

The Plactic Monoid (M.P. Schützenberger)

A totally ordered alphabet $\{a_1 < a_2 < \dots < a_n\}$, $\{1 < 2 < \dots < n\}$

Motivation Given a $w \in A^*$, what is the length of the longest nondecreasing subword of w ? (complexity of finding such a subword is $O(n \ln n)$)

132541, longest nondecreasing subword has length 3.

C. Schensted (now $E_n E_n$) gave an algorithm that finds the length without finding the sequence using Young tableau

Definition (tableau)

* a nondecreasing word $u \in A^*$ is called a row

* rows $u = x_1 \dots x_r$, $v = y_1 \dots y_s$, $u \triangleleft v$ if $r \geq s$, $x_i < y_i$

$$\begin{matrix} & \leq & \\ \wedge & x_1, x_2, \dots, x_r & \\ & y_1, y_2, \dots, y_s & \end{matrix}$$

* every w has a unique factorisation $w = u_1 u_2 \dots u_k$ into rows u_i of maximal length

a tableau is a word w such that $u_1 \triangleright u_2 \triangleright \dots \triangleright u_k$

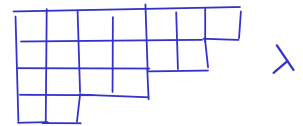
ex 68 4556 223357 112444 \mapsto

1	1	2	4	4	4
2	2	3	3	5	7
4	5	5	6		
6	8				

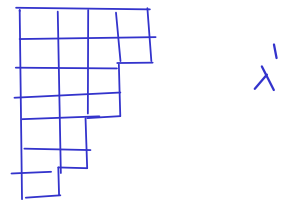
 tableau has shape $(7, 6, 4, 2)$ (length of rows)

 ↑
 columns are increasing subwords

* A Young tableau is a tableau with partition shape λ

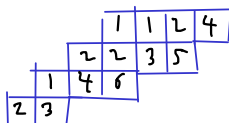


* λ' denotes the conjugate partition of λ ($\lambda' = 4433221$)



* A skew Young tableau is a tableau of skew shape

λ/μ ($\mu \leq \lambda$) ex $\lambda = 7642$, $\mu = 321$



Schensted's insertion algorithm

input $w \in A^*$ output a tableau $t = P(w)$

(length of first row
= length of longest nondecreasing
subword of w)

elementary step:

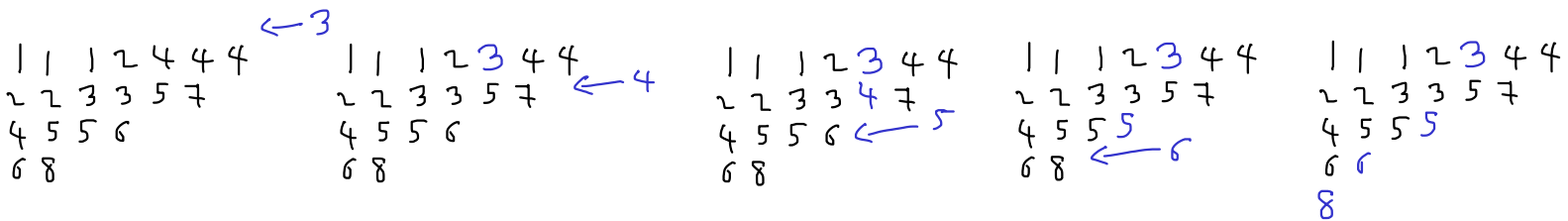
- $v = y_1 y_2 \dots y_r$ row, insert x

$y_1 y_2 \dots y_r \leftarrow x$

$$P(v) = \begin{cases} y_1 y_2 \dots y_r x & \text{if } vx \text{ is a row } (y_r \leq x) \\ y_1 y_2 \dots x y_{i+1} \dots y_r & \text{if } y_i \text{ leftmost letter of } v, y_i > x \\ & \text{"x bumps } y_i \text{"} \end{cases}$$

- insert x in a tableau: insert x in first row, if y_i is bumped, insert y_i second row ----

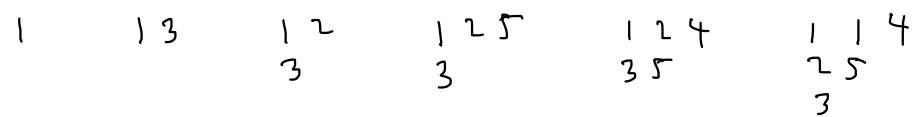
ex insert $x=3$ in $\begin{matrix} 1 & 1 & 1 & 2 & 4 & 4 & 4 \\ 2 & 2 & 3 & 3 & 5 & 7 \\ 4 & 5 & 5 & 6 \\ 6 & 8 \end{matrix} \leftarrow 3$



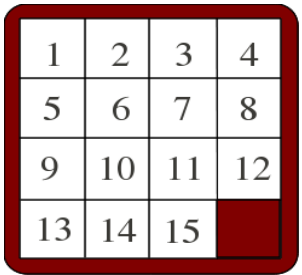
General case: given $w = x_1 \dots x_n$

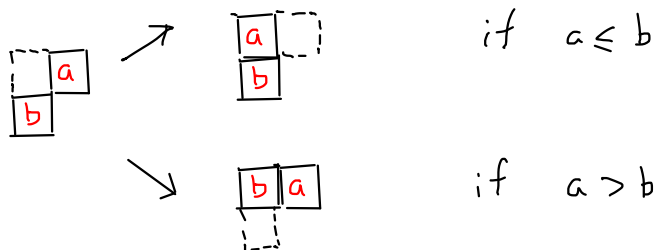
insert x_1 to \emptyset get $P(x_1)$
insert x_2 to $P(x_1) \dots P(x_1 x_2)$
:
insert x_n to $P(x_1 \dots x_{n-1})$ get $P(x_1 x_2 \dots x_n)$

example $P(132541)$



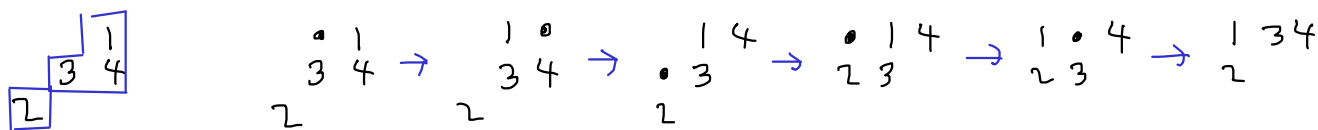
Schensted's insertion algorithm (gen-de-taquin version)





(do a move to preserve a tableau)

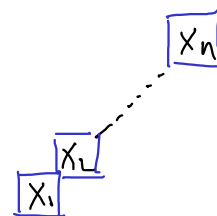
"jeu de taquin" is a set of rules to transform a skew tableau into a Young tableau



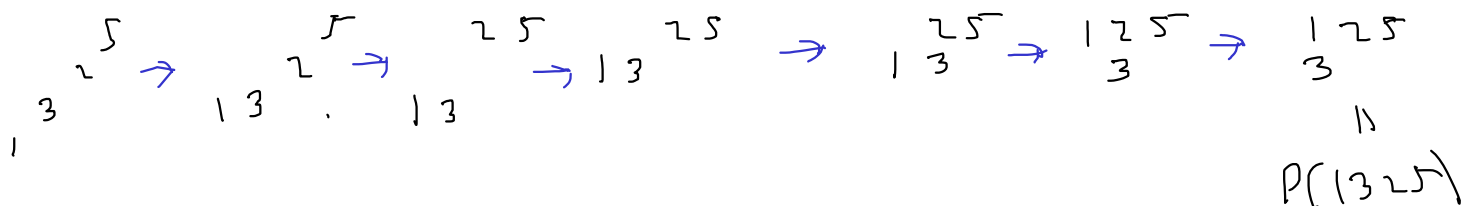
ex jeu de taquin models a "corporate restructuring"



Claim Schensted insertion $w = x_1 x_2 \dots x_n$ \equiv jeu de taquin



ex 1 3 2 5



Thm 1 - The maximal length of a nondecreasing subword of w is the length of the first row of $P(w)$

- The maximal length of a decreasing subword of w is the length of the first column of $P(w)$

We will prove something stronger

let $l_k(w)$ be the maximum of the length of k disjoint nondecreasing subwords of w .

$l_1(w)$ = maximal length of a nondecreasing subword of w

($l'_k(w)$ same for decreasing subwords)

ex 13254 $l_1(w) = 3$, $l_2(w) = 5$, $l_3(w) = 5$, $l_4(w) = 5$

Thm 2 (Greene) For $w \in A^*$, If $P(w)$ has shape $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_r)$
 & $\lambda' = (\lambda'_1, \lambda'_2, \dots, \lambda'_s)$ then for $k=1, \dots, r$ $\lambda_k = l_k(w) - l_{k-1}(w)$
 $k=1, \dots, s$ $\lambda'_k = l'_k(w) - l'_{k-1}(w)$

ex $w = 13254$, $P(w) = \begin{array}{ccc} 1 & 2 & 4 \\ & 3 & 5 \end{array}$, $\lambda = (3, 2)$, $\lambda'_1 = 2$, $\lambda'_2 = 4$, $\lambda'_3 = 5$.

, $l_1 = 3$, $l_2 = 3+2$, $l_3 = 5$

, $l'_1 = 2$, $l'_2 = 4$, $l'_3 = 5$.

Thm 2 implies Thm 1.