

# Math 116 Spring 2022

## Homework 2

Due Friday, April 15th

### Instructions for SageMath

You may use `factor` to check your answer to Problem 1.

You can use `k.<a> = FiniteField(p, impl='modn')` to set  $k$  to be  $\mathbb{F}_p$  and  $a$  the multiplicative unit in  $k$  usually represented by the integer 1.

If  $b \in \mathbb{F}_p$  then the `multiplicative_order` method of  $b$  returns the multiplicative order of  $b$  in  $\mathbb{F}_p$ . For example, if you want to find the multiplicative order of 2 in  $\mathbb{F}_7$ , then

```
k.<a> = FiniteField(7, impl='modn')
(2*a).multiplicative_order()
```

Here we write `(2*a)` to coerce 2 into an element of  $\mathbb{F}_7$ . This can be used to check your answer to Problem 4.

The `log` method of  $b$  returns an integer  $x$  such that  $m = b^x \pmod{p}$ . For example, if you want to find the discrete logarithm  $\log_3(5)$  in  $\mathbb{F}_7$ , then use the code

```
k.<a> = FiniteField(7, impl='modn')
(5*a).log(3)
```

This can be used to check your answer for Problem 7, and is needed for Problem 8.

### An aside on syntax for people with Python experience

You may be upset that `k.<a> = blah` is not syntactically valid Python. Sage uses an IPython input transformation hook which rewrites

```
x.<y, z, w> = blah
```

into the syntactically valid Python code

```
x = blah; (y, z, w) = x._first_ngens(3))
```

If you are ever confused by the creepy IPython input rewrite hooks, my suggestion is to intentionally force a syntax error. The syntax error will show the rewritten code. For instance typing `726 +` results in the syntax error:

```
Integer(726) +
SyntaxError: invalid syntax
```

we see that Sage wrapped 726 with `Integer`. Likewise `726.2 +` results in:

```
RealNumber("726.2") +
SyntaxError: invalid syntax
```

This can make it a lot easier to understand what is going on.

## Primes and Finite Fields

**Problem 1.** (1.30(c), by hand) Compute  $\text{ord}_p(46375)$  for  $p$  equal to 3, 5, 7, and 11.

**Problem 2.** (1.31) Let  $p$  be a prime number. Prove that  $\text{ord}_p$  has the following properties.

- (a)  $\text{ord}_p(ab) = \text{ord}_p(a) + \text{ord}_p(b)$ . (Thus  $\text{ord}_p$  resembles the logarithm function, since it converts multiplication into addition!)
- (b)  $\text{ord}_p(a + b) \geq \min\{\text{ord}_p(a), \text{ord}_p(b)\}$ .
- (c) If  $\text{ord}_p(a) \neq \text{ord}_p(b)$ , then  $\text{ord}_p(a + b) = \min\{\text{ord}_p(a), \text{ord}_p(b)\}$ .

A function satisfying these properties is called a valuation.

**Problem 3.** (1.32(a), by hand) Let  $p = 47$  and let  $a = 11$ . Compute  $a^{-1} \pmod p$  in two ways:

- (a) Use the extended Euclidean algorithm.
- (b) Use the fast power algorithm and Fermat's little theorem.

**Problem 4.** (1.34, by hand) Recall that  $g$  is called a primitive root modulo  $p$  if the powers of  $g$  give all nonzero elements of  $\mathbb{F}_p$ .

- (a) For which of the following is 2 a primitive root mod  $p$ ?
  - (i)  $p = 7$    (ii)  $p = 13$    (iii)  $p = 19$    (iv)  $p = 23$
- (b) Find a primitive root for  $p = 29$  and for  $p = 41$ .
- (c) Find all primitive roots modulo 11. Verify that there are exactly  $\phi(10)$  of them, as asserted in Remark 1.32.

## Symmetric Ciphers

**Problem 5.** (1.43(a)(c), by hand (you can use Sage for `xgcd`)) This problem is about the affine cipher. The affine cipher has key given by a pair of integers  $k = (k_1, k_2)$ . The encryption and decryption functions are given by

$$e_k(m) \equiv k_1 \cdot m + k_2 \pmod p, \quad \text{and} \quad d_k(c) \equiv k_1^{-1} \cdot (c - k_2) \pmod p.$$

- (a) Let  $p = 541$  and let the key be  $k = (34, 71)$ . Encrypt the message  $m = 204$ . Decrypt the ciphertext  $c = 431$ .
- (b) Alice and Bob decide to use the prime  $p = 601$  for their affine cipher. The value of  $p$  is public knowledge, and Eve intercepts the ciphertexts  $c_1 = 324$  and  $c_2 = 381$  and also manages to find out that the corresponding plaintexts are  $m_1 = 387$  and  $m_2 = 491$ . Determine the private key and then use it to encrypt the message  $m_3 = 173$ .

**Problem 6.** (1.44(a)(c)) Consider the Hill cipher, defined by the same equations as the affine cipher except where now  $m$ ,  $c$ , and  $k_2$  are vectors of dimension  $n$  and  $k_1$  is an  $n \times n$  matrix.

- (a) Let  $p = 7$ ,  $k_1 = \frac{1}{2} \begin{pmatrix} 3 & \\ & 2 \end{pmatrix}$  and  $k_2 = \begin{pmatrix} 5 \\ 4 \end{pmatrix}$ .
- (i) Encrypt  $m = \begin{pmatrix} 2 \\ 1 \end{pmatrix}$ .
  - (ii) What is the matrix  $k^{-1}$  used for decryption?
  - (iii) Decrypt the message  $c = \begin{pmatrix} 3 \\ 5 \end{pmatrix}$ ?
- (b) The following plaintext/ciphertext pairs were generated using a Hill cipher with the prime  $p = 11$ . Find the keys  $k_1$  and  $k_2$ .

$$m_1 = \begin{pmatrix} 5 \\ 4 \end{pmatrix}, \quad c_1 = \begin{pmatrix} 1 \\ 8 \end{pmatrix}, \quad m_2 = \begin{pmatrix} 8 \\ 10 \end{pmatrix}, \quad c_2 = \begin{pmatrix} 8 \\ 5 \end{pmatrix}, \quad m_3 = \begin{pmatrix} 7 \\ 1 \end{pmatrix}, \quad c_3 = \begin{pmatrix} 8 \\ 7 \end{pmatrix},$$

## Diffie–Hellman

**Problem 7.** (2.4(a), by hand) Solve the congruence  $2^x \equiv 13 \pmod{23}$ .

**Problem 8.** (2.6, with sage) Alice and Bob agree to use the prime  $p = 1373$  and the base  $g = 2$  for a Diffie–Hellman key exchange. Alice sends Bob the value  $A = 974$ . Bob asks your assistance, so you tell him to use the secret exponent  $b = 871$ . What value  $B$  should Bob send to Alice, and what is their secret shared value? Can you figure out Alice’s secret exponent?

**Problem 9.** (2.7) Let  $p$  be a prime and let  $g$  be an integer. The *Decision Diffie–Hellman Problem* is as follows. Suppose you are given three numbers  $A$ ,  $B$ , and  $C$ . Suppose that

$$A \equiv g^a \pmod{p} \quad \text{and} \quad B \equiv g^b \pmod{p}$$

but you do not know  $a$  and  $b$ . The goal is to determine whether  $C \equiv g^{ab} \pmod{p}$ .

- (a) Prove that an algorithm to solve Diffie–Hellman can be used to solve Decision Diffie–Hellman.
- (b) Do you think that the decision Diffie–Hellman problem is hard or easy? Why? See Exercise 6.40 for a related example in which the decision problem is easy, but it is believed that the associated computational problem is hard.