\mathcal{M},n}\}_{n=1}^{\infty}$ is not a family of expanders [Lub, 3.3.7]. We already used this fact in Example 10.3 to show that the Schreier graphs of the automaton defined there are not expanders.

11.8. On the other hand, sometimes the structure of $G_{\mathcal{M}}$ is enough to guarantee that $\{\Gamma_{\mathcal{M},n}\}_{n=1}^{\infty}$ is a family of expanders. Notably, if $\Gamma_1, \Gamma_2, \dots$ are Schreier graphs (with respect to a fixed generating set) of a group with *Kazhdan property (T)*, and $|\Gamma_i| \rightarrow \infty$, then these graphs must form a family of expanders [Lub, 3.3.4]. This fact was used by Margulis to give the first explicit construction of expanders [Mar].

In [GM], it was shown that there are Mealy automata \mathcal{M} for which $G_{\mathcal{M}}$ has property (T), so Mealy automata can be used to construct expander families. The authors do not state their construction in our language, it is implicit in their results. Of course, the corresponding expander Schreier graphs would be strongly explicit. However, their construction is significantly more complicated and perhaps less suitable to generalizations. Note also that the groups $G_{\mathcal{A}}$ and $G_{\mathcal{B}}$ do not have property (T), so this approach is not sufficient to prove our conjectures 1.2 and 1.3.

11.9. As mentioned in section 4, the groups $G_{\mathcal{A}}$ and $G_{\mathcal{B}}$ corresponding to the Aleshin automaton and Bellaterra automaton are, respectively, the free group of 3 generators and the free product of 3 copies of the cyclic group C_2 . There are generalizations of the Aleshin and Bellaterra automata, described in [SV] and [VV2], whose automata groups are free groups of higher order and free products of more copies of C_2 . The Schreier graphs of these automata are likely to have similar properties to the graphs A_n and B_n .

Problem 11.4. *Are the Schreier graphs of the generalized Aleshin and Bellaterra automata described in [SV] and [VV2] expanders? Do their diameters grow polynomially?*

11.10. In [MSS], the ideas of [BL] were extended to construct families of bipartite Ramanujan graphs (i.e., expander graphs with optimal spectral gap) of arbitrary degree. The construction uses a new technique to pick a particular 2-lift of a graph which does not introduce any new large eigenvalues. We should note that this construction is not *very explicit*, in the sense given above.

11.11. In [G2, Section 10], Grigorchuk shows that in a certain formal sense, the Aleshin and Bellaterra automata are examples of *asymptotic expanders*, thus giving further evidence to Conjectures 1.2 and 1.3. He also states these conjectures as open problems, and suggests that a sequence of Schreier graphs constructed by a finite automaton cannot be Ramanujan.

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