

# HEDETNIEMI'S CONJECTURE

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## 1. INTRODUCTION

A **graph**  $G$  is a set  $V(G)$  equipped with a symmetric relation  $E(G)$  (note that  $G$  may have loops). A function  $\phi : V(G) \rightarrow V(H)$  is a **homomorphism** if it preserves the edge relation, i.e. if  $(v, w) \in E(G)$ , then  $(\phi(v), \phi(w)) \in E(H)$ . Let  $\text{Hom}(G, H)$  denote the set of homomorphisms from  $G$  to  $H$ . We write  $G \leq H$  if  $\text{Hom}(G, H)$  is nonempty (this is usually denoted  $G \rightarrow H$  in finite graph theory).

Given a cardinal number  $n$  (which may be infinite), let  $K_n$  denote the **complete graph** on  $n$  vertices:  $V(K_n) = n$ , and  $(k, l) \in E(K_n)$  iff  $k \neq l$  (we follow the set-theoretic convention of writing  $n$  for the set of  $k$  with  $0 \leq k < n$ ). Given a graph  $G$ , a **proper  $n$ -colouring** of  $G$  is a homomorphism from  $G$  to  $K_n$ , and the **chromatic number** of  $G$ , denoted  $\chi(G)$ , is the minimal  $n$  such that  $G \leq K_n$  (we write  $\chi(G) = \infty$  if no such  $n$  exists, which occurs exactly when  $G$  is not simple).

Given graphs  $G$  and  $H$ , the **categorical product**  $G \times H$  is the graph with  $V(G \times H) = V(G) \times V(H)$  such that  $((v, w), (v', w')) \in E(G \times H)$  iff  $(v, v') \in E(G)$  and  $(w, w') \in E(H)$ . There are projections  $G \times H \rightarrow G$  and  $G \times H \rightarrow H$ , so  $\chi(G \times H)$  is bounded above by both  $\chi(G)$  and  $\chi(H)$ . Hedetniemi asked in his PhD thesis if this bound is optimal:

**Conjecture 1** (Hedetniemi [Hed66]). *Let  $G$  and  $H$  be finite graphs. Then*

$$(*) \quad \chi(G \times H) = \min\{\chi(G), \chi(H)\}.$$

Let us make a few remarks.

- (1) If either  $G$  or  $H$  is not simple, then  $(*)$  holds trivially, since if  $G$  is not simple, then  $H \leq G \times H$ . So the conjecture is really about simple graphs.
- (2) The equality  $(*)$  is known to hold in the following cases:
  - (a)  $\chi(G \times H) \leq 2$  (easy).
  - (b)  $\chi(G \times H) = 3$ , due to El-Zahar and Sauer [ES85].
  - (c) When  $\chi(G)$  is infinite and  $\chi(H)$  is finite, due to Hajnal [Haj85] (for possibly infinite graphs).

- (d) For Borel chromatic numbers of analytic graphs, (\*) holds when  $\chi_B(G)$  and  $\chi_B(H)$  are uncountable, as a corollary of the  $G_0$ -dichotomy [KST99, 6.11].
- (3) The generalisation to infinite graphs does not hold in general. This was first explicitly published by Hajnal [Haj85] (in the same issue of *Combinatorica* as El-Sahar and Zauer!).

In May 2019, Yaroslav Shitov refuted Hedetniemi's conjecture:

**Theorem 1** (Shitov [Shi19]). *Hedetniemi's conjecture is false.*

We present his proof below.

## 2. PRELIMINARIES

Given graphs  $G$  and  $H$ , the **exponential graph**  $H^G$  is the graph with  $V(H^G) = V(H)^{V(G)}$ , such that  $(\alpha, \beta) \in E(H^G)$  iff for all  $(v, w) \in E(G)$ , we have  $(\alpha(v), \beta(w)) \in E(H)$ . It satisfies the following property:

$$\text{Hom}(G \times H, K) \cong \text{Hom}(G, K^H)$$

We can use this to simplify Hedetniemi's conjecture. The conjecture says that if  $G$  is a finite graph with  $\chi(G) > n$ , then any of the following equivalent statements holds:

- For every graph  $H$ , if  $\chi(G \times H) \leq n$ , then  $\chi(H) \leq n$ .
- For every graph  $H$ , if  $G \times H \leq K_n$ , then  $H \leq K_n$ .
- For every graph  $H$ , if  $H \leq (K_n)^G$ , then  $H \leq K_n$ .
- $(K_n)^G \leq K_n$ .
- $\chi((K_n)^G) \leq n$ .

So we have the following equivalent formulation of Hedetniemi's conjecture, first observed by El-Zahar and Sauer [ES85, Conjecture 2]:

**Conjecture 2** (Hedetniemi v2). *Let  $G$  be a finite graph with  $\chi(G) > n$ . Then  $\chi((K_n)^G) \leq n$ .*

We will be dealing with this version of the conjecture, and thus we will be interested in the graph  $(K_n)^G$ . This is the graph with vertex set  $n^{V(G)}$ , where  $\alpha$  and  $\beta$  are adjacent iff for all  $(v, w) \in E(G)$ , we have  $\alpha(v) \neq \beta(w)$ . Let  $\bar{k}$  denote the constant function taking the value  $k$ .

A **suited colouring** of  $(K_n)^G$  is a proper  $n$ -colouring  $\Phi : (K_n)^G \rightarrow K_n$  such that  $\Phi(\bar{k}) = k$  for all  $k \in n$ . If  $\chi((K_n)^G) \leq n$ , then it is easily seen that there is a suited colouring of  $(K_n)^G$ , so we will work exclusively with suited colourings. The main convenience offered by suited colourings is the following fact, which we will use freely without further mention:

**Proposition 1.** *Let  $\Phi$  be a suited colouring of  $(K_n)^G$ . Then for every  $\alpha \in (K_n)^G$ , we have  $\Phi(\alpha) \in \text{im}(\alpha)$ .*

*Proof.* Since  $\Phi$  is suited, we have  $\Phi(\alpha) = \Phi(\overline{\Phi(\alpha)})$ , so since  $\Phi$  is a proper colouring,  $\alpha$  and  $\overline{\Phi(\alpha)}$  are not adjacent in  $(K_n)^G$ . In particular, their images must intersect, and thus  $\Phi(\alpha) \in \text{im}(\alpha)$ .  $\square$

Given graphs  $G$  and  $H$ , the **strong product**  $G \boxtimes H$  is the graph with vertex set  $V(G) \times V(H)$  such that  $(v, w)$  and  $(v', w')$  are adjacent iff one of the following holds:

- (1)  $(v, v') \in E(G)$  and  $(w, w') \in E(H)$ .
- (2)  $(v, v') \in E(G)$  and  $w = w'$ .
- (3)  $v = v'$  and  $(w, w') \in E(H)$ .

We can now state the main theorem:

**Theorem 2.** *Let  $G$  be a finite graph with girth  $> 5$  and radius  $> 2$ , and let  $n > 6m$ . If  $n$  is sufficiently large, then  $\chi((K_n)^{G \boxtimes K_m}) > n$ .*

To use this to refute Hedetniemi's conjecture, we will use the **fractional chromatic number**, which is defined as follows:

$$\chi_f(G) := \inf_m \frac{\chi(G \boxtimes K_m)}{m}$$

*Proof of Theorem 1.* Fix a finite graph  $G$  with girth  $> 5$  and  $\chi_f(G) > 7$  (for example, any graph with girth  $> 5$  and independence number  $< \frac{|G|}{7}$ , see [Die17, 11.2.2]).  $G$  has radius  $> 2$ , since otherwise  $G$  would be a tree, which has  $\chi_f(G) \leq 2$ . Then for any  $m$ , we have

$$\chi(G \boxtimes K_m) \geq \chi_f(G) \cdot m > 7m,$$

Let  $n = 7m$ . By Theorem 2, if  $n$  is sufficiently large, then  $\chi((K_n)^{G \boxtimes K_m}) > n$ , refuting Conjecture 2.  $\square$

### 3. CONDENSED PROOF

We present a version of the proof requiring minimal overhead.

We will write  $\mathbb{P}_k$  to denote the probability taken over  $k \in n$  uniformly distributed, and similarly for  $\mathbb{P}_\alpha$ , where  $\alpha$  ranges over  $n^G$ .

*Proof of Theorem 2.* For every  $\alpha \in n^G$ , write  $\bar{\alpha} \in (K_n)^{G \boxtimes K_m}$  for the function  $\bar{\alpha}(v, k) = \alpha(v)$ .

Suppose that  $\chi((K_n)^{G \boxtimes K_m}) \leq n$ , and fix a suited colouring  $\Psi$  of  $(K_n)^{G \boxtimes K_m}$ .

First we find for each  $v \in G$ , some  $\mu_v \in (K_n)^{G \boxtimes K_m}$  such that

- (1)  $|\text{im}(\mu_v)| = 3m + 1$ , and
- (2)  $\mu_v(w, i) = \Psi(\mu_v)$  iff  $d(w, v) > 2$ .

For  $k \in \{3m, \dots, n-1\}$ , define  $\mu_k \in (K_n)^{G \boxtimes K_m}$  via

$$\mu_k(w, i) = \begin{cases} i & w = v \\ i + m & d(w, v) = 1 \\ i + 2m & d(w, v) = 2 \\ k & \text{otherwise} \end{cases}$$

Since  $G$  has radius  $> 2$ , the set  $\{\mu_{3m}, \dots, \mu_{n-1}\}$  has size  $n - 3m$ , and since  $G$  has girth  $> 5$ , it is a clique in  $(K_n)^{G \boxtimes K_m}$ . Thus since  $n - 3m > 3m$ , there is some  $\mu_k$  such that  $\Psi(\mu_k) \notin \{0, \dots, 3m-1\}$ . But since  $\Psi(\mu_k) \in \text{im}(\mu_k) = \{0, \dots, 3m-1, k\}$ , we must have  $\Psi(\mu_k) = k$ . Set  $\mu_v = \mu_k$ .

Next, we claim that for every  $\alpha \in n^G$ , if  $\Psi(\bar{\alpha}) = \alpha(v) \notin \text{im}(\mu_v)$ , then there is some  $v' \neq v$  such that  $\alpha(v') \in \{\Psi(\mu_v), \alpha(v)\}$ . To see this, define  $\beta \in n^G$  as follows:

$$\beta(w) = \begin{cases} \Psi(\mu_v) & d(w, v) \leq 1 \\ \alpha(v) & \text{otherwise} \end{cases}$$

Since  $\alpha(v) \notin \text{im}(\mu_v)$ , we must have  $\bar{\beta}$  and  $\mu_v$  adjacent, and thus  $\Psi(\bar{\beta}) \neq \Psi(\mu_v)$ . Thus  $\Psi(\bar{\beta}) = \alpha(v) = \Psi(\bar{\alpha})$ . So  $\bar{\beta}$  and  $\bar{\alpha}$  are not adjacent, and thus there is some  $v' \neq v$  such that  $\alpha(v') \in \text{im}(\beta) = \{\Psi(\mu_v), \alpha(v)\}$ .

The claim, combined with the fact that  $|\text{im}(\mu_v)| = 3m + 1 < \frac{n}{2} + 1$ , gives the following inequality:

$$\begin{aligned} \left(\frac{1}{2} - \frac{1}{n}\right)^{|G|} &< \prod_v \mathbb{P}_k[k \notin \text{im}(\mu_v)] \\ &= \mathbb{P}_\alpha[\forall v [\alpha(v) \notin \text{im}(\mu_v)]] \\ &\leq \sum_v \mathbb{P}_\alpha[\Psi(\bar{\alpha}) = \alpha(v) \notin \text{im}(\mu_v)] \\ &\leq \sum_v \mathbb{P}_\alpha[\exists v' \neq v [\alpha(v') \in \{\Psi(\mu_v), \alpha(v)\}]] \\ &\leq \sum_v \sum_{v' \neq v} \mathbb{P}_\alpha[\alpha(v') \in \{\Psi(\mu_v), \alpha(v)\}] \\ &\leq \sum_v \sum_{v' \neq v} \frac{2}{n} \\ &= \frac{2|G|(|G| - 1)}{n} \end{aligned}$$

This only holds for finitely many  $n$ , so we are done.  $\square$

## APPENDIX A. UNCONDENSED PROOF

We will need the following important definition:

**Definition 1** (Shitov). Fix a suited colouring  $\Phi$  of  $(K_n)^G$ . Then for a vertex  $v \in G$ , a colour  $k \in n$  is  *$v$ -robust* if for every  $\alpha \in (K_n)^G$  with  $\Phi(\alpha) = k$ , there is a vertex  $w \in G$  with  $d(w, v) \leq 1$  and  $\alpha(w) = k$ .

The proof strategy goes as follows. We will show that for any suited colouring, there is always a vertex with many robust colours. However, Hedetniemi's conjecture will provide us with too many suited colourings, in particular, one which does not have a vertex with many robust colours.

We first prove the existence of a vertex with many robust colours:

**Proposition 2.** *Let  $G$  be a finite graph and fix a suited colouring  $\Phi$  of  $(K_n)^G$ . Then there is some  $v \in G$  such that*

$$\mathbb{P}_k[k \text{ is not } v\text{-robust}] < \sqrt[|G|]{\frac{|G|^3}{n}}$$

*Proof.* For every  $v$  and  $k$ , define  $\beta_{v,k} \in (K_n)^G$  as follows: if  $k$  is  $v$ -robust, pick  $\beta_{v,k}$  arbitrarily; otherwise, pick  $\beta_{v,k}$  witnessing the non-robustness, i.e. such that  $\Phi(\beta_{v,k}) = k$  and for every  $v'$  with  $d(v', v) \leq 1$ , we have  $\beta_{v,k}(v') \neq k$ .

We claim that if  $\Phi(\alpha) = \alpha(v)$  and  $\alpha(v)$  is not  $v$ -robust, then  $\exists v' \neq v$  such that  $\alpha(v') \in \text{im}(\beta_{v,\alpha(v)})$ . If not, then  $\alpha$  and  $\beta$  would be adjacent in  $(K_n)^G$ , contradicting the fact that  $\Phi(\alpha) = \Phi(\beta)$ .

This gives the following bound, from which the proposition follows immediately:

$$\begin{aligned} & \prod_v \mathbb{P}_k[k \text{ is not } v\text{-robust}] \\ &= \mathbb{P}_\alpha[\forall v [\alpha(v) \text{ is not } v\text{-robust}]] \\ &\leq \sum_v \mathbb{P}_\alpha[\Phi(\alpha) = \alpha(v) \text{ and } \alpha(v) \text{ is not } v\text{-robust}] \\ &\leq \sum_v \sum_{v' \neq v} \mathbb{P}_\alpha[\alpha(v') \in \text{im } \beta_{v,\alpha(v)}] \\ &\leq \sum_v \sum_{v' \neq v} \frac{|G|}{n} \\ &< \frac{|G|^3}{n} \end{aligned}$$

□

Given a graph  $G$ , the **reflexive closure**  $G^\circ$ , is the graph obtained from  $G$  by adding every loop.

We will now see the consequences of having too many suited colourings:

**Proposition 3.** *Let  $G$  be a graph with girth  $> 5$  and radius  $> 2$ , and let  $n > 6m$ . Suppose that  $\chi((K_n)^{G \boxtimes K_m}) \leq n$ . Then there is a suited colouring of  $(K_n)^{G^\circ}$  such that for every  $v \in G^\circ$ ,*

$$\mathbb{P}_k[k \text{ is } v\text{-robust}] < \frac{1}{2} + \frac{1}{n}$$

*Proof.* Note that there is a natural map  $G \boxtimes K_m \rightarrow G^\circ$ , and this induces a natural map  $(K_n)^{G^\circ} \rightarrow (K_n)^{G \boxtimes K_m}$ . For every  $\alpha \in (K_n)^{G^\circ}$ , let  $\bar{\alpha} \in (K_n)^{G \boxtimes K_m}$  be the image of  $\alpha$  under this map. Let  $\Psi$  be a suited colouring of  $(K_n)^{G \boxtimes K_m}$ , and let  $\Phi$  be the suited colouring of  $(K_n)^{G^\circ}$  obtained by composing  $\Psi$  with the map  $(K_n)^{G^\circ} \rightarrow (K_n)^{G \boxtimes K_m}$ .

Now fix  $v \in G^\circ$ . For  $k \in \{3m, \dots, n-1\}$ , define  $\mu_k \in (K_n)^{G \boxtimes K_m}$  via

$$\mu_k(w, i) = \begin{cases} i & w = v \\ i + m & d(w, v) = 1 \\ i + 2m & d(w, v) = 2 \\ k & \text{otherwise} \end{cases}$$

Since  $G$  has radius  $> 2$ , the set  $\{\mu_{3m}, \dots, \mu_{n-1}\}$  has size  $n - 3m$ , and since  $G$  has girth  $> 5$ , it is a clique in  $(K_n)^{G \boxtimes K_m}$ . Thus since  $n - 3m > 3m$ , there is some  $\mu_k$  such that  $\Psi(\mu_k) \notin \{0, \dots, 3m-1\}$ . But since  $\Psi(\mu_k) \in \text{im}(\mu_k) = \{0, \dots, 3m-1, k\}$ , we must have  $\Psi(\mu_k) = k$ .

It suffices to show that every  $v$ -robust colour is contained in  $\text{im}(\mu_k)$ , since  $|\text{im}(\mu_k)| = 3m+1 < \frac{n}{2} + 1$ . To this end, let  $l$  be a  $v$ -robust colour. Define  $\beta \in (K_n)^{G^\circ}$  as follows:

$$\beta(w) = \begin{cases} k & d(w, v) \leq 1 \\ l & \text{otherwise} \end{cases}$$

We must have  $\Phi(\beta) = k$ , since if  $\Phi(\beta) = l$ , then by robustness, we would have  $\beta(w) = k$  for some  $w$  with  $d(v, w) \leq 1$ , and thus  $l = k$ . Thus  $\Psi(\beta) = \Phi(\beta) = k$ . Since  $\Psi(\mu_k) = k$  and  $\Psi$  is a proper colouring,  $\mu_k$  and  $\beta$  are not adjacent, and thus we must have  $l \in \text{im}(\mu_k)$ .  $\square$

*Proof of Theorem 2.* Suppose that  $\chi((K_n)^{G \boxtimes K_m}) \leq n$ . By Proposition 3, there is a suited colouring  $\Phi$  of  $(K_n)^{G^\circ}$  such that for every  $v \in G^\circ$ , we have

$$\mathbb{P}_k[k \text{ is } v\text{-robust}] < \frac{1}{2} + \frac{1}{n}.$$

But by [Proposition 2](#), there is some  $v \in G$  such that

$$\mathbb{P}_k[k \text{ is not } v\text{-robust}] < \sqrt[|G|]{\frac{|G|^3}{n}}.$$

There are only finitely many  $n$  such that both of these hold (since their sum is  $\frac{1}{2} + o(1)$ ).  $\square$

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