

OH T44, 30

Th - go over mid exam

Counting edges

$K_n =$ Complete graph on n -vertices

one edge between all pairs of vertices

simple, undirected graph



K_1

$$E_1 = 0$$



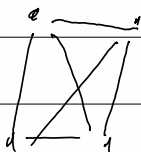
K_2

$$E_2 = 1$$

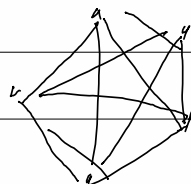


K_3

$$E_3 = 3$$



K_4

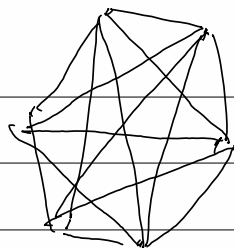


K_5

$$E_5 = 10$$

$$E_4 = 6$$

$$E_n = \# \text{ edges in } K_n$$



formula for E_n ?

$$n = 2, 3, 4, 5$$

$$E_n = 1, 3, 6, 10, 15$$

$$\begin{array}{cccc} \cup & \cup & \cup & \cup \\ 2 & 3 & 4 & 5 \end{array}$$

$$10 \pm 1 + 3 + 6$$

$$0 \neq 3 + 1$$

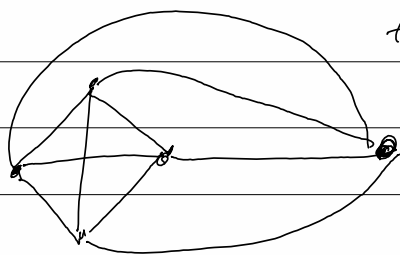
$$3 \neq 1$$

$$1, 2, 4, 8, 16$$

$$1, 2, 3, 4, 5$$

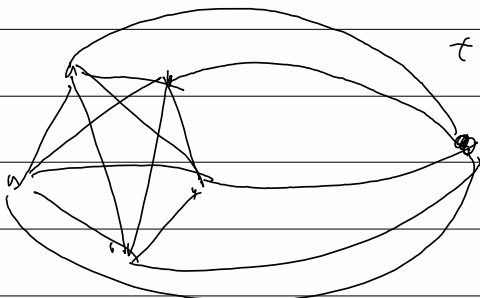
K_5

K_{4+1}



+4

K_6



+5

$$E_{n+1} = n + E_n$$

$$E_2 = 1$$

→

$$E_1 = 0$$

$$E_2 = 1$$

$$E_{n+1} = n + E_n$$

$$1, 2, 4, 8, 16$$

$$q_n = 2q_{n-1}, \quad q_0 > 1$$

$$q_n \geq 2^n$$

$$E_n = ? \quad \text{not involving } E_{n-1}$$

$$E_n = (n-1) + E_{n-1}$$

$$= (n-1) + ((n-2) + E_{n-2})$$

$$= (n-1) + (n-2) + \left(\underbrace{(n-3) + E_{n-3}}_{n-1} \right)$$

$$= (n-1) + (n-2) + \dots + 1 + E_1$$

$$= \sum_{k=1}^{n-1} k$$

$$\text{Prop. } E_n = \sum_{k=1}^{n-1} k$$

pf, induction

Gauss adding

up $1+2+\dots+n$

$$E_n = \sum_{k=1}^{n-1} k$$

$$= \frac{n(n-1)}{2}$$

$$1 + 2 + 3 + 4 + \dots + (n-1) + (n-2) + (n-1)$$

$$\begin{aligned} 1 + (n-1) &= n \\ 2 + (n-2) &= n \\ 3 + (n-3) &= n \\ &\vdots \\ &\vdots \end{aligned}$$

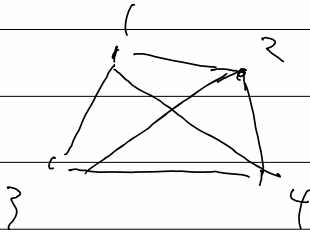
$$\frac{n \cdot \frac{n-1}{2}}$$

$$E_n = \frac{n(n-1)}{2} = C(n, 2) = \frac{n!}{2!(n-2)!} = \frac{n(n-1)}{2}$$

$$E_n = \underline{C(n, 2)}$$

K_n — vertices, $\{1, \dots, n\}$

— edges $\leftrightarrow$ pairs of vertices



n people, all have to shake each other's hand

$$\sum_{k=1}^{n-1} k = C(n, 2)$$

✓

#edges in K_n

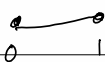
$Q_n =$ hypercube of dimension n

$n=0$



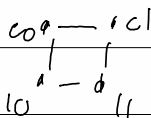
$E_0 = 0$

$n=1$



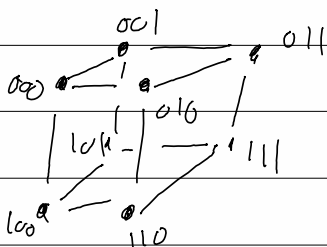
$E_1 = 1$

$n=2$



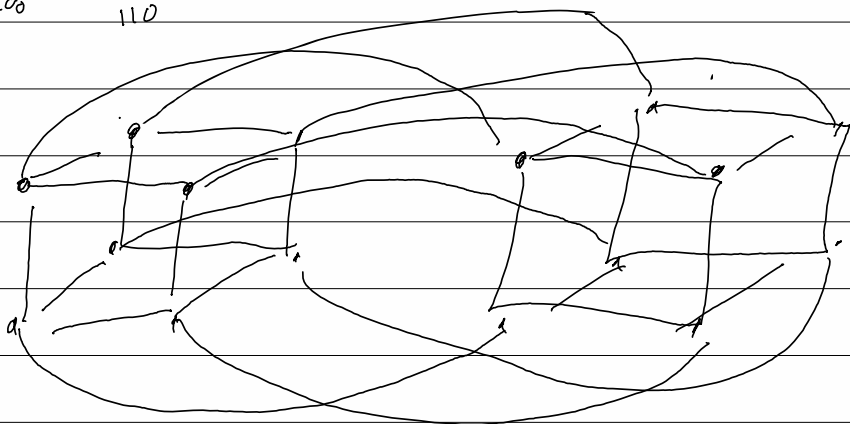
$E_2 = 4$

$n=3$

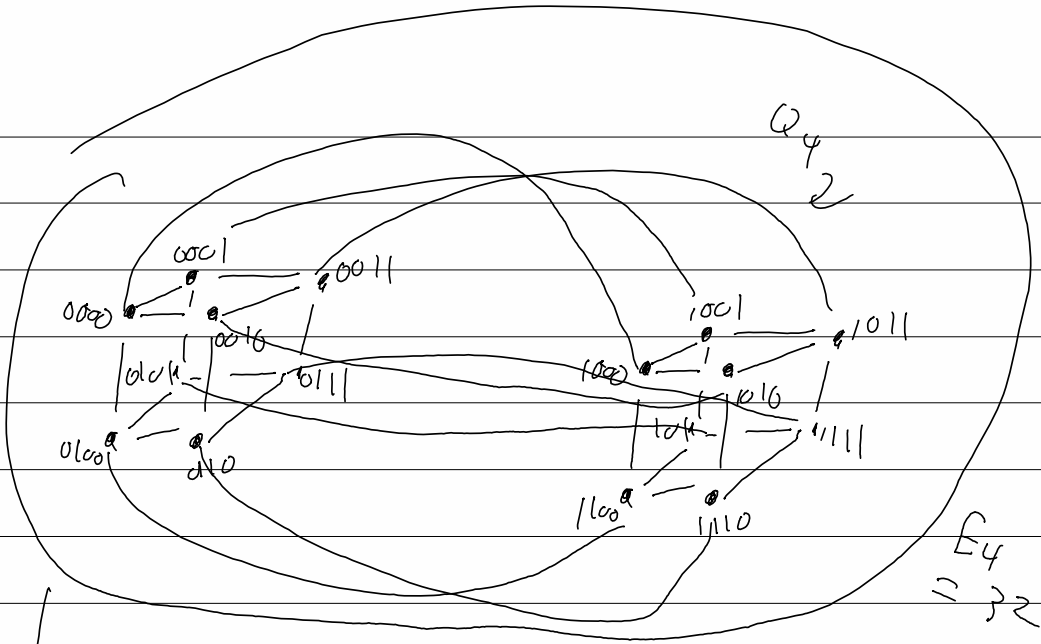


$E_3 = 12$

$n=4$



$Q_n = \begin{cases} \text{vertices} & n\text{-bit strings } 000\dots 01 \\ \text{edges} & \text{edges between 2 } n\text{-bit strings, if they} \\ & \text{differ by one bit} \\ & \text{differ in exactly one place.} \end{cases}$



$V_n = \# \text{ vertices in } Q_n$

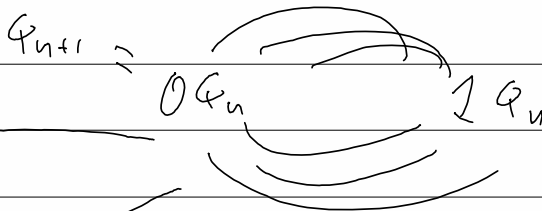
$E_n = \# \text{ edges in } Q_n$

$$V_n = \# \text{ n-bit strings} = 2^n$$

n	0	1	2	3	4
E_n	0	1	4	12	32

prop. $E_n = n2^{n-1}$

$3 \rightarrow 4$



$$E_{n+1} = E_n + E_n + V_n$$

$$E_{n+1} = 2E_n + V_n$$

$$= 2E_n + 2^n$$

prove

$n 2^{n-1}$ edges

$$E_{n+1} = 2E_n + 2^n$$

$$\text{and } E_1 = 1$$

$$E_n = 2E_{n-1} + 2^{n-1}$$

$$= 2(2E_{n-2} + 2^{n-2}) + 2^{n-1}$$

$$= 2^2 E_{n-2} + 2^{n-1} + 2^{n-1}$$

$$= 2^2 (2E_{n-3} + 2^{n-3}) + 2^{n-1} + 2^{n-1}$$

$$= 2^3 E_{n-3} + 2^{n-1} + 2^{n-1} + 2^{n-1}$$

$= \dots$

$$= 2^k E_{n-k} + k 2^{n-1}$$

$$\downarrow k = (n-1)$$

$$= 2^{n-1} E_1 + (n-1) 2^{n-1}$$

$$\approx n 2^{n-1}$$

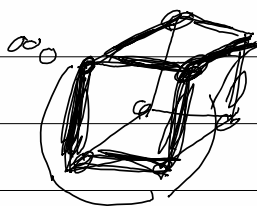
each vertex

in G_n has

n edges out of it

coll 7 edges

$$n 2^n / 2 = n 2^{n-1}$$

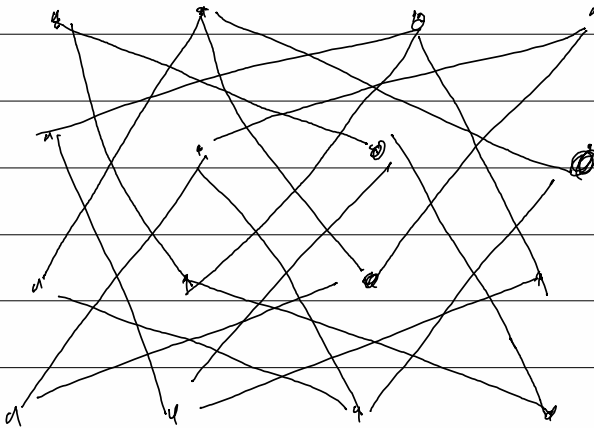


Hamming error correction
Hamming distance

Knight's graph

$V =$ hex- grid squares
pairs (i, j) $1 \leq j \leq 4$

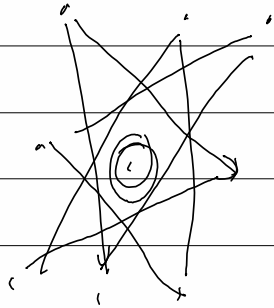
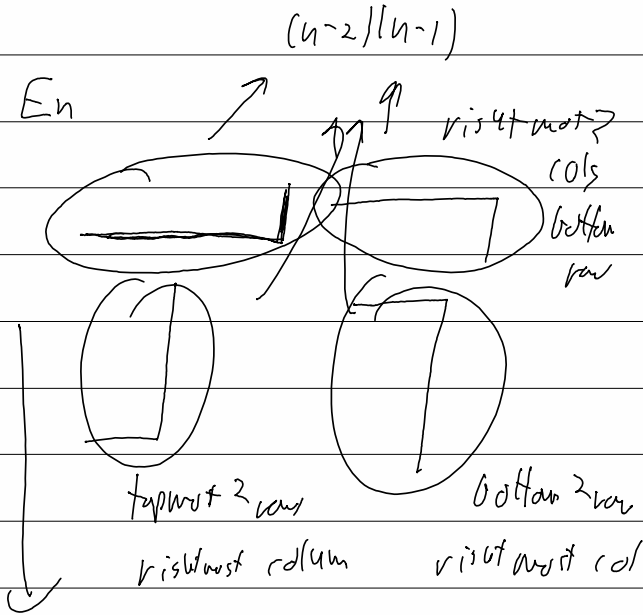
$E =$ 2 vert conn if they differ by a
knight's move



N_n , Hamiltonian path $\leftrightarrow$ Knight's tour

$$n \geq 3$$

$$V_n = n^2$$



he Kurat's four!

vertices have this as a legal move?



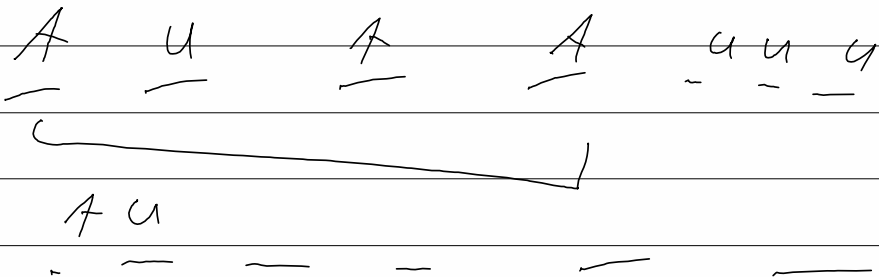
$$E_n \approx 4(n-1)(n-2)$$

Transfers available/unavailable

Inventory w/ 8 items
either available or unavailable

50 available

wt₃, There are at least 2 available items, which
are exactly 2 places apart



AAAAAAAAUUUUUUUU

A A A U A A A

f is item, $f(1) f(2)$
 so available $\begin{matrix} 1 \\ 1 \\ 0 \end{matrix}$

$A = \text{Set of available items}$

$\Rightarrow$ so

$f: A \rightarrow \{1, \dots, 88\}$

$a \mapsto \text{its position}$

Goal: there are a, a' in A

$$f(a) - f(a') = 9$$

~~$A \times A \rightarrow \{1, \dots, 88\}$
 $(a, b) \mapsto f(a) - f(b)$~~

ATMATAA
ATMATAA

label available items

$1, 2, \dots, f_0$

$f: \{1, \dots, f_0\} \rightarrow \{1, \dots, f_0\}$

$f(i) = \text{pos of item } i$

Goal. $f(i) - f(j) = g$ some $i \neq j$

$S_1 = \{f(1), f(2), \dots, f(f_0)\}$ all positions of available items
 $S_2 = \{f(1) + g, f(2) + g, \dots, f(f_0) + g\}$ all possible available items

want to show $S_1 \cap S_2 \neq \emptyset$ Shurafa by P

let's say n in $S_1 \cap S_2$

$n \in S_1 \therefore n = f(i)$ some i

$n \in S_2 \therefore n = f(j) + g$ some j

$f(i) = f(j) + g$

holes
range $(1, \dots, g)$
between them
hika

paid } times

3rd form

$f(\text{pigeon})$ the output of this pigeon is in

$$x \xrightarrow{f} f(x)$$

$$X \xrightarrow{f} Y$$

pigeons

holes

$$|X| = n$$

$$|Y| = m$$

there are $\lceil \frac{n}{m} \rceil$ distinct elts of X

$$x_1, \dots, x_{\lceil \frac{n}{m} \rceil}$$

$$\text{s.t. } f(x_1) = \dots = f(x_{\lceil \frac{n}{m} \rceil})$$

$$n > m \quad \lceil \frac{n}{m} \rceil \geq 2$$

$$n = 17$$

$$\lceil \frac{17}{5} \rceil = 4$$

$$m = 5$$

