

# Noisy Signal Recovery via Iterative Reweighted L1-Minimization

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# Setup

- ① Suppose  $x$  is an unknown signal in  $\mathbb{R}^d$ .
- ② Design measurement matrix  $\Phi : \mathbb{R}^d \rightarrow \mathbb{R}^m$ .
- ③ Collect noisy measurements  $u = \Phi x + e$ .

$$\begin{bmatrix} u \end{bmatrix} = \begin{bmatrix} \Phi \end{bmatrix} \begin{bmatrix} x \end{bmatrix} + \begin{bmatrix} e \end{bmatrix}$$

- ④ **Problem:** Reconstruct signal  $x$  from measurements  $u$
- ⑤ Wait, isn't this impossible?
  - Assume  $x$  is  $s$ -sparse:  $\|x\|_0 \stackrel{\text{def}}{=} |\text{supp}(x)| \leq s \ll d$ .

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# How can we reconstruct?

Obvious way:

Suppose the matrix  $\Phi$  is one-to-one on the set of sparse vectors and  $e = 0$ . Set

$$\hat{x} = \operatorname{argmin} \|z\|_0 \quad \text{such that} \quad \Phi z = u.$$

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Bad news:

This would require a search through  $\binom{d}{s}$  subspaces! **Not numerically feasible.**

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# How else can we reconstruct?

## Geometric Idea

Minimizing the  $\ell_0$ -ball is too hard, so let's try a different one.

Our favorites...

- Least Squares
- L1-Minimization (using Linear Programming)

Which one?

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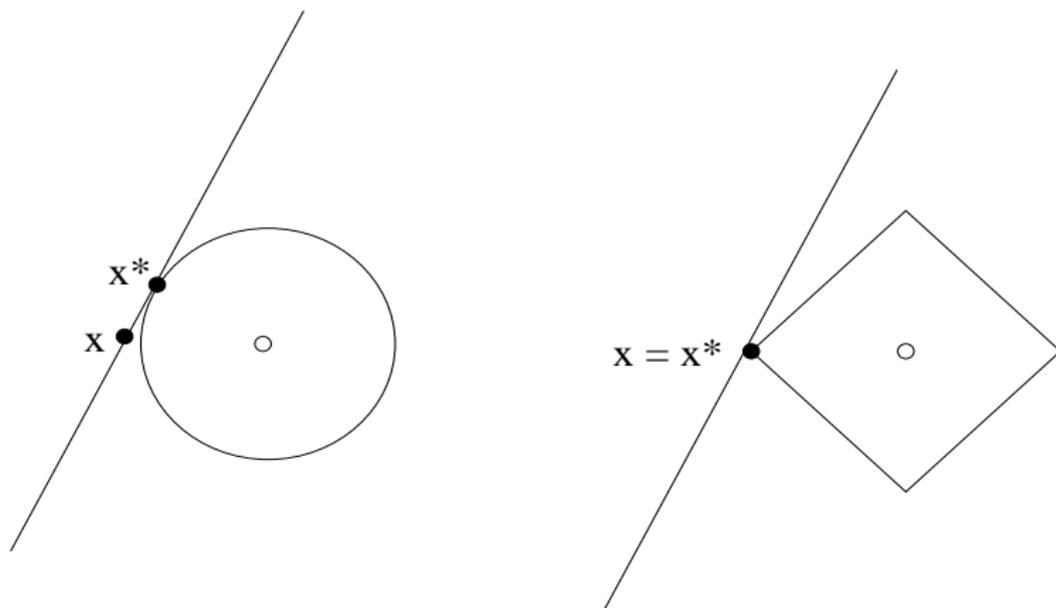


Figure: Minimizing the  $l_2$  versus the  $l_1$  balls.

# What do we assume about $\Phi$ ?

## Restricted Isometry Property (RIP)

- The  $s^{\text{th}}$  **restricted isometry constant** of  $\Phi$  is the smallest  $\delta_s$  such that

$$(1 - \delta_s)\|x\|_2 \leq \|\Phi x\|_2 \leq (1 + \delta_s)\|x\|_2 \quad \text{whenever } \|x\|_0 \leq s.$$

- For Gaussian or Bernoulli measurement matrices, with high probability

$$\delta_s \leq c < 1 \quad \text{when } m \gtrsim s \log d.$$

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# Proven Results

## L1-Minimization [Candès-Tao]

Assume that the measurement matrix  $\Phi$  satisfies the RIP with  $\delta_{2s} < \sqrt{2} - 1$ . Then every  $s$ -sparse vector  $x$  can be exactly recovered from its measurements  $u = \Phi x$  as a unique solution to the linear optimization problem:

$$\hat{x} = \operatorname{argmin} \|z\|_1 \quad \text{such that} \quad \Phi z = u.$$

# Numerical Results

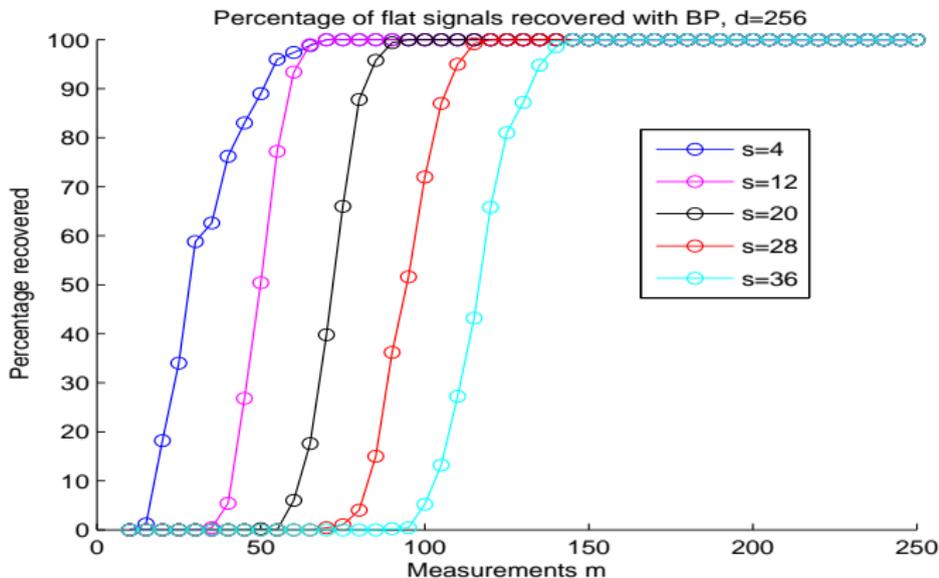


Figure: The percentage of sparse flat signals exactly recovered by Basis Pursuit as a function of the number of measurements  $m$  in dimension  $d = 256$  for various levels of sparsity  $s$ .

# What about noise?

## Noisy Formulation

For a non-sparse vector  $x$  with noisy measurements  $u = \Phi x + e$ ,

$$\hat{x} = \operatorname{argmin} \|z\|_1 \quad \text{such that} \quad \|\Phi z - u\|_2 \leq \varepsilon. \quad (1)$$

## L1-Minimization [Candès-Romberg-Tao]

Let  $\Phi$  be a measurement matrix satisfying the RIP with  $\delta_{2s} < \sqrt{2} - 1$ . Then for any *arbitrary* signal and corrupted measurements  $u = \Phi x + e$  with  $\|e\|_2 \leq \varepsilon$ , the solution  $\hat{x}$  to (1) satisfies

$$\|\hat{x} - x\|_2 \leq C_s \cdot \varepsilon + C'_s \cdot \frac{\|x - x_s\|_1}{\sqrt{s}}.$$

Note: As  $\delta_{2s} \rightarrow \sqrt{2} - 1$ ,  $C_s, C'_s \rightarrow \infty!!$

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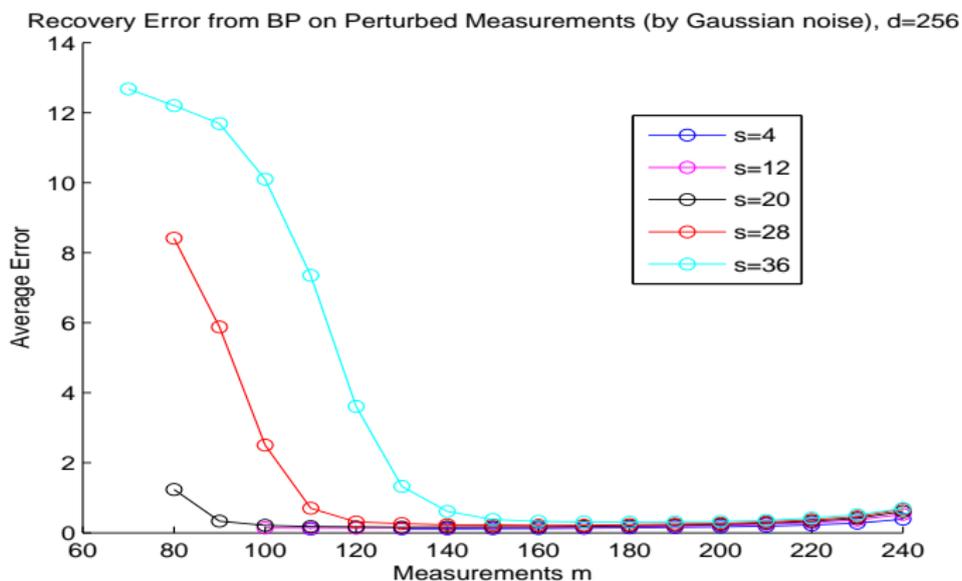


Figure: The recovery error of L1-Minimization under perturbed measurements ( $\|e\|_2 = 0.5$ ) as a function of the number of measurements  $m$  in dimension  $d = 256$  for various levels of sparsity  $s$ .

# What if we are close?

- Suppose we recover  $\hat{x} \approx x$
- Most likely, this means  $\hat{x}_i \approx x_i$
- In particular,  $\hat{x}_i$  is small/large when  $x_i$  is small/large

## Weighted L1

$$\hat{x}^{(2)} = \underset{z}{\operatorname{argmin}} \sum_{i=1}^d \left| \frac{z_i}{\hat{x}_i} \right| \quad \text{such that} \quad \|\Phi z - u\|_2 \leq \varepsilon$$

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$$\hat{x}^{(2)} = \operatorname{argmin}_z \sum_{i=1}^d \left| \frac{z_i}{\hat{x}_i + a} \right| \quad \text{such that} \quad \|\Phi z - u\|_2 \leq \varepsilon$$

# Weighted Geometry

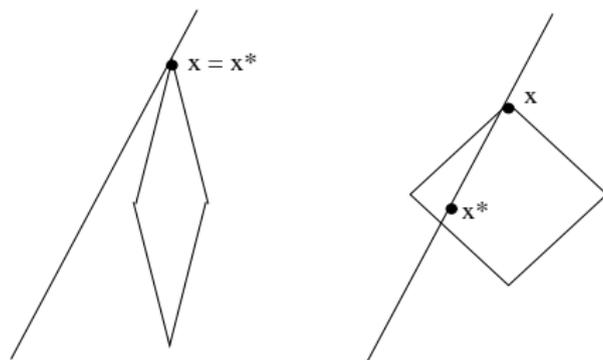


Figure: The geometry of the weighted  $\ell_1$ -ball.

- Noise-free case: In cases where  $\hat{x} \neq x$ , we should have that  $\hat{x}^{(2)}$  is closer to  $x$ , or even *equal*.
- Noisy case: This implies  $\hat{x}^{(2)}$  should be closer to  $x$  than  $\hat{x}$  was.
- Can we repeat this again?

# Reweighted $\ell_1$ -minimization (RWL1)

INPUT: Measurement vector  $u \in \mathbb{R}^m$ , stability parameter  $a$

OUTPUT: Reconstructed vector  $\hat{x}$

Initialize Set the weights  $w_i = 1$  for  $i = 1 \dots d$ .

Approximate Solve the reweighted  $\ell_1$ -minimization problem:

$$\hat{x} = \operatorname{argmin}_{\hat{x} \in \mathbb{R}^d} \sum_{i=1}^d w_i \hat{x}_i \text{ subject to } \|\Phi \hat{x} - u\|_2 \leq \varepsilon.$$

Update Reset the weights:

$$w_i = \frac{1}{|\hat{x}_i| + a}.$$

# Numerical Results

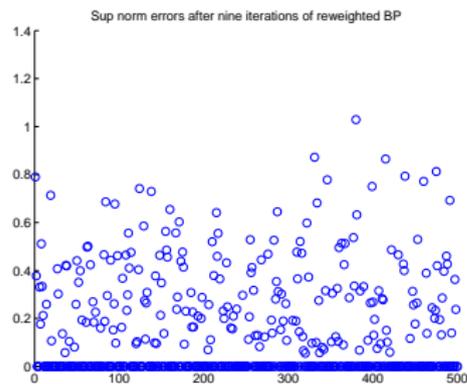
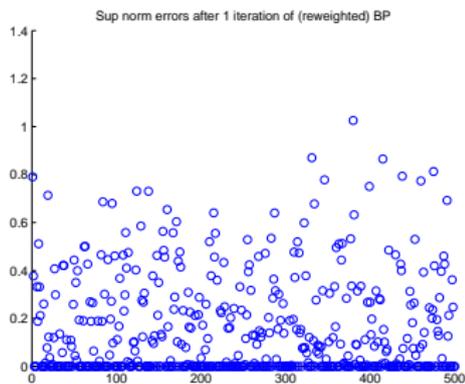


Figure:  $l_\infty$ -norm error for reweighted L1 in the noise-free case



# Numerical Results with noise

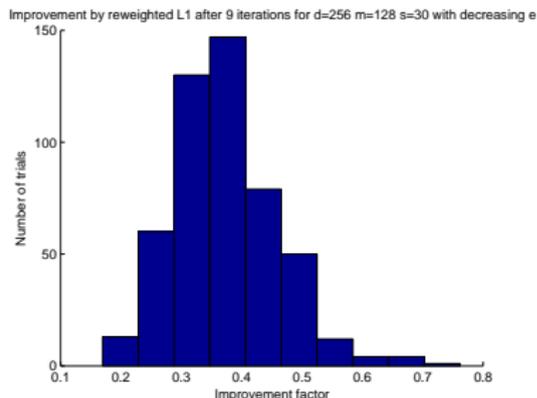
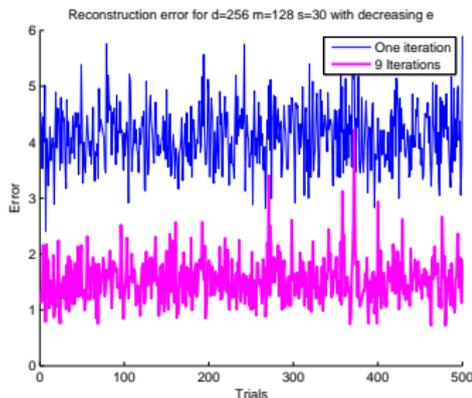


Figure: Improvements in the  $\ell_2$  reconstruction error using reweighted  $\ell_1$ -minimization versus standard  $\ell_1$ -minimization for sparse Gaussian signals.

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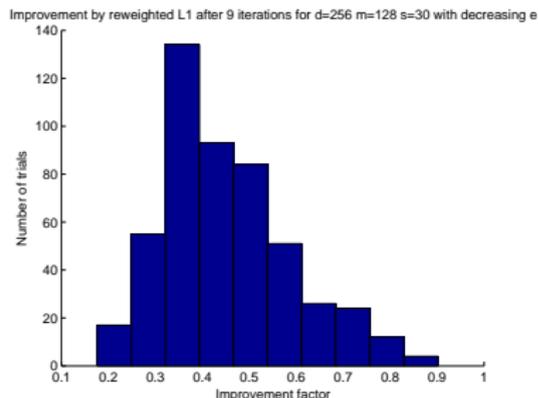
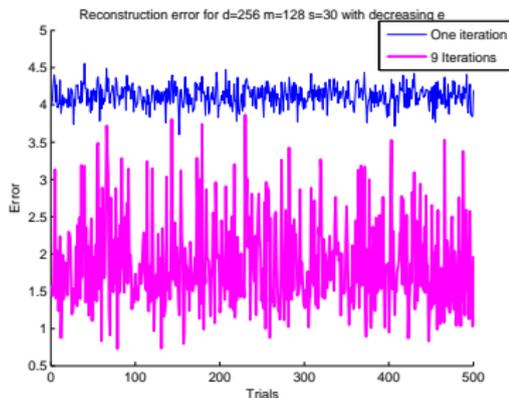


Figure: Improvements in the  $\ell_2$  reconstruction error using reweighted  $\ell_1$ -minimization versus standard  $\ell_1$ -minimization for sparse Bernoulli signals.

# Observations

- The noiseless case suggests that an  $\ell_\infty$ -norm bound may be required for RWL1 to succeed.
- In the noisy case it is clear that we cannot recover signal coordinates that are below some threshold.
- If each iteration of RWL1 improves the error, perhaps we should take  $a \rightarrow 0$ . (Recall  $w_i = \frac{1}{|\hat{x}_i| + a}$ ).

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# Main Results

## RWL1 - Sparse case [N]

Assume  $\Phi$  satisfies the RIP with  $\delta_{2s} \leq \delta$  where  $\delta < \sqrt{2} - 1$ . Let  $x$  be an  $s$ -sparse vector with noisy measurements  $u = \Phi x + e$  where  $\|e\|_2 \leq \varepsilon$ . Assume the smallest nonzero coordinate  $\mu$  of  $x$  satisfies  $\mu \geq \frac{4\alpha\varepsilon}{1-\rho}$ . Then the limiting approximation from reweighted  $\ell_1$ -minimization satisfies

$$\|x - \hat{x}\|_2 \leq C'' \varepsilon,$$

where  $C'' = \frac{2\alpha}{1+\rho}$ ,  $\rho = \frac{\sqrt{2}\delta}{1-\delta}$  and  $\alpha = \frac{2\sqrt{1+\delta}}{1-\delta}$ .

# Remarks

- Without noise, this result coincides with previous results on L1.
- The key improvement: As  $\delta \rightarrow \sqrt{2} - 1$ ,  $C''$  remains bounded.
- The error bound is the *limiting* bound, but a recursive relation in the proof gives exact improvements per iteration. We show in practice it is attained quite quickly.
- For signals whose smallest non-zero coefficient  $\mu$  does not satisfy the condition of the theorem, we may apply the theorem to those coefficients that do satisfy this requirement, and treat the others as noise...

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# Extension

## RWL1 - non-sparse extension [N]

Assume  $\Phi$  satisfies the RIP with  $\delta_{2s} \leq \sqrt{2} - 1$ . Let  $x$  be an arbitrary vector with noisy measurements  $u = \Phi x + e$  where  $\|e\|_2 \leq \varepsilon$ . Assume the smallest nonzero coordinate  $\mu$  of  $x_s$  satisfies  $\mu \geq \frac{4\alpha\varepsilon_0}{1-\rho}$ , where  $\varepsilon_0 = 1.2(\|x - x_s\|_2 + \frac{1}{\sqrt{s}}\|x - x_s\|_1) + \varepsilon$ . Then the limiting approximation from reweighted  $\ell_1$ -minimization satisfies

$$\|x - \hat{x}\|_2 \leq \frac{4.1\alpha}{1+\rho} \left( \frac{\|x - x_{s/2}\|_1}{\sqrt{s}} + \varepsilon \right),$$

and

$$\|x - \hat{x}\|_2 \leq \frac{2.4\alpha}{1+\rho} \left( \|x - x_s\|_2 + \frac{\|x - x_s\|_1}{\sqrt{s}} + \varepsilon \right),$$

where  $\rho$  and  $\alpha$  are as before.

# Theoretical Results

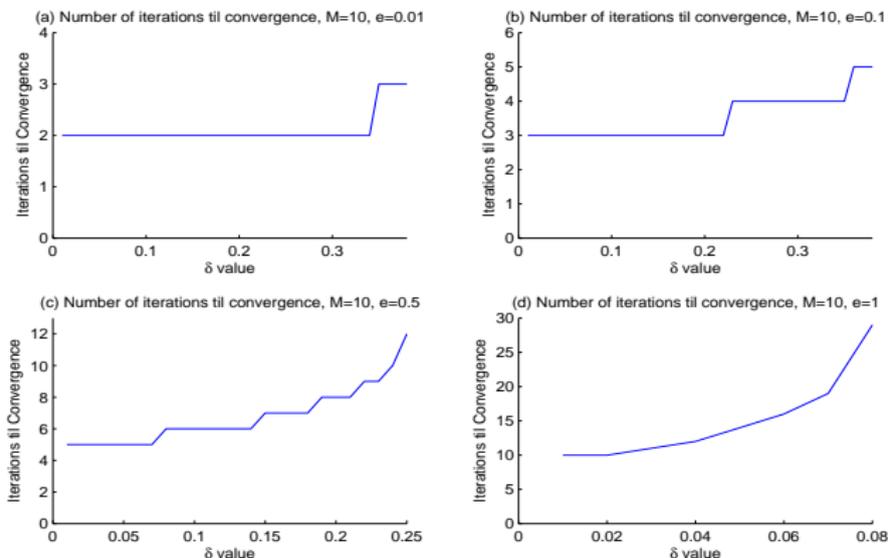


Figure: Number of iterations required for theoretical error bounds to reach limiting theoretical error when (a)  $\mu = 10$ ,  $\varepsilon = 0.01$ , (b)  $\mu = 10$ ,  $\varepsilon = 0.1$ , (c)  $\mu = 10$ ,  $\varepsilon = 0.5$ , (d)  $\mu = 10$ ,  $\varepsilon = 1.0$ .

# Recent work

- Wipf-Nagarajan elaborate on convergence and show connections to reweighted  $\ell_2$ -minimization.
- Wipf-Nagarajan also show that a non-separable variant has desirable properties.
- Xu-Khajehnejad-Avestimehr-Hassibi provide a theoretical foundation for the analysis of RWL1 and show that for a nontrivial class of signals, a variant of RWL1 indeed can improve upon L1 in the noiseless case.

Thank you

# For more information

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## References:

- Candes, Wakin, Boyd, “Enhancing sparsity by reweighted  $\ell_1$  minimization”, *J. Fourier Anal. Appl.*, **14** 877-905.
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- Wipf and Nagarajan, “Solving Sparse Linear Inverse Problems: Analysis of Reweighted  $\ell_1$  and  $\ell_2$  Methods,” *J. of Selected Topics in Signal Processing, Special Issue on Compressive Sensing*, 2010.
- Xu, Khajehnejad, Avestimehr, and Hassibi, “Breaking through the Thresholds: an Analysis for Iterative Reweighted  $\ell_1$  Minimization via the Grassmann Angle Framework,” *Proc. Allerton Conference on Communication, Control, and computing*, Sept. 2009.