

L24

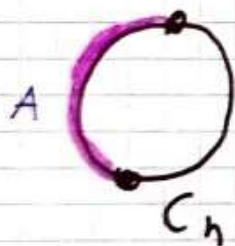
Expander graphs

Igor Pak

Def $G = (V, E) \leftarrow$ conn. graph, d-regular, $\varepsilon > 0$
 G is ε -expander if $\forall A \subset V, A \neq \emptyset$
 $\#E(A, V-A) \geq \varepsilon |A|, |A| \leq \frac{|V|}{2}$

$n = |V|, \{G_n\} \leftarrow$ expander if $d = \text{const}, \exists \varepsilon > 0$
s.t. G_n is ε -expander $\forall n$.

Ex 1) n -cycle C_n is $(\frac{4}{n})$ -expander $\leftarrow \text{Exc}$



2) d -hypercube is 1-expander $\leftarrow \text{Exc}$

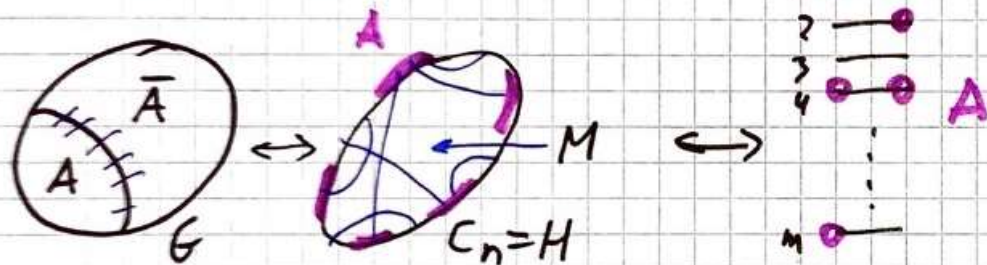


3) random 3-regular graph
is $> \frac{1}{100}$ -expander w.h.p. $\leftarrow \text{Exc}^*$

$\leftarrow$ uniform
sum of 3 random matchings

Th $G = (V, E) \leftarrow$ random 3-regular graph, $|V| = n = 2m$
 formed by $G = H \cup M$, $H = n$ -cycle, $M \leftarrow$ uniform PM
then G is ϵ -expander w.h.p. $\epsilon =$

D (sketch)



$$|A| = k \quad 1 \leq k \leq m$$

$$E[\# \text{ edges from } A \text{ to } \bar{A}] = 2 \frac{k(n-k)}{n(n-1)} \cdot m$$

$$\geq 2 \frac{k}{n} \frac{n}{2} = k$$

$$P[E(A, \bar{A}) \geq \frac{k}{2}] > 1 - n^{-ck}$$

$$\Rightarrow P[E(A, \bar{A}) \geq \frac{|A|}{2}] > 1 - \sum_{k=1}^m \binom{n}{k} n^{-ck} > 1 - \delta > 0$$

$\forall A \subset V$
 $1 \leq |A| \leq n/2$

Motivation: [AKS] $\leftarrow$ Ajtai-Komlós-Szemerédi (1983)

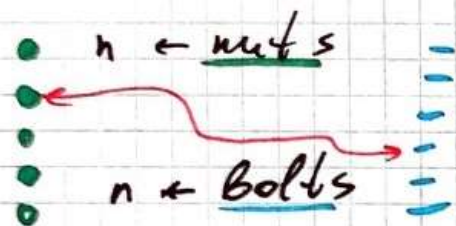
Sorting
the network

$O(n \log n)$ sorting
w/ depth $O(\log n)$

①, ②, ... ①
 deterministic algorithm.

②

Nuts & Bolts [Alon-Blum-Fiat-Kannan-Naor-Ostrovsky] 1994



Random Alg

- (1) pick random bolt, compare w/ all nuts
 - (2) compare matching nut w/ all bolts
- split the problem into 2 roughly $\frac{n}{2}$ size

cost $O(n \log n) \leftarrow f(n) = 2n + 2f(\frac{n}{2}) =$
 $= 2n + 2n + 4f(\frac{n}{2}) = \dots$

Th [Alon ++]

$\exists O(n (\log n)^{3+\epsilon})$ deterministic Alg

[AKS], [Alon ++] required

Th [Margulis, 1973]

$\exists$ explicit construction of ϵ -expanders, d, ϵ -fixed

$n \rightarrow \infty$

Def $G = (V, E)$ has an explicit construction if $V \subseteq \{0,1\}^L$, $L = O(\log n)$
 $\forall i \in V, (v_i, v_j) \in E$ in $\text{Poly}(\log n)$ time

Nuts & Bolts det Alg



- (1) compare all bolts w/ corr. nuts
 - (2) delete half nuts whose rank far from $\frac{d}{2}$
- Repeat until middle nut is found.
Split the problem.

Exc get $O(n \log^4 n)$ alg.

Margulis - Gabber-Galil expanders $G_p = (V_p, E_p)$

$$V_p = \mathbb{Z}_p^2$$

$$d=8$$

$$E_p = \{ (a, b) \leftrightarrow (a \pm 1, b), (a, b \pm 1), (a, b \pm a), (a \pm b, b) \}$$

Th [M, G-G'81]

$\{ G_p = (V_p, E_p) \} \leftarrow \frac{1}{160}$ -expanders.

Margulis '73

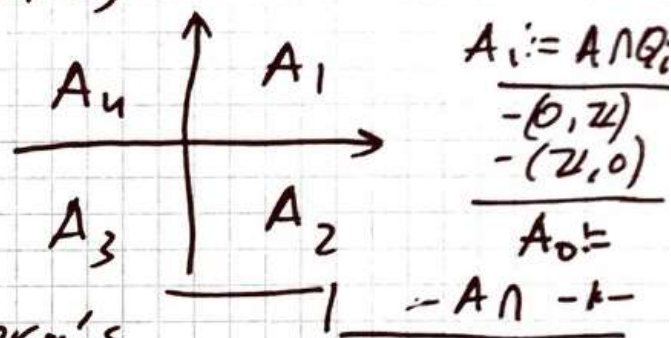
Gabber-Galil

$\leftarrow$ Algebraic proof using Kazhdan (Γ)

$\leftarrow$ first "elementary proof"

$$D \quad S(a, b) := (a, a+b) \quad T(a, b) := (a+b, b)$$

$$\underline{\quad} \quad \forall A \subset \mathbb{Z}^2, O=(0,0) \\ |E(A, \bar{A})| \geq \frac{1}{7} |A|$$



$$\boxed{DL} \quad |S(A_1)| = |T(A_1)| = |A| \quad \text{since } S, T \in \text{perm's}$$

$$S(A_1) \cap T(A_1) = \emptyset \quad \leftarrow \quad (a, a+b) = (a'+b', b') \Rightarrow b = -b' \times$$

$$S(A_1), T(A_1) \subset Q_1 \Rightarrow \text{at least } \frac{1}{2} \text{ of } E(A_1, *) \ni A$$

$$\Rightarrow \underline{\text{claim 1}} \quad \#E[A - A_0, \bar{A}] \geq |A - A_0| = |A| - |A_0|$$

$$\underline{\text{claim 2}} \quad \#E[A_0, \bar{A}] \geq 4|A_0| - 3|A - A_0| = 7|A_0| - 3|A|$$

EXC

$$\underline{\text{claim 1}} + \underline{\text{claim 2}} \Rightarrow \underline{\quad} \quad (\text{By balancing ineq's}) \quad \square$$

Now $\mathbb{Z}^2 \rightarrow \mathbb{Z}_p$ is difficult but routine \(\square\)

Th (Margulis '74) $\forall k$

$\text{Cayley}(G_p, R_p) \leftarrow \text{expander}$

$$G_p = \text{SL}(k, \mathbb{F}_p), \quad R_p = \left\{ \begin{pmatrix} 1 & \pm 1 \\ & 1 \end{pmatrix}_{k \times k} \right\}$$

$G = G \leftarrow \text{Margulis } k=3$

$$\left(\text{SL}(3, p) \right) \begin{pmatrix} a \\ b \\ c \end{pmatrix} \leftarrow \text{exp}, \oplus \text{ proj onto } (a, b), \underline{c=1}$$

Cor $\Gamma_k(\mathbb{Z}_p^k)$ are expanders $\forall k$

$$\Rightarrow \text{Mix time of PRA on } \Gamma_k(\mathbb{Z}_p^k) \\ = c_k (\log p)$$

vs $O(k^6 (\log p)^3)$ last time.

Th [P-Zuk $\oplus$ Kossobov]

$$\text{Mix} = O(k^3 \log p)$$

Next week:

Cor/Th [Kaluza-Kielak - Nowak-Ozawa '17 + '18]

$\forall k$ PRA mixes on $\Gamma_k(G)$ $\forall k \geq 5$
in $O_k(\log |G|)$ time.

Real Th $\Gamma_k(G)$ are ε_k -expanders $\forall k \geq 5$

/ conjecture Lubotzky-P '01 /

Note: 1) every conn. comp. of $\Gamma_k(G)$

2) proof $\leftarrow$ computer-assisted