Problem Set 3 Solutions

Foundations of Number Theory

Math 435, Fall 2006

1. (20 pts.) Let p be prime, $\alpha \in \mathbb{N}$. We have $\sigma(p^{\alpha}) = \frac{p^{\alpha}-1}{p-1} + p^{\alpha} < 2p^{\alpha}$, hence p^{α} is not perfect. Let q be a prime different from p, and suppose pq is perfect. Then

$$2pq = \sigma(pq) = (p+1)(q+1),$$

hence p|(q+1) and q|(p+1). Write q+1 = pk, p+1 = ql with $k, l \in \mathbb{N}^{>0}$. Then (kl)pq = (p+1)(q+1), hence kl = 2. Thus either k = 1, l = 2, or k = 2, l = 1; in the first case p = q + 1, 2q = p + 1, hence q = 2, p = 3, and in the second case p = 2, q = 3. In both cases pq = 6.

2. (20 pts.) Suppose $n \in \mathbb{N}^{>0}$ is perfect. Then

$$\sum_{d|n} d = 2n,$$

hence, dividing by n on both sides of this equation:

$$\sum_{d|n} \frac{1}{n/d} = 2$$

Now note that

$$\sum_{d|n} \frac{1}{n/d} = \sum_{d|n} \frac{1}{d}$$

3. (20 pts.) Suppose f * g = 0, that is,

$$\sum_{d|n} f(d)g(n/d) = 0 \quad \text{for all } n \in \mathbb{N}^{>0}.$$

Assume that $f, g \neq 0$; we then need to show that $f * g \neq 0$. Since $f \neq 0$, there is some $k \in \mathbb{N}^{>0}$ with $f(k) \neq 0$; take k minimal with this property. Similarly, let $l \in \mathbb{N}^{>0}$ be minimal with $g(l) \neq 0$, and put n := kl. We claim that $(f * g)(n) \neq 0$. To see this, we study each term in the sum

$$(f * g)(n) = \sum_{d|n} f(d)g(n/d).$$

If d < k then f(d)g(n/d) = 0 (by minimality of k); if d > k then n/d = kl/d < l, hence f(d)g(n/d) = 0 (by minimality of l). Therefore $(f*g)(n) = f(k)g(l) \neq 0$. Thus $f * g \neq 0$.

4. (20 pts.) Recall that

$$\sigma(n) = \prod_{p} (1 + p + p^2 + \dots + p^{\alpha_p}).$$

where $n = \prod_p p^{\alpha_p}$ is the prime factorization of n. Suppose $\sigma(n)$ is odd. Then each of the factors $1 + p + p^2 + \cdots + p^{\alpha_p}$ is odd. If p is a prime > 2, then each of the $\alpha_p + 1$ summands in this sum is odd; hence there has to be an odd number of summands, so α_p is even. Thus n is a square (if α_2 is even) or twice a square (if α_2 is odd).

- 5. (5+5 pts.) Let f be a number-theoretic function.
 - (a) Suppose first that g is a number-theoretic function with $f * g = \varepsilon$. Then $1 = \varepsilon(1) = (f * g)(1) = f(1)g(1)$, hence $f(1) \neq 0$. Conversely, suppose $f(1) \neq 0$. We define g(n) by recursion on n. For n = 1 we set g(1) := 1/f(1). If n > 1 and $g(1), \ldots, g(n-1)$ have already been defined, we put

$$g(n) := -\frac{1}{f(1)} \sum_{\substack{d|n \\ d>1}} f(d)g(n/d).$$

One checks immediately that then $f * g = \varepsilon$.

(b) Suppose g, h are number-theoretic functions with $f * g = \varepsilon$ and $f * h = \varepsilon$. Then

$$g=\varepsilon\ast g=(f\ast h)\ast g=f\ast (h\ast g)=f\ast (g\ast h)=(f\ast g)\ast h=\varepsilon\ast h=h.$$

6. (10 pts.) Let f be a multiplicative number-theoretic function with $f(1) \neq 0$. As shown in class, we have f(1) = 1, hence

$$(\mu f * f)(1) = \mu(1) \cdot f(1)^2 = 1 = \varepsilon(1).$$

Note also that $f^{-1}(1) = 1$ since

$$1 = \varepsilon(1) = (f^{-1} * f)(1) = f^{-1}(1) \cdot f(1)$$

Now suppose first that f is completely multiplicative. If p is a prime and $\alpha \in \mathbb{N}, \alpha > 0$, then

$$(\mu f * f)(p^{\alpha}) = \sum_{i=0}^{\alpha} \mu(p^{i}) f(p^{i}) \cdot f(p^{\alpha-i}) = f(1) \cdot f(p^{\alpha}) - f(p) \cdot f(p^{\alpha-1}),$$

and this equals $0 = \varepsilon(p^{\alpha})$, since $f(p^{\alpha}) = f(p)f(p^{\alpha-1})$ by complete multiplicativity of f. Since both $\mu f * f$ and ε are multiplicative, this suffices to show $\mu f * f = \varepsilon$, hence $f^{-1} = \mu f$. Next suppose $f^{-1} = \mu f$; then clearly for every prime p and every $\alpha > 1$ we have $f^{-1}(p^{\alpha}) = \mu(p^{\alpha})f(p^{\alpha}) = 0$ since $\mu(p^{\alpha}) = 0$. Finally, suppose $f^{-1}(p^{\alpha}) = 0$ for all prime numbers p

and $\alpha \in \mathbb{N}$, $\alpha \geq 2$. In order to check that f is completely multiplicative, it is enough to verify that $f(p^{\alpha}) = f(p)f(p^{\alpha-1})$ for every prime p and every $\alpha > 1$. (Why?) To see this we note that for those p and α we have

$$0 = \varepsilon(p^{\alpha}) = (f^{-1} * f)(p^{\alpha}) = \sum_{i=0}^{\alpha} f^{-1}(p^{i})f(p^{\alpha-i}),$$

and this sum simplifies to $f(p^{\alpha}) + f^{-1}(p)f(p^{\alpha-1})$. Also

$$0 = \varepsilon(p) = (f^{-1} * f)(p) = f(p) + f^{-1}(p),$$

hence $f^{-1}(p) = -f(p)$ and thus

$$f(p^{\alpha}) - f(p)f(p^{\alpha-1}) = f(p^{\alpha}) + f^{-1}(p)f(p^{\alpha-1}) = 0,$$

so $f(p^{\alpha}) = f(p)f(p^{\alpha-1})$ as required.

7. (20 pts. extra credit.) Since all the functions involved are multiplicative, it is enough to show, for every prime p and every $\alpha \in \mathbb{N}$, that

$$(\tau^3 * 1)(p^{\alpha}) = ((\tau * 1)(p^{\alpha}))^2.$$

Now

$$(\tau * 1)(p^{\alpha}) = \sum_{i=0}^{\alpha} \tau(p^i) = \sum_{i=0}^{\alpha} i + 1 = \sum_{i=1}^{\alpha+1} i.$$

Also

$$(\tau^3 * 1)(p^{\alpha}) = \sum_{i=0}^{\alpha} \tau(p^i)^3 = \sum_{i=0}^{\alpha} (i+1)^3 = \sum_{i=1}^{\alpha+1} i^3.$$

So (rewriting $\alpha + 1$ as m) we have to show, for every m > 0, that

 $(1 + 2 + \dots + m)^2 = 1^3 + 2^3 + 3^3 + \dots + m^3.$

We do this by induction on m. The base case m = 1 is trivial. Suppose we have shown the claim for some value of m. Then

$$(1+2+\dots+m+(m+1))^2 = (1+2+\dots+m)^2 + (m+1)^2 + 2(1+2+\dots+m)(m+1) = (1^3+2^3+3^3+\dots+m^3) + (m+1)((m+1)+2(1+2+\dots+m)) = (1^3+2^3+3^3+\dots+m^3) + (m+1)((m+1)+m(m+1)) = 1^3+2^3+3^3+\dots+m^3+(m+1)^3,$$

where in the second equality we used the inductive hypothesis, and in the third the well-known formula for $1 + 2 + \cdots + m$.

Total: 100 pts. + 20 pts. extra credit.