

Lecture 18 Elliptic Curves § 6.1

Why elliptic Curves

- They are groups w/ more confusing addition laws than $(\mathbb{Z}/p\mathbb{Z})^*$ \Rightarrow harder DLP.

Most modern cryptography uses Elliptic curve Diffie-Hellman to exchange a symmetric encryption key.

- Curves with a B -smooth number of points can be used to factor integers & compute discrete logs.
Lenstra's Algorithm

For instance, if we know of an elliptic curve over \mathbb{F}_p w/ a B -smooth # of points then we can use an oracle for Diffie-Hellman to solve Discrete Log.

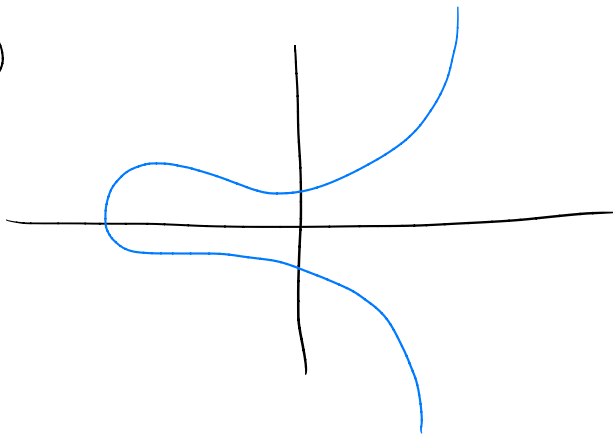
Elliptic curves

Def A Weierstrass cubic is an eqn of the form

$$y^2 = x^3 + Ax + B$$

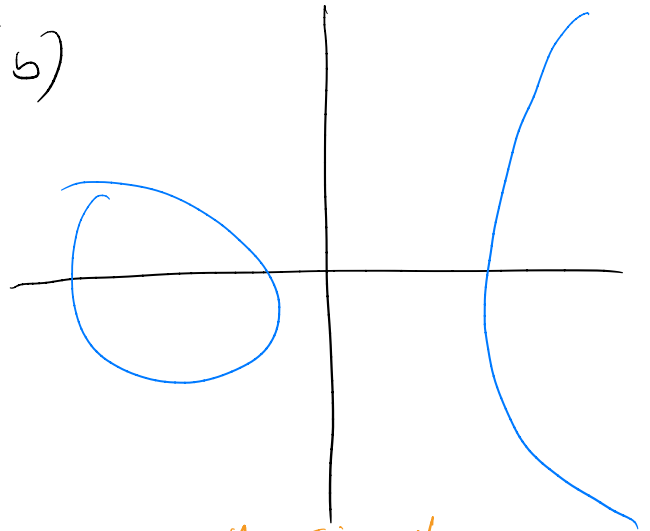
The pairs $(x, y) \in \mathbb{R}^2$ satisfying this eqn look like one of the following:

(a)



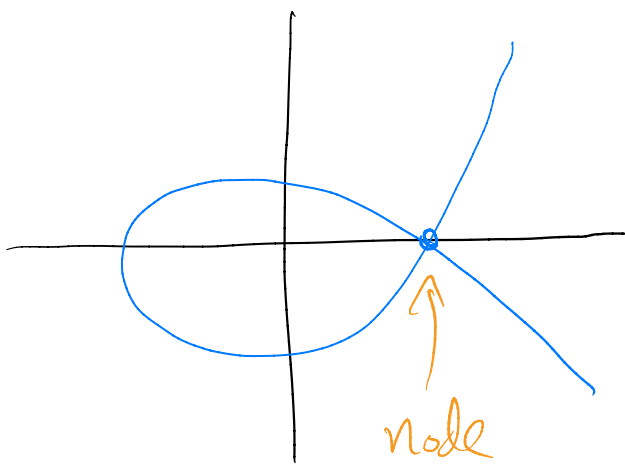
nonsingular

(b)



nonsingular

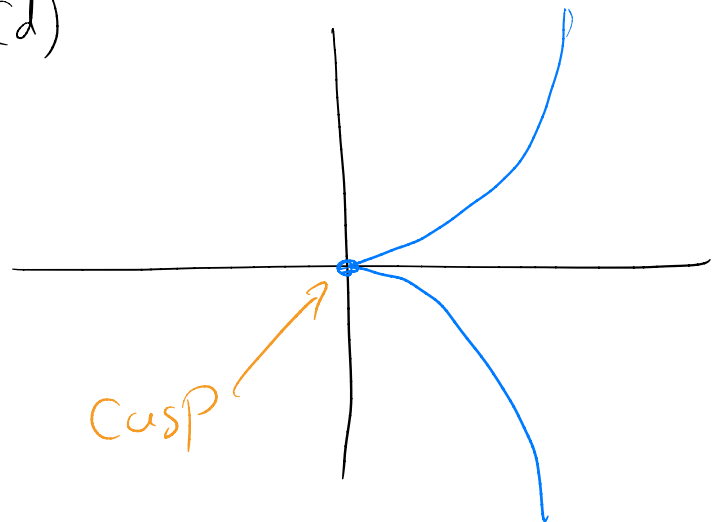
(c)



node

singular

(d)



cusp

singular

The set of solutions of a nonsingular Weierstrass eqn is an elliptic curve.

(a) & (b) are nonsingular & are elliptic curves

(c) has a node called a "nodal cubic"

(d) has a cusp called a "cuspidal cubic"

(c) & (d) are not elliptic curves.

A Weierstrass equation $y^2 = x^3 + Ax + B$

is nonsingular if and only if $f(x) = x^3 + Ax + B$ has no repeated roots.

Def Suppose $p(x)$ is a polynomial which factors as $p(x) = (x-r_1) \cdots (x-r_n)$. The discriminant of $p(x)$ is $D = \prod_{i \neq j} (r_i - r_j)$.

If the roots r_i are distinct, $D \neq 0$.

If there is a repetition, then $D = 0$.

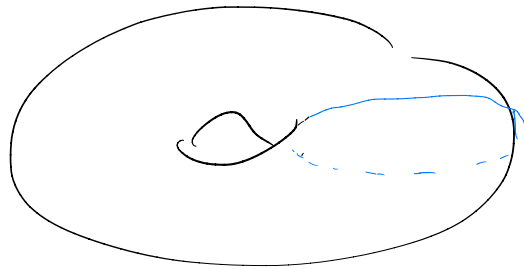
Conveniently, D can be expressed in terms of the coefficients of $p(x)$!

Fact If $f(x) = x^3 + Ax + B$ then

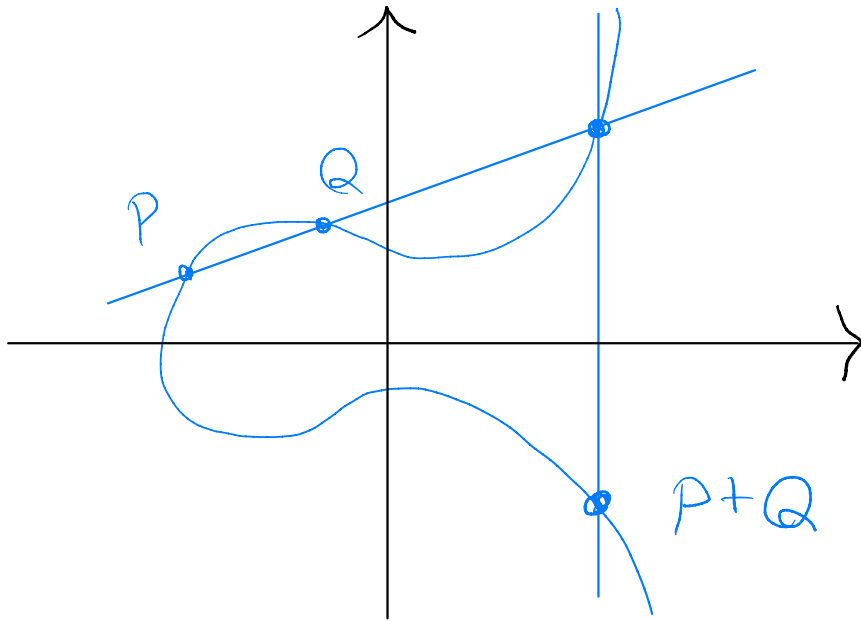
$$D = -4A^3 - 27B^3.$$

Thus, any Weierstrass equ w/ $D = -4A^3 - 27B^3 \neq 0$ is an elliptic curve.

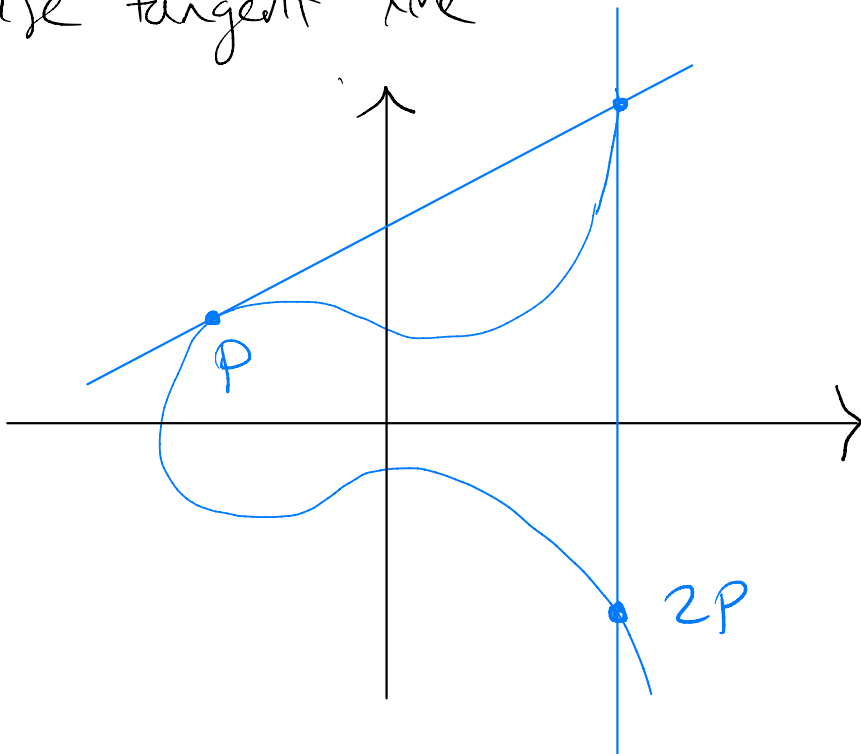
Geometrically, the complex solutions look like a donut:



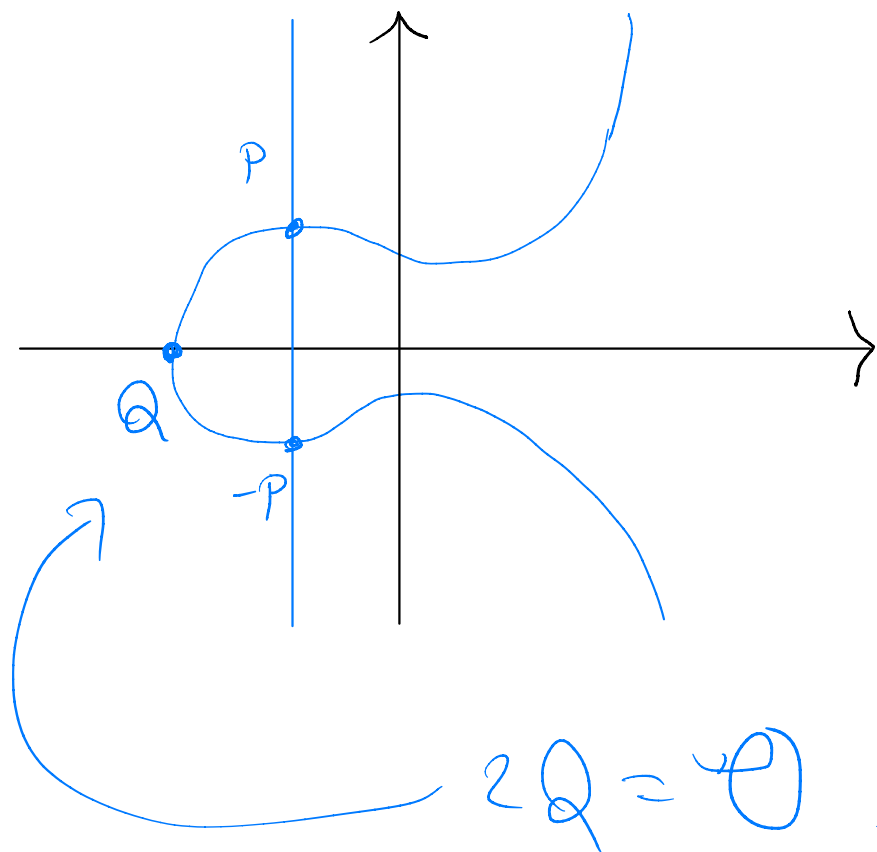
We can add points on an elliptic curve:



Each line intersects the curve in exactly 3 points counting repetitions. To double a point, use tangent line instead of secant line.



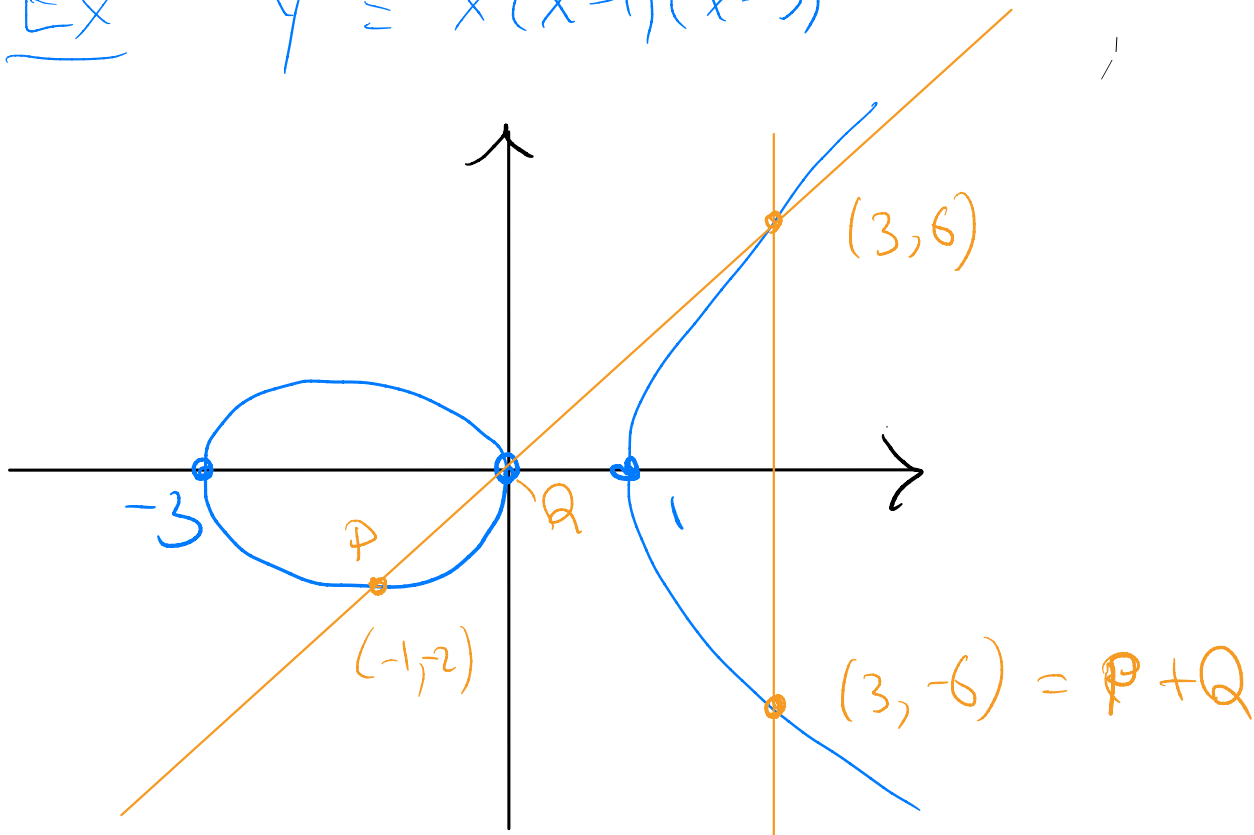
Also, vertical lines only intersect curve twice, so we add an extra point \mathcal{O} and say it is on every vertical line. $P + \mathcal{O} = P$, so \mathcal{O} is the identity.



Fact This addition is a group structure on the points of the curve.

Rmk We don't need a Weierstrass cubic, any cubic w/ nonvanishing discriminant works.

Ex $y^2 = x(x-1)(x-3)$



$P = (-1, -2)$ $Q = (0, 0)$. The line through P & Q is $y = 2x$, so the third point of intersection is

$$\Rightarrow 4x^2 = x^3 + 2x^2 - 3x$$

$$x^3 - 2x^2 - 3x = 0$$

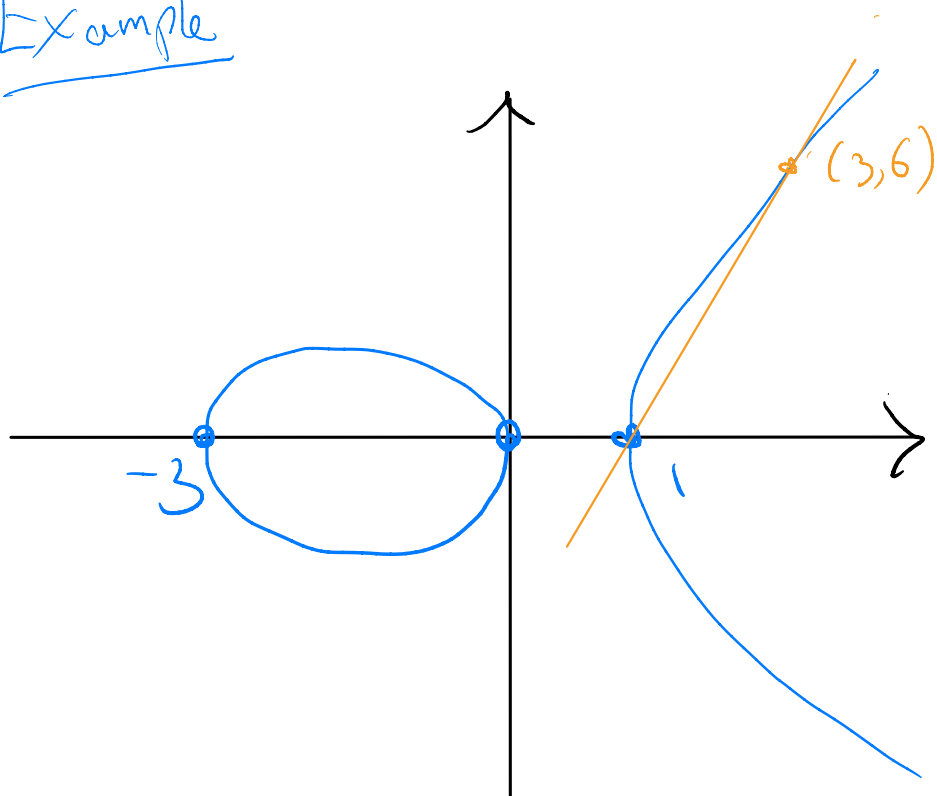
$$\underbrace{x(x+1)}_{\text{known roots}} \underbrace{(x-3)}_{\text{third intersection point}} = 0$$

known roots

third intersection point

$$y = 2x \Rightarrow (3, 6) \rightarrow P + Q = (3, -6)$$

Example



$$2. (3, 6) = (1, 0)$$

Tangent line

$$y^2 = x(x-1)(x+3) = x^3 + 2x^2 - 3x$$

implicit differentiation $\Rightarrow 2yy' = 3x^2 + 4x - 3$

plug in $(x, y) = (3, 6) \Rightarrow 12y' = 27 + 12 - 3 = 36$

$$\Rightarrow y' = 3$$

$$\Rightarrow \text{tangent line is } y = 3(x-3) + 6 \\ = 3x - 3$$

Plug tangent line back into elliptic curve eqn:

$$(3x-3)^2 = x^3 + 2x^2 - 3x$$

$$9x^2 - 18x + 9 = x^3 + 2x^2 - 3x$$

$$x^3 - 7x^2 + 15x - 9 = 0$$

Root sum must be 7. b/c

$$(x-a)(x-b)(x-c) = x^3 - (a+b+c)x^2 + \dots$$

$$3 + 3 + - = 7 \Rightarrow \text{third root is } 1.$$

$$\Rightarrow x=1 \quad \& \quad y = 3(1) - 3 = 0.$$

Reflect wrt x-axis gives $(1, -0) = \boxed{(1, 0)}$