

MIDTERM 2 (MATH 61, SPRING 2017)

Your Name: .....

UCLA id: .....

Math 61 Section: .....

Date: .....

**The rules:**

You MUST simplify completely and BOX all answers with an **INK PEN**.  
You are allowed to use only this paper and pen/pencil. No calculators.  
No books, no notebooks, no phones, no web access. You MUST write your name and UCLA id. Except for the last problem, you MUST write out your logical reasoning and/or proof in full. You have exactly 50 minutes.

**Warning:** those caught violating the rules get automatic 10% score deduction.

**Points:**

1 |  
2 |  
3 |  
4 |  
5 |

.....  
**Total:** (out of 100)

**Problem 1.** (15 points)

Solve the following LHR:  $a_{n+1} = 5a_n - 6a_{n-1}$ ,  $a_1 = 9$ ,  $a_2 = 20$ .

**Solution:**

characteristic polynomial:  $x^2 = 5x - 6 \Rightarrow x_1 = 2, x_2 = 3$

We have:  $a_n = A \cdot 2^n + B \cdot 3^n$

Plug in  $a_1$  and  $a_2$  we get  $A = \frac{7}{2}, B = \frac{2}{3}$

**Problem 2.** (20 points)

Find the number of subgraphs of  $G$  isomorphic to  $H$ , where

a)  $G = C_{12}$ ,  $H = P_5$ .

**Answer:** 12

b)  $G = K_{9,9}$ ,  $H = C_6$ .

**Answer:**  $\binom{9}{3}^2 \times 6$

**Hint:** First choose  $K_{3,3}$ , then there are 6  $C_6$  in it

c)  $G = K_{9,9}$ ,  $H = C_7$ .

**Answer:** 0

d)  $G = H_3$  (the 3-cube),  $H = C_4$ .

**Answer:** 6 (only the 6 faces of the cube count)

e)  $G = H_3$  (the 3-cube),  $H = P_4$ .

**Answer:** 48 (Fix the central edge of  $P_4$ : There are 12 different central edges for  $P_4$ , each of them could induce 4  $P_4$ )

**Problem 3.** (20 points)

For each of these sequences, either draw a simple graph with this score (degree sequence), or explain why there is no such graph.

a)  $(3,3,3,1,1,1)$

**Exist:** Just draw a triangle and 3 extra vertices (one to one) connecting to each vertex of the triangle.

b)  $(4,4,4,1,1,1)$

**Not Exist:** sum of degree is odd.

c)  $(5,5,5,3,2,1)$

**Not Exist:** sum of degree is odd.

d)  $(4,4,\dots,4) \leftarrow 16$  numbers.

**Exist:** You can draw 2 copies of  $K_{4,4}$

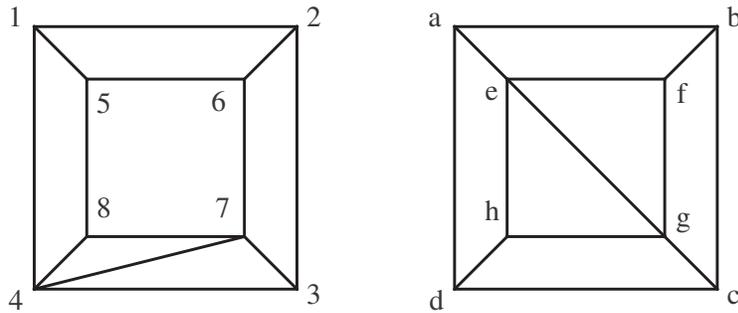
e)  $(3,3,3,3,2,2,2,2)$

**Exist:** You can draw a  $C_8$  and add 2 chords

**Problem 4.** (15 points)

Decide whether the following pairs of graphs on 8 vertices are isomorphic or non-isomorphic.

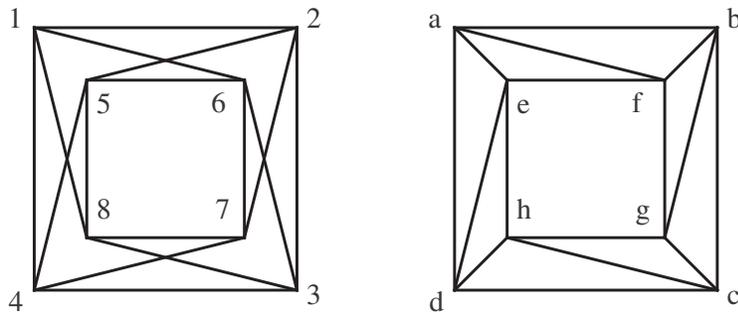
a)



**Isomorphic:** for example:  $\{1,2,3,4,5,6,7,8\} \rightarrow \{a,b,f,e,d,c,g,h\}$

You need to make sure that if 4 vertices construct a  $C_4$  then their mapped vertices also form a  $C_4$ .

b)



**Not isomorphic:** the left graph does not have  $C_3$  as subgraph but the right one does.

**Important:** In case of isomorphism, you must present a bijection. In case non-isomorphism, you must present an argument.

**Problem 5.** (30 points, 2 points each) **TRUE or FALSE?**

Circle correct answers with ink. No explanation is required or will be considered.

- T F** (1) The hypercube graph  $H_4$  contains an Eulerian cycle.
- T F** (2) The hypercube graph  $H_4$  contains a Hamiltonian cycle.
- T F** (3) A subgraph of a connected graph is always connected.
- T F** (4) A subgraph of a disconnected graph is always disconnected.
- T F** (5) The sum of degrees of  $K_{\ell,\ell}$  is  $\ell^2$ .
- T F** (6) The sum of degrees of a graph on  $n \geq 3$  vertices is smaller than  $n^2 - 1$ .
- T F** (7) Computing the number of walks of given length in a graph can be done efficiently.
- T F** (8) Deciding whether a graph has Hamiltonian cycle can be done efficiently.
- T F** (9) The number of (shortest) grid walks  $(0,0) \rightarrow (5,5)$  which do not go through  $(1,3), (3,4), (4,1), (1,5)$  is  $> 100$ .
- T F** (10) Graph  $H$  is a subgraph of  $G$ . Graph  $H$  contains a Hamiltonian cycle. Then  $G$  contains a Hamiltonian cycle.
- T F** (11) Graph  $H$  is a subgraph of  $G$ . Graph  $G$  contains a Hamiltonian cycle. Then  $H$  contains a Hamiltonian cycle.
- T F** (12) Graph  $K_{50,50}$  contains a subgraph isomorphic to  $K_{10}$ .
- T F** (13) Graph  $K_{50,50}$  contains a subgraph isomorphic to  $K_{10,10}$ .
- T F** (14) The number of walks  $1 \rightarrow 1$  of length  $k$  in a graph  $G$  on  $n$  vertices can be computed via matrix  $(A_G)^n$ .
- T F** (15) Isomorphic graphs have the same number of Eulerian circuits.

**T T F F F T T F F F F F T F T**