

## 5. ARITHMETIC OF THE NATURALS

In order to bring the abstract treatment of the naturals closer to our intuition, we will now define the basic operations of *addition*, *multiplication*, *powers* etc on the naturals and prove the standard relations between them.

**5.1 Addition.**

We will spend most of the time on addition as other operations are handled analogously. Pick  $m \in \mathbb{N}$  and invoke the recursion principle in Theorem 4.5 for the choice  $E := \mathbb{N}$ ,  $a := m$  and  $h := S$  to define  $\{X_n : n \in \mathbb{N}\}$  such that

$$X_0 = m \quad \text{and} \quad \forall n \in \mathbb{N}: X_{S(n)} = S(X_n). \quad (5.1)$$

Then we denote

$$m + n := X_n. \quad (5.2)$$

As consequence of the construction (5.1) we get a symbol  $m + n$  satisfying

- (1)  $\forall m \in \mathbb{N}: m + 0 = m$ , and
- (2)  $\forall m, n \in \mathbb{N}: m + S(n) = S(m + n)$ .

From these observation we now derive further facts about addition relying, predominantly, on the Induction Principle.

We will now prove that the operation  $m, n \mapsto m + n$  is commutative. We begin by:

**Lemma 5.1**  $\forall m \in \mathbb{N}: 0 + m = m$

*Proof.* Let  $P_m$  denote the logical proposition  $0 + m = m$ . Then  $P_0$  is TRUE because (1) above implies  $0 + 0 = 0$ . Next assume  $P_m$  holds for some  $m \in \mathbb{N}$ . Then

$$0 + S(m) \stackrel{(2)}{=} S(0 + m) \stackrel{P_m}{=} S(m). \quad (5.3)$$

It follows that  $P_m \Rightarrow P_{S(m)}$ . By the Induction Lemma,  $\{m \in \mathbb{N}: 0 + m = m\} = \mathbb{N}$ .  $\square$

Next we need:

**Lemma 5.2**  $\forall m, n \in \mathbb{N}: m + S(n) = S(m) + n$

*Proof.* Fix  $m \in \mathbb{N}$  and let  $P_n$  be the statement  $m + S(n) = S(m) + n$ . Since

$$m + S(0) \stackrel{(2)}{=} S(m + 0) \stackrel{(1)}{=} S(m) \stackrel{(1)}{=} S(m) + 0 \quad (5.4)$$

we get that  $P_0$  is TRUE. Next assume that  $P_n$  is TRUE for some  $n \in \mathbb{N}$ . Then

$$m + S(S(n)) \stackrel{(2)}{=} S(m + S(n)) \stackrel{P_n}{=} S(S(m) + n) \stackrel{(2)}{=} S(m) + S(n) \quad (5.5)$$

implying  $P_{S(n)}$ . Hence,  $\forall n \in \mathbb{N}: P_n \Rightarrow P_{S(n)}$  and, by the Induction lemma,  $\{n \in \mathbb{N}: m + S(n) = S(m) + n\} = \mathbb{N}$ . As this holds for all  $m \in \mathbb{N}$ , we are done.  $\square$

Hence we finally conclude:

**Proposition 5.3** (Commutativity of addition)

$$\forall m, n \in \mathbb{N}: m + n = n + m \quad (5.6)$$

*Proof.* Let  $Q_m$  be the statement  $\forall n \in \mathbb{N}: m + n = n + m$ . Then  $Q_0$  is TRUE by (1) and Lemma 5.1. Assume now that  $Q_m$  is TRUE. Then for any  $n \in \mathbb{N}$ ,

$$S(m) + n \stackrel{\text{Lemma 5.2}}{=} m + S(n) \stackrel{(2)}{=} S(m + n) \stackrel{Q_m}{=} S(n + m) \stackrel{(2)}{=} n + S(m). \quad (5.7)$$

It follows that  $Q_m \Rightarrow Q_{S(m)}$ . By induction,  $\{m \in \mathbb{N}: Q_m\} = \mathbb{N}$ .  $\square$

Similarly we also prove that the operation  $m, n \mapsto m + n$  is associative:

**Proposition 5.4** (Associativity of addition)

$$\forall m, n, k \in \mathbb{N}: m + (n + k) = (m + n) + k \quad (5.8)$$

*Proof.* Left as a homework exercise. Commutativity should not be required.  $\square$

## 5.2 Ordering of the naturals.

A useful property of addition is that it acts injectively:

**Lemma 5.5**  $\forall m, n, \ell \in \mathbb{N}: m + n = m + \ell \Rightarrow n = \ell$

*Proof.* Let  $P_m$  be the statement  $\forall n, \ell \in \mathbb{N}: m + n = m + \ell \Rightarrow n = \ell$ . By (1) and Lemma 5.1,  $P_0$  is TRUE. Now assume  $P_m$  is TRUE for some  $m \in \mathbb{N}$ . For  $n, \ell \in \mathbb{N}$  are such that

$$S(m) + n = S(m) + \ell \quad (5.9)$$

then Lemma 5.2 implies  $S(m) + n = m + S(n)$  and  $S(m) + \ell = \ell + S(n)$

$$m + S(n) \stackrel{\text{Lemma 5.2}}{=} S(m) + n = S(m) + \ell \stackrel{\text{Lemma 5.2}}{=} m + S(\ell) \quad (5.10)$$

thus implying  $S(n) = S(\ell)$  via  $P_m$ . But  $S$  is injective by P4 and so we get  $n = \ell$ . Hence  $P_m \Rightarrow P_{S(m)}$  and, by induction,  $P_m$  is TRUE for all  $m \in \mathbb{N}$ .  $\square$

This property implies that, given  $m \in \mathbb{N}$ , for each  $n \in \mathbb{N}$  the equation  $n = m + s$  has at most one solution for  $s$  in  $\mathbb{N}$ . We can formally describe the pairs  $(m, n) \in \mathbb{N} \times \mathbb{N}$  for which the solution exists by way of the *less than or equal* relation  $\leq$  defined by

$$m \leq n \quad \Leftrightarrow \quad \exists s \in \mathbb{N}: n = m + s. \quad (5.11)$$

Here are some properties of this relation:

**Lemma 5.6** *The relation  $\leq$  is reflexive, antisymmetric and transitive.*

*Proof.* Reflexivity is immediate from (1) and transitivity follows from the associativity of addition. So the main point to check is antisymmetry. For that assume  $m, n \in \mathbb{N}$  are such that  $m \leq n \wedge n \leq m$ . Then there are  $r, s \in \mathbb{N}$  such that  $m = n + s \wedge n = m + r$ . Putting these together and invoking the associativity of addition, we get  $n = n + (s + r)$ . Lemma 5.5 and (1) then force  $s + r = 0$ . By P3 and (2) above,  $r$  cannot be a successor and so  $r = 0$ . Then also  $s = 0$  whereby we conclude

$$\forall m, n \in \mathbb{N}: m \leq n \wedge n \leq m \Rightarrow m = n \quad (5.12)$$

meaning that  $\leq$  is antisymmetric.  $\square$

It easy to check that the following properties of  $\leq$  are true:

**Lemma 5.7** *We have*

$$\forall n \in \mathbb{N}: 0 \leq n \quad (5.13)$$

$$\forall n \in \mathbb{N}: n \leq S(n) \quad (5.14)$$

and

$$\forall m, n \in \mathbb{N}: m \leq n \Rightarrow S(m) \leq S(n) \quad (5.15)$$

*Proof.* Left to a homework exercise.  $\square$

An important point of the relation  $\leq$  is that it is *connex*, meaning that every pair of naturals are ordered one or the other way. This is usually phrased by saying that  $\leq$  is a *total ordering* in the following sense:

**Lemma 5.8** (Total-ordering of  $\mathbb{N}$ )

$$\forall m, n \in \mathbb{N}: m \leq n \vee n \leq m \quad (5.16)$$

*Proof.* Let  $P_m$  be the statement  $\forall n \in \mathbb{N}: m \leq n \vee n \leq m$ . Then  $P_0$  is TRUE by (5.13) so assume that  $P_m$  holds for some  $m \in \mathbb{N}$  and let  $n \in \mathbb{N}$ . If  $n \leq m$  or  $n = m$  then (5.14) and transitivity imply  $n \leq S(m)$ . In the opposite case we must have  $m \leq n$  (as  $P_m$  was assumed to hold) and  $m \neq n$ . The definition (5.11) and Lemma 4.2 then show existence of an  $r \in \mathbb{N}$  such that

$$n = m + S(r) \stackrel{\text{Lemma 5.2}}{=} S(m) + r \quad (5.17)$$

proving  $S(m) \leq n$  and thus also  $P_m \Rightarrow P_{S(m)}$ . Hence,  $P_m$  is TRUE for all  $m \in \mathbb{N}$ .  $\square$

Note that we can re-state Lemma 5.8 as saying that at least one of  $n = m + r$  or  $m = n + r$  has a solution for  $r$  in the naturals.

### 5.3 Multiplication, powers, factorial.

Moving to a definition of multiplication, pick  $m \in \mathbb{N}$  and use Theorem 4.5 with the choices  $E := \mathbb{N}$ ,  $h(r) := r + m$  and  $a := 0$  to construct  $\{X_n: n \in \mathbb{N}\}$  such that

$$X_0 = 0 \quad \wedge \quad \forall n \in \mathbb{N}: X_{S(n)} = X_n + m. \quad (5.18)$$

We will write  $n \cdot m$  for  $X_n$  and thus get

$$0 \cdot m = 0 \quad \wedge \quad \forall n \in \mathbb{N}: S(n) \cdot m = n \cdot m + m \quad (5.19)$$

We also define the *unity* in  $\mathbb{N}$  by

$$1 := S(0) \quad (5.20)$$

and observe that

$$S(n) = S(n + 0) = n + S(0) = n + 1 \quad (5.21)$$

which will eventually allow us to drop the notation using the successor function and write it as “plus one” instead. The following properties are then checked:

**Proposition 5.9** (Properties of multiplication on  $\mathbb{N}$ ) *We have:*

- (1) (Commutative law)  $\forall m, n \in \mathbb{N}: m \cdot n = n \cdot m$ ,
- (2) (Associative law)  $\forall m, n, k \in \mathbb{N}: (m \cdot n) \cdot k = m \cdot (n \cdot k)$ ,
- (3) (Distributive law)  $\forall m, n, k \in \mathbb{N}: (n + k) \cdot m = (n \cdot m) + (k \cdot m)$
- (4) (Zero and unity)  $\forall m \in \mathbb{N}: 0 \cdot m = 0 \wedge 1 \cdot m = m$
- (5) (Injectivity)  $\forall m, n, k \in \mathbb{N}: k \neq 0 \wedge k \cdot m = k \cdot n \Rightarrow m = n$

*Proof.* A somewhat tedious but doable exercise that we leave to the reader. □

Multiplication also behaves nicely around the total ordering relation:

**Lemma 5.10** *We have*

$$\forall m, n, r \in \mathbb{N}: m \leq n \Rightarrow r \cdot m \leq r \cdot n \quad (5.22)$$

*Proof.* Left to homework exercise. □

With multiplication in place, we can now define natural *powers*. Here we pick  $m \in \mathbb{N}$  and use Theorem 4.5 to construct  $\{m^n : n \in \mathbb{N}\}$  satisfying

$$m^0 = 1 \quad \wedge \quad \forall n \in \mathbb{N}: m^{S(n)} = m \cdot m^n. \quad (5.23)$$

Note that this entails  $m^0 = 1$  (even for  $m = 0$ ) while  $0^n = 0$  for  $n \neq 0$ . Similarly,  $1^n = 1$  for all  $n \in \mathbb{N}$ . The following properties will again be of relevance:

**Lemma 5.11 (Powers)** *Let  $m \in \mathbb{N} \setminus \{0\}$ . Then*

- (1)  $\forall r, s \in \mathbb{N}: m^{r+s} = m^r \cdot m^s,$
- (2)  $\forall r, s \in \mathbb{N}: m^{r \cdot s} = (m^r)^s.$

*Proof.* Proved readily by induction. □

A related construction permits us to construct the *factorial* of  $n$ , with notation  $n!$ , by imposing

$$0! = 1 \quad \text{and} \quad \forall n \in \mathbb{N}: S(n)! = S(n) \cdot n! \quad (5.24)$$

By (5.21), the statement in the second part can be written as  $(n+1)! = (n+1) \cdot n!$ , which is the recursive form of the informal expression  $n! = n \cdot (n-1) \cdots 1$ .

Factorials appear frequently in combinatorial arguments (indeed,  $n!$  is the number of permutations of  $n$  elements) but also appears in analytic expressions (thanks to, for instance, Taylor's theorem).