Move-reduced graphs on a torus vs positroid Catalan numbers

Pavel Galashin (UCLA) based on joint works with Terrence George and Thomas Lam



(scan to play the game)

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Let n≥ 2. Š<sub>n</sub> := {f : Z → Z bijection | f(i + n) = f(i) + n ∀i ∈ Z}.
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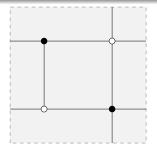
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## Bipartite graphs on a torus

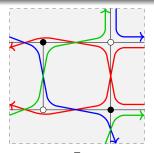
### Definition

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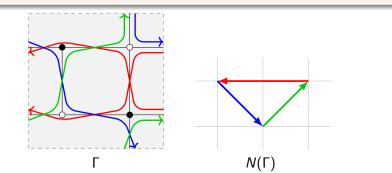


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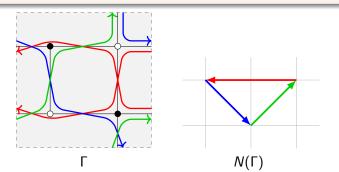
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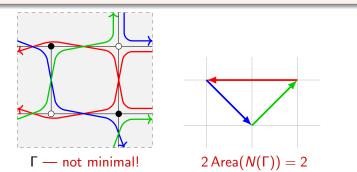
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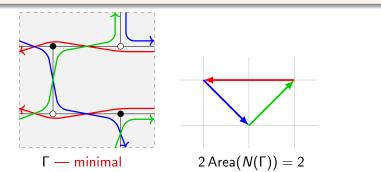
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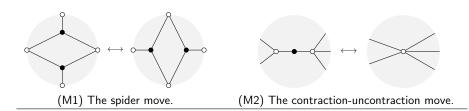
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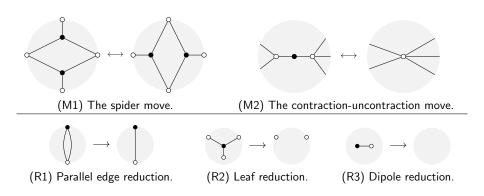
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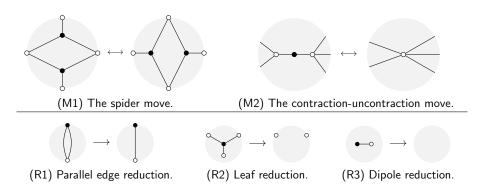
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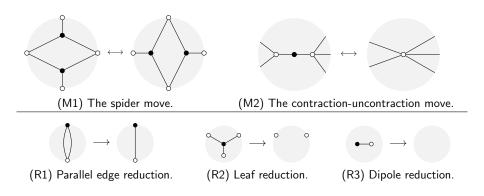
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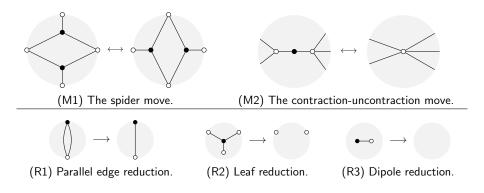
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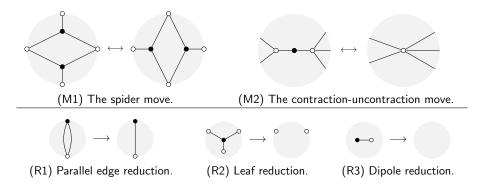
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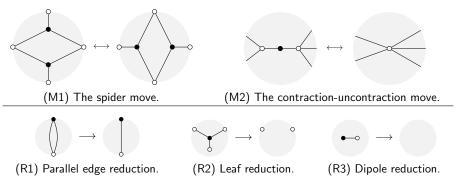
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- Question: When are two move-reduced graphs move-equivalent?



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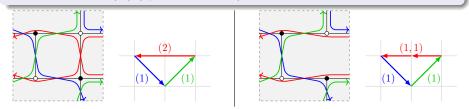
faces, where the sum is over zig-zag paths of  $\boldsymbol{\Gamma}.$ 

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  - µ(Γ) = µ(Γ'), where µ(Γ) ∈ ℤ/dℤ—modular invariant,
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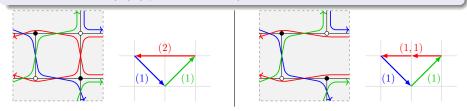
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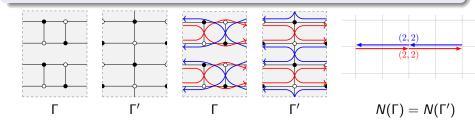


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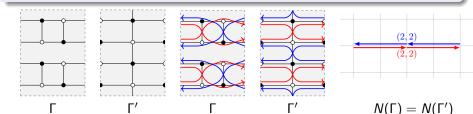
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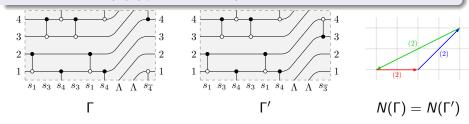
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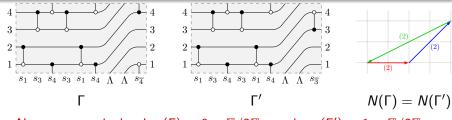
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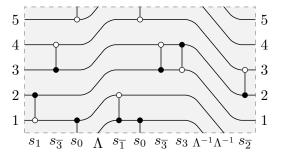
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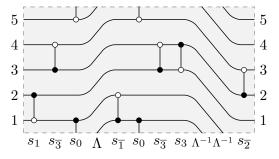
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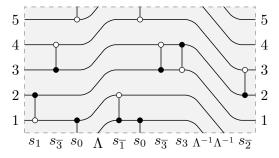


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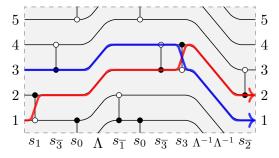
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Associate a white-black bridge to s<sub>i</sub> and a black-white bridge to s<sub>i</sub>. Get a bicolored graph Γ on a torus.



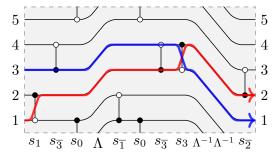
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#### Theorem (G.–George; see also Fock–Marshakov '16)

- If f, f' are c-reduced then  $\Gamma$  is move-reduced.
- Any move-reduced graph  $\Gamma$  can be obtained in this way.

Recall:  $\tilde{S}_{k,n} := \{ f \in \tilde{S}_n \mid \frac{1}{n} \sum_{i=1}^n (f(i) - i) = k \}.$ 

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Bounded affine permutations:  $\mathcal{B}_{k,n} := \{ f \in \tilde{S}_{k,n} \mid i \leq f(i) \leq i+n \text{ for all } i \in \mathbb{Z} \}.$ 

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Let  $f = \Lambda_n^k \in \mathcal{B}_{k,n}$ , i.e., f(i) = i + k for all  $i \in \mathbb{Z}$ . Assume ncyc(f) = gcd(k, n) = 1.

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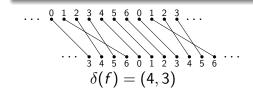
#### Question

What does this have to do with Dyck paths?

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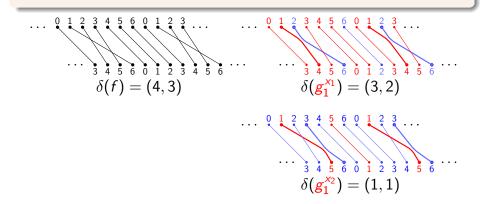
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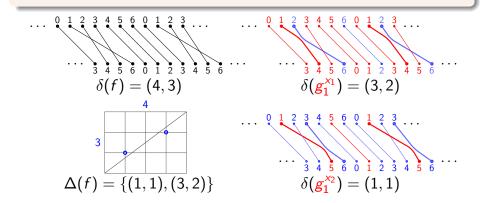
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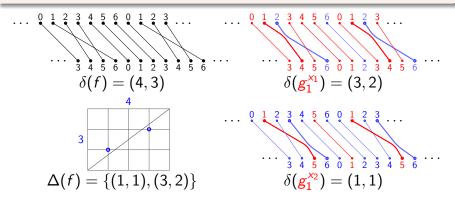
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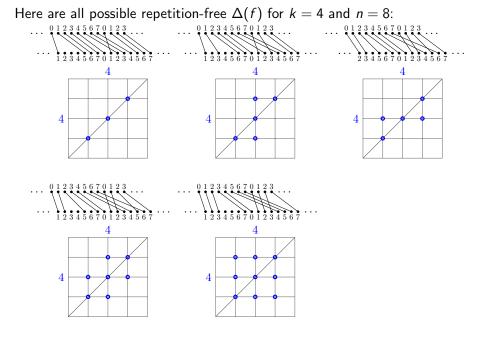
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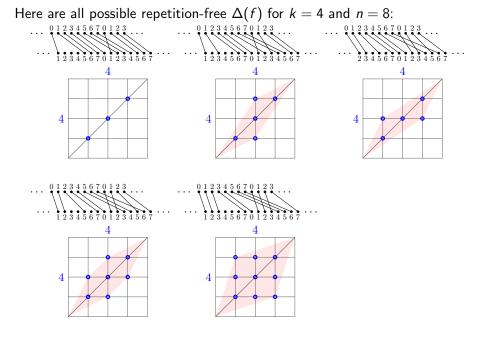
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- f is called repetition-free if all points in  $\Delta(f)$  are distinct.

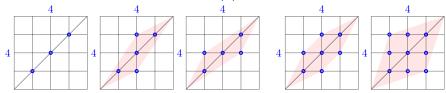


Here are all possible repetition-free  $\Delta(f)$  for k = 4 and n = 8:

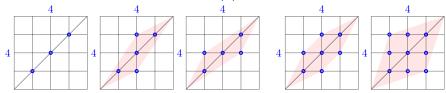




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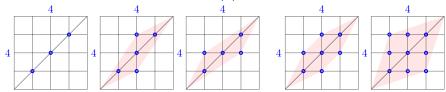


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 Δ ⊆ [n − k − 1] × [k − 1] arises as Δ(f) for some repetition-free f ∈ B<sub>k,n</sub> if and only if Δ ⊔ {(0,0), (n − k, k)} = P ∩ Z<sup>2</sup>

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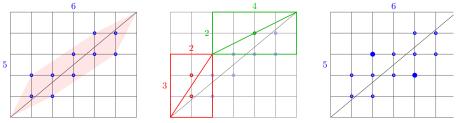
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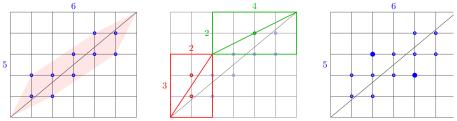
 $\# \operatorname{Dyck}(\Delta(f)) = \# \operatorname{Dyck}(\Delta(g_1)) \cdot \# \operatorname{Dyck}(\Delta(g_2)) + \# \operatorname{Dyck}(\Delta(s_i f s_i))$ 

### If f is repetition-free, then $C_f = \# \operatorname{Dyck}(\Delta(f))$ .

Recall: If  $\ell(s_i f s_i) = \ell(f) + 2$  then

$$C_f = C_{g_1} \cdot C_{g_2} + C_{s_i f s_i},$$

where  $g_1, g_2$  are the two cycles of  $s_i f \stackrel{c}{\sim} fs_i$ .



 $\#\operatorname{Dyck}(\Delta(f)) = \#\operatorname{Dyck}(\Delta(g_1)) \cdot \#\operatorname{Dyck}(\Delta(g_2)) + \#\operatorname{Dyck}(\Delta(s_i f s_i))$ 

# **Thanks!**