

# The conjugate gradient method

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The conjugate gradient method is an iterative method for solving the linear system  $Ax = b$ , where  $A \in \mathbb{R}^{n \times n}$  is symmetric positive-definite.

Let  $\langle x, y \rangle_A := \langle Ax, y \rangle$  be the inner product defined by  $A$  and  $\|x\|_A := \sqrt{\langle x, x \rangle_A}$  be the induced norm. Given an initial guess  $x^{(0)}$  for the solution  $x^*$ , the  $k^{\text{th}}$  iterate of the method is defined as

$$x^{(k)} = \arg \min_{x \in x^{(0)} + \mathcal{K}_k(A, r^{(0)})} \|x^* - x\|_A,$$

where  $\mathcal{K}_k(A, r^{(0)})$  is the Krylov subspace  $\text{span}\{A^j r^{(0)}\}_{j=0}^{k-1}$  and  $r^{(0)} = b - Ax^{(0)}$ . (In other words, the  $A$ -norm of the error is minimized over the  $k^{\text{th}}$  affine Krylov subspace generated by the initial residual and translated by the initial guess.)<sup>1</sup>

Let us abbreviate  $\mathcal{K}_k(A, r^{(0)})$  as  $\mathcal{K}_k$  and write  $r^{(k)} = b - Ax^{(k)}$  for the residual of the  $k^{\text{th}}$  iterate. The iterate  $x^{(k)}$  is therefore the  $A$ -orthogonal projection of  $x^*$  onto  $x^{(0)} + \mathcal{K}_k$ , defined by the Galerkin conditions  $x^{(k)} - x^{(0)} \in \mathcal{K}_k$  and  $x^* - x^{(k)} \perp_A \mathcal{K}_k$ ; we note that the orthogonality condition is equivalent to  $r^{(k)} \perp \mathcal{K}_k$ .

Now suppose that  $\{p^{(j)}\}_{j < k}$  is a basis of  $\mathcal{K}_k$  and let  $P_k = [p^{(0)} \ \dots \ p^{(k-1)}]$ . Then  $x^{(k)} = x^{(0)} + P_k y^{(k)}$ , where

$$y^{(k)} = \arg \min_{y \in \mathbb{R}^k} \|x^* - (x^{(0)} + P_k y)\|_A.$$

If  $p^{(k)}$  is such that  $\{p^{(j)}\}_{j < k+1}$  is a basis of  $\mathcal{K}_{k+1}$ , we can express the next iterate  $x^{(k+1)}$  in an analogous manner – that is,  $x^{(k+1)} = x^{(0)} + P_{k+1} y^{(k+1)}$ , where  $P_{k+1} = [P_k \ p^{(k)}]$ . Writing  $y^{(k+1)} = \begin{bmatrix} \tilde{y}^{(k)} \\ \alpha_k \end{bmatrix}$  for some  $\tilde{y}^{(k)} \in \mathbb{R}^k$  and  $\alpha_k \in \mathbb{R}$ , we see that

$$\begin{aligned} x^* - (x^{(0)} + P_{k+1} y^{(k+1)}) &= [x^{(k)} - (x^{(0)} + P_k \tilde{y}^{(k)})] + [(x^* - x^{(k)}) - \alpha_k p^{(k)}] \\ &= P_k (y^{(k)} - \tilde{y}^{(k)}) + [(x^* - x^{(k)}) - \alpha_k p^{(k)}]. \end{aligned}$$

Thus, if we select  $p^{(k)}$  to be  $A$ -orthogonal to  $p^{(j)}$  for all  $j < k$ , then by the Pythagorean theorem,

$$\|x^* - (x^{(0)} + P_{k+1} y^{(k+1)})\|_A^2 = \|P_k (y^{(k)} - \tilde{y}^{(k)})\|_A^2 + \|(x^* - x^{(k)}) - \alpha_k p^{(k)}\|_A^2,$$

so the solution to the least squares problem for  $y^{(k+1)}$  is given recursively by  $\tilde{y}^{(k)} = y^{(k)}$  and  $\alpha_k p^{(k)} = \text{proj}_{p^{(k)}}^A (x^* - x^{(k)})$ . It follows that

$$\begin{aligned} x^{(k+1)} &= x^{(0)} + P_k y^{(k)} + \alpha_k p^{(k)} \\ &= x^{(k)} + \alpha_k p^{(k)}, \end{aligned} \tag{X}$$

where

$$\begin{aligned}\alpha_k &= \frac{\langle x^* - x^{(k)}, p^{(k)} \rangle_A}{\langle p^{(k)}, p^{(k)} \rangle_A} \\ &= \frac{\langle r^{(k)}, p^{(k)} \rangle}{\langle p^{(k)}, p^{(k)} \rangle_A}.\end{aligned}\tag{A}$$

This also implies that

$$r^{(k+1)} = r^{(k)} - \alpha_k A p^{(k)}.\tag{R}$$

To generate  $A$ -orthogonal vectors  $p^{(j)}$  such that  $\{p^{(j)}\}_{j < k}$  is a basis of  $\mathcal{K}_k$  for each  $k$ , we notice that  $r^{(k+1)} \perp_A \mathcal{K}_k = \text{span}\{p^{(j)}\}_{j < k}$  because  $r^{(k+1)} \perp \mathcal{K}_{k+1}$  and  $A\mathcal{K}_k \subseteq \mathcal{K}_{k+1}$ . As a result, when  $r^{(k+1)}$  is  $A$ -orthogonalized against  $p^{(k)}$ , the resulting vector will automatically be  $A$ -orthogonal to  $p^{(j)}$  for all  $j < k + 1$ , suggesting that we define

$$\begin{aligned}p^{(k+1)} &= r^{(k+1)} - \text{proj}_{p^{(k)}}^A r^{(k+1)} \\ &= r^{(k+1)} + \beta_k p^{(k)},\end{aligned}\tag{P}$$

where  $p^{(0)} = r^{(0)}$  and

$$\beta_k = -\frac{\langle r^{(k+1)}, p^{(k)} \rangle_A}{\langle p^{(k)}, p^{(k)} \rangle_A}.\tag{B}$$

Referring back to the residual equation (R), we can show by induction that the  $p^{(j)}$  thus defined will also constitute bases of successive Krylov subspaces. More precisely, suppose that the solution has not been found by the beginning of the  $k^{\text{th}}$  iteration, in the sense that  $r^{(j)} \neq 0$  for all  $j < k$ . We claim then that  $r^{(k-1)} \in \mathcal{K}_k$  and that  $\{p^{(j)}\}_{j < k}$  is an  $A$ -orthogonal basis of  $\mathcal{K}_k$ .

Indeed, if  $r^{(0)} \neq 0$ , then  $r^{(0)} \in \mathcal{K}_1 = \text{span}\{r^{(0)}\}$  and  $\{p^{(0)}\} = \{r^{(0)}\}$  is an  $A$ -orthogonal basis of  $\mathcal{K}_1$ . Now suppose that the claim holds up to the  $k^{\text{th}}$  iteration and that its hypothesis is satisfied at the beginning of the  $(k + 1)^{\text{th}}$  iteration. Then

$$r^{(k)} = r^{(k-1)} - \alpha_{k-1} A p^{(k-1)} \in \mathcal{K}_k + A\mathcal{K}_k \subseteq \mathcal{K}_{k+1},$$

so

$$p^{(k)} = r^{(k)} + \beta_{k-1} p^{(k-1)} \in \mathcal{K}_{k+1} + \mathcal{K}_k \subseteq \mathcal{K}_{k+1}.$$

In addition,  $p^{(k)} \neq 0$  because  $r^{(k)} \perp \mathcal{K}_k$  and  $r^{(k)} \neq 0$ . Hence, by construction,  $\{p^{(j)}\}_{j < k+1}$  is an  $A$ -orthogonal set of nonzero vectors in  $\mathcal{K}_{k+1}$  and is moreover a basis thereof, since  $\dim(\mathcal{K}_{k+1}) \leq k + 1$ .

An immediate consequence is that  $\{r^{(j)}\}_{j < k}$  will be an orthogonal basis of  $\mathcal{K}_k$  for all such iterations: if (say)  $i < j < k$ , then  $r^{(i)} \in \mathcal{K}_{i+1} \subseteq \mathcal{K}_j$ , and we know that  $r^{(j)} \perp \mathcal{K}_j$ . Furthermore, the iteration will break down exactly when  $r^{(k)} \in \mathcal{K}_k$ , or equivalently,  $r^{(k)} = 0$ , meaning that the solution was attained in the  $k^{\text{th}}$  iteration.

We can also derive alternative formulas for the scalars  $\alpha_k$  and  $\beta_k$  that reduce the number of inner products in each iteration. First, using the fact that  $r^{(k)} \perp \mathcal{K}_k = \text{span}\{p^{(j)}\}_{j < k}$ , we obtain

$$\begin{aligned}
\alpha_k &= \frac{\langle r^{(k)}, p^{(k)} \rangle}{\langle p^{(k)}, p^{(k)} \rangle_A} \\
&= \frac{\langle r^{(k)}, r^{(k)} + \beta_{k-1} p^{(k-1)} \rangle}{\langle p^{(k)}, p^{(k)} \rangle_A} \\
&= \frac{\langle r^{(k)}, r^{(k)} \rangle}{\langle p^{(k)}, p^{(k)} \rangle_A}.
\end{aligned} \tag{A}$$

Hence

$$\begin{aligned}
\beta_k &= -\frac{\langle r^{(k+1)}, p^{(k)} \rangle_A}{\langle p^{(k)}, p^{(k)} \rangle_A} \\
&= -\alpha_k \frac{\langle r^{(k+1)}, p^{(k)} \rangle_A}{\langle r^{(k)}, r^{(k)} \rangle} \\
&= \frac{\langle r^{(k+1)}, r^{(k+1)} - r^{(k)} \rangle}{\langle r^{(k)}, r^{(k)} \rangle} \\
&= \frac{\langle r^{(k+1)}, r^{(k+1)} \rangle}{\langle r^{(k)}, r^{(k)} \rangle}.
\end{aligned} \tag{B}$$

In summary,

$r^{(0)}$	$= b - Ax^{(0)}$	
$p^{(0)}$	$= r^{(0)}$	
$\alpha_k$	$= \frac{\langle r^{(k)}, r^{(k)} \rangle}{\langle p^{(k)}, p^{(k)} \rangle_A}$	<b>A</b>
$x^{(k+1)}$	$= x^{(k)} + \alpha_k p^{(k)}$	<b>X</b>
$r^{(k+1)}$	$= r^{(k)} - \alpha_k A p^{(k)}$	<b>R</b>
$\beta_k$	$= \frac{\langle r^{(k+1)}, r^{(k+1)} \rangle}{\langle r^{(k)}, r^{(k)} \rangle}$	<b>B</b>
$p^{(k+1)}$	$= r^{(k+1)} + \beta_k p^{(k)}$	<b>P</b>

1. The choice of this minimization problem can be partially motivated as follows. In view of the fact that  $x^* = x^{(0)} + A^{-1}r^{(0)}$  and that  $A^{-1}$  is a polynomial in  $A$  of degree at most  $n - 1$ , in the  $k^{\text{th}}$  iteration of the method, we seek an approximation to the solution of the form  $x^{(0)} + p_{k-1}(A)r^{(0)}$ , where  $p_{k-1}$  is a polynomial of degree at most  $k - 1$ . This guarantees that the  $A$ -norm of the error decreases monotonically and that the solution is found in at most  $n$  iterations (in exact arithmetic). Although the choice of the objective function is not canonical, it turns out that this choice leads to a particularly tractable method. 