

Lattices from equiangular tight frames with applications to lattice sparse recovery

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The compressed sensing problem

1. Signal of interest $f \in \mathbb{C}^d (= \mathbb{C}^{N \times N})$
2. Measurement operator $\mathcal{A} : \mathbb{C}^d \rightarrow \mathbb{C}^m$ ($m \ll d$)
3. Measurements $y = \mathcal{A}f + \xi$

$$\begin{bmatrix} y \end{bmatrix} = \begin{bmatrix} \mathcal{A} \end{bmatrix} \begin{bmatrix} f \end{bmatrix} + \begin{bmatrix} \xi \end{bmatrix}$$

4. **Problem:** Reconstruct signal f from measurements y

Sparsity

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Assume f is *sparse*:

- In the coordinate basis: $\|f\|_0 \stackrel{\text{def}}{=} |\text{supp}(f)| \leq s \ll d$
- In orthonormal basis: $f = Bx$ where $\|x\|_0 \leq s \ll d$

In practice, we encounter *compressible* signals.

- ★ f_s is the best s -sparse approximation to f

Many applications

- Radar, Error Correction
- Computational Biology, Geophysical Data Analysis
- Data Mining, classification
- Neuroscience
- Imaging
- Sparse channel estimation, sparse initial state estimation
- Topology identification of interconnected systems
- ...

Reconstruction approaches

- ★ ℓ_1 -minimization [Candès-Romberg-Tao]

Let A satisfy the *Restricted Isometry Property* and set:

$$\hat{f} = \underset{g}{\operatorname{argmin}} \|g\|_1 \quad \text{such that} \quad \|Af - y\|_2 \leq \varepsilon,$$

where $\|\xi\|_2 \leq \varepsilon$. Then we can stably recover the signal f :

$$\|f - \hat{f}\|_2 \lesssim \varepsilon + \frac{\|x - x_s\|_1}{\sqrt{s}}.$$

This error bound is optimal.

- ★ Other methods (iterative, greedy) too (OMP, ROMP, StOMP, CoSaMP, IHT, ...)

Restricted Isometry Property

- \mathcal{A} satisfies the Restricted Isometry Property (RIP) when there is $\delta < c$ such that

$$(1 - \delta)\|f\|_2 \leq \|\mathcal{A}f\|_2 \leq (1 + \delta)\|f\|_2 \quad \text{whenever } \|f\|_0 \leq s.$$

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★ Related to dimension reduction and the Johnson-Lindenstrauss Lemma (dimension reduction with preserved geometry).

Sparsity plus other structures?

What if signal is also *lattice*-valued?

- Wireless communications
- Radar (massive MIMO) [Rossi et.al.]
- Wideband spectrum sensing [Axell et.al.]
- Error correcting codes [Candès et.al.]
- ...

Lattices

What is a lattice?

Lattices

What is a lattice?

A **lattice** $\Lambda \subset \mathbb{R}^n$ of rank $1 \leq k \leq n$ is a free \mathbb{Z} -module of rank k , which is the same as a discrete co-compact subgroup of $V := \text{span}_{\mathbb{R}} \Lambda$. If $k = n$, i.e. $V = \mathbb{R}^n$, we say that Λ is a lattice of **full rank** in \mathbb{R}^n . Hence

$$\Lambda = \text{span}_{\mathbb{Z}}\{\mathbf{a}_1, \dots, \mathbf{a}_k\} = A\mathbb{Z}^k,$$

where $\mathbf{a}_1, \dots, \mathbf{a}_k \in \mathbb{R}^n$ are \mathbb{R} -linearly independent **basis** vectors for Λ and $A = (\mathbf{a}_1 \ \dots \ \mathbf{a}_k)$ is the corresponding $n \times k$ basis matrix.

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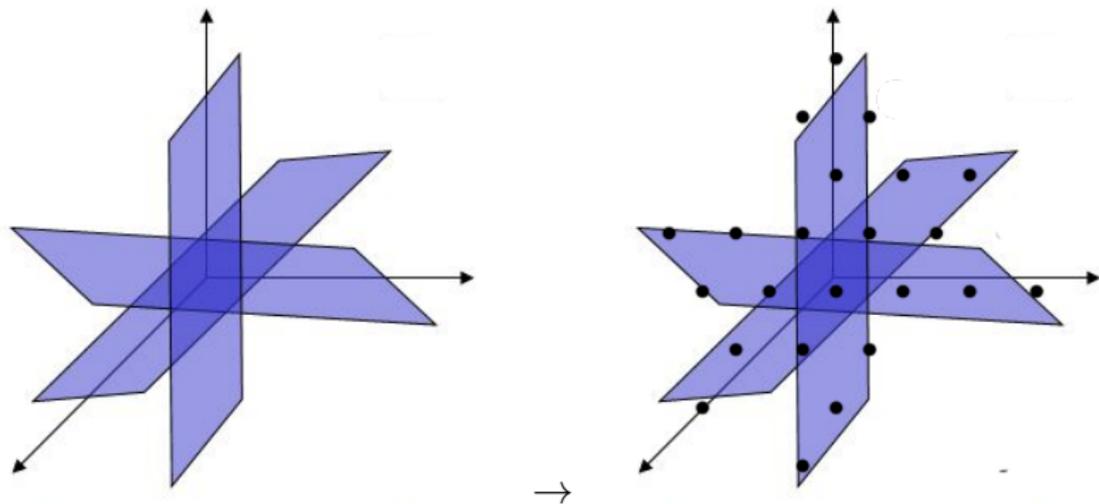
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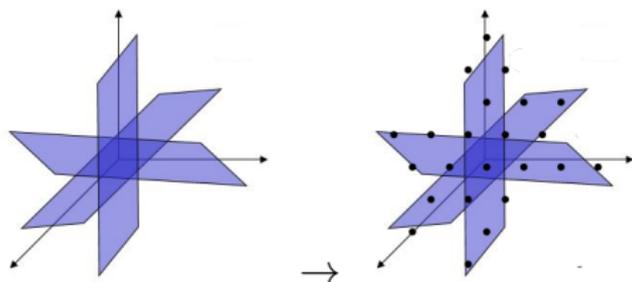
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★ A *sparse lattice-valued signal* is $\mathbf{v} \in \Lambda$ with $\|\mathbf{v}\|_0 \leq s$. Alternatively can consider $\mathbf{v} = A\mathbf{w}$ where $\mathbf{w} \in \mathbb{Z}^k$ and $\|\mathbf{w}\|_0 \leq s$.

Lattices

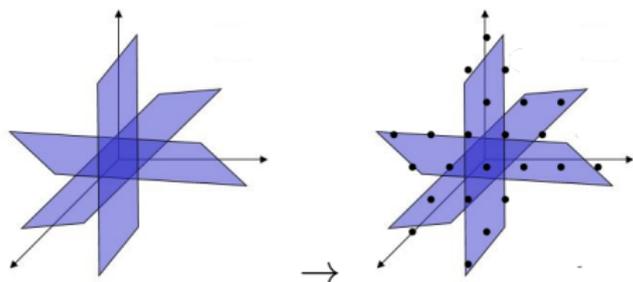


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- Question: when is lattice knowledge helpful??

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- Suppose the signal x is 1-sparse. Need $\approx \log(d)$ RIP measurements.

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- Or suppose instead that $x \in \Lambda = \mathbb{Z}^d$. ??
- The point: sometimes lattice info can give a huge savings. Sometimes maybe not?

Some results

- Dense ± 1 signals [Mangasarian-Recht '11] :
$$\min \|x\|_{\infty} \quad \text{s.t.} \quad \mathcal{A}x = y$$

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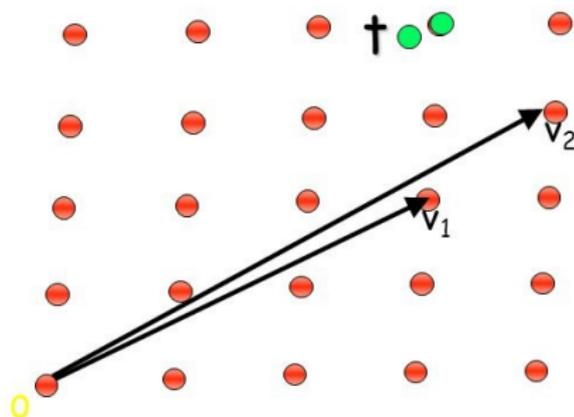
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- Sparse lattice signals [Flinth-Kutyniok '16] : OMP with initialization step (PROMP)

Some results [Sphere decoders]

- ★ The closest point problem:

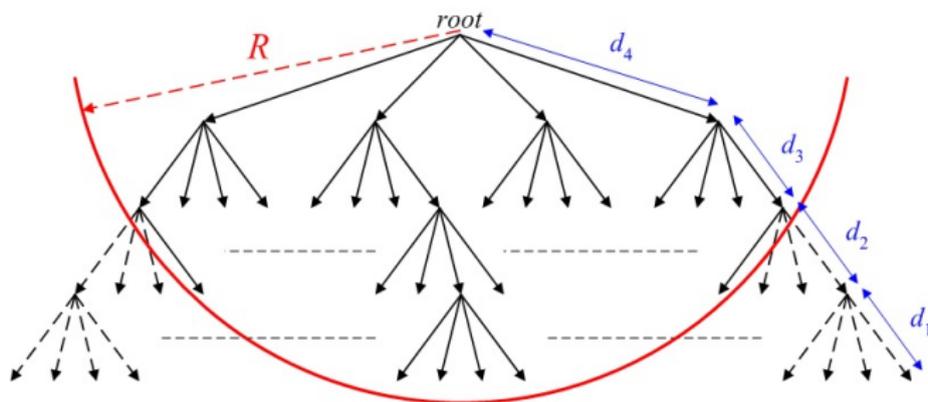
Closest Vector Problem (CVP)



- ★ Find point in lattice closest to a given vector in some metric (e.g. $\|x - y\|_2$).

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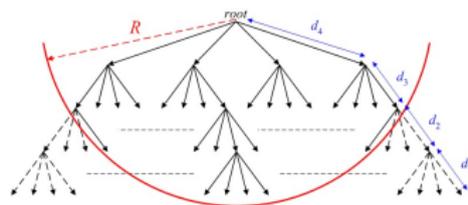
- ★ Sphere decoder:



- ★ Using some ordering of the lattice (recursively), prune the search tree using spheres of specified radius.

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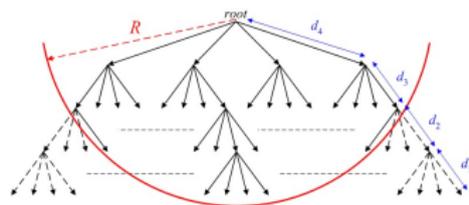
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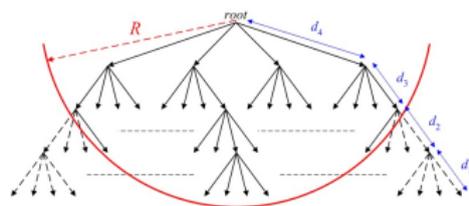
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- Use sphere decoder method with metric $\|y - Ax\|_2 + \lambda\|x\|_0$
- Lattice pruning/ordering no longer clear in this metric
- Lack of rigorous theory

Some results [Flinth-Kutyniok '16]

- ★ PROMP:
 - Run least squares $\hat{x} = \operatorname{argmin}_x \|\mathcal{A}x - y\|_2$
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- ★ Some theory about accuracy of initialization

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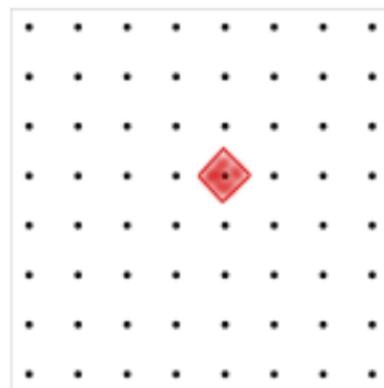
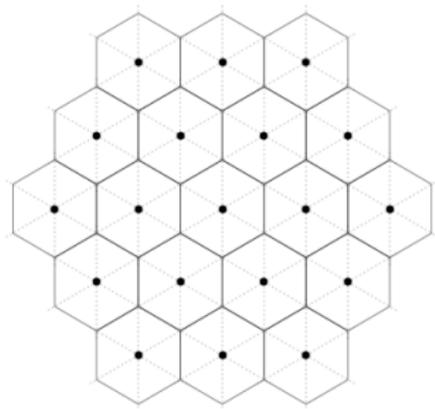
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★ Same is true for lattices whose Voronoi region Ω satisfies $A^{-1}\Omega \subset (-1, 1)^k$.

Voronoi: $\Omega \stackrel{\text{def}}{=} \{v : \forall z \in A\mathbb{Z}^k, \|v\|_2 \leq \|v - z\|_2\}$. (e.g. diamond)



Lattices: minimal vectors

Minimal norm of a lattice Λ is

$$|\Lambda| = \min \{ \|\mathbf{x}\| : \mathbf{x} \in \Lambda \setminus \{\mathbf{0}\} \},$$

where $\|\cdot\|$ is Euclidean norm. The set of **minimal vectors** of Λ is

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- If $\text{rk } \Lambda > 4$, a strictly stronger condition is that Λ is **generated by minimal vectors**, i.e. $\Lambda = \text{span}_{\mathbb{Z}} S(\Lambda)$.
- It has been shown by Conway & Sloane (1995) and Martinet & Schürmann (2011) that there are lattices of rank ≥ 10 generated by minimal vectors which do not contain a **basis of minimal vectors**.

Lattices: eutaxy and perfection

Let $k = \text{rk } \Lambda$ and

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$$\|\mathbf{v}\|^2 = \sum_{i=1}^m c_i \langle \mathbf{v}, \mathbf{x}_i \rangle^2$$

for every vector $\mathbf{v} \in \text{span}_{\mathbb{R}} \Lambda$, where $\langle \cdot, \cdot \rangle$ is the usual inner product.

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This lattice is called **perfect** if the set of symmetric matrices

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spans the space of $k \times k$ symmetric matrices.

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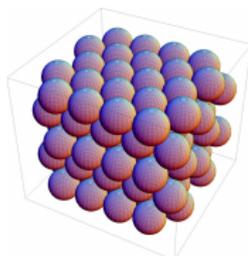
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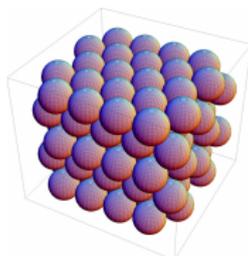
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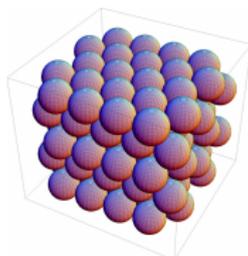
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★ If a lattice is strongly eutactic, but not perfect, then it is a local minimum of the packing density function. Extremal lattices are local maxima.

Equiangular frames

Another interesting construction of lattices comes from *frames*. A collection of $n \geq k$ unit vectors $\mathbf{f}_1, \dots, \mathbf{f}_n \in \mathbb{R}^k$ is called an (real) (k, n) -**equiangular tight frame** (ETF) if it spans \mathbb{R}^k and

1. $|\langle \mathbf{f}_i, \mathbf{f}_j \rangle| = c$ for all $1 \leq i \neq j \leq n$, for some constant $c \in [0, 1]$,
2. $\sum_{i=1}^n \langle \mathbf{f}_i, \mathbf{x} \rangle^2 = \gamma \|\mathbf{x}\|^2$ for each $\mathbf{x} \in \mathbb{R}^k$, for some absolute constant $\gamma \in \mathbb{R}$.

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If this is the case, it is known that

$$k \leq n \leq \frac{k(k+1)}{2}, \quad c = \sqrt{\frac{n-k}{k(n-1)}}, \quad \gamma = \sqrt{\frac{n}{k}}.$$

Mercedes

Here is a (2, 3)-ETF $\mathcal{F} := \left\{ \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \begin{pmatrix} -1/2 \\ -\sqrt{3}/2 \end{pmatrix}, \begin{pmatrix} 1/2 \\ -\sqrt{3}/2 \end{pmatrix} \right\}$:



Mercedes

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Notice that $\pm\mathcal{F} = S(\Lambda_h)$, the set of minimal vectors of the hexagonal lattice $\Lambda_h = \begin{pmatrix} 1 & 1/2 \\ 0 & \sqrt{3}/2 \end{pmatrix} \mathbb{Z}^2$.

Lattices: questions

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- If so, does it have a basis of minimal vectors?
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Consequences:

- If the span is a lattice, the frame viewed as a sensing matrix yields an image that is a discrete set.
- If the frame atoms are minimal vectors, we can guarantee separation between sample vectors in its image.
- Johnson-Lindenstrauss may then be used for reconstruction guarantees?
- When is reconstruction impossible?

Lattice construction

Let $\mathcal{F} = \{\mathbf{f}_1, \dots, \mathbf{f}_n\} \subset \mathbb{R}^k$ be a (k, n) -ETF, and define

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Proposition 2 (Böttcher, Fukshansky, Garcia, Maharaj, N- '16)

If $\Lambda(\mathcal{F})$ is a lattice and

$$S(\Lambda(\mathcal{F})) = \{\pm \mathbf{f}_1, \dots, \pm \mathbf{f}_n\},$$

then $\Lambda(\mathcal{F})$ is strongly eutactic.

Main results on ETF lattices

Theorem 3 (Böttcher, Fukshansky, Garcia, Maharaj, N- '16)

1. *For every $k \geq 2$, there are $(k, k + 1)$ -ETFs \mathcal{F} such that $\Lambda(\mathcal{F})$ is a full-rank lattice. This lattice has a basis of minimal vectors, is non-perfect and strongly eutactic.*

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3. *There are $(3, 6)$, $(7, 14)$, and $(9, 18)$ -ETFs \mathcal{F} for which $\Lambda(\mathcal{F})$ is not a lattice.*

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2. *There are infinitely many k for which there exist $(k, 2k)$ -ETFs \mathcal{F} such that $\Lambda(\mathcal{F})$ is a full-rank lattice, e.g. $(5, 10)$, $(13, 26)$.*
3. *There are $(3, 6)$, $(7, 14)$, and $(9, 18)$ -ETFs \mathcal{F} for which $\Lambda(\mathcal{F})$ is not a lattice.*
4. *There is a $(7, 28)$ -ETF \mathcal{F} for which $\Lambda(\mathcal{F})$ is a full-rank lattice that has a basis of minimal vectors, is a perfect strongly eutactic lattice, and hence extreme.*

Main results on ETF lattices

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5. *There is a $(6, 16)$ -ETF \mathcal{F} for which $\Lambda(\mathcal{F})$ is a full-rank lattice that has a basis of minimal vectors.*

Remarks

- There are often multiple ETFs with the same parameters (k, n) . For instance, we exhibit two lattices from $(5, 10)$ -ETFs, three lattices from $(13, 26)$ -ETFs, and ten lattices from $(25, 50)$ -ETFs. We also compute determinants of all our examples.

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- Perfection of the lattice from $(7, 28)$ -ETF was previously (2015) established by Roland Bacher, however he constructed this lattice differently and then remarked that its minimal vectors comprise a set of equiangular lines.
- Minimal vectors of ETF lattices often are precisely \pm frame vectors (this is the case with all our examples). In this case, the set of corresponding symmetric matrices has *at most* $k(k+1)/2$ matrices, which is the *least* possible number required to span all symmetric matrices. Hence ETF lattices are unlikely to be perfect (and hence extremal) – the $(7, 28)$ case is likely an exception.

Future directions

- Further study *geometric* properties of ETFs to decipher when they create a lattice.
- Given a lattice ETF whose atoms are minimal vectors, *how* can we reconstruct lattice signals?
- How can we incorporate sparsity? → need Johnson-Lindenstrauss
- Computationally efficient reconstructions that beat classical CS methods?

References

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Thank you!

Now "lattice" take any questions...