## **Closure operators**

**Definition.** Let *S* be a set. A **closure operator** on *S* is a function  $cl : \mathcal{P}(S) \to \mathcal{P}(S)$  such that for all  $A, B \subseteq S$ ,

- $A \subseteq cl(A)$  (cl is *extensive*)
- $A \subseteq B \implies cl(A) \subseteq cl(B)$  (cl is increasing)
- cl(cl(A)) = cl(A) (cl is *idempotent*).

We call cl(A) the **closure** of *A*; a set  $C \subseteq S$  that is the closure of some subset of *S* is called a **closed set** of cl.

These three defining conditions are in fact equivalent to a single biconditional, which is sometimes easier to verify.

**Theorem 1.** A function  $f : \mathcal{P}(S) \to \mathcal{P}(S)$  is a closure operator on *S* if and only if  $A \subseteq f(B) \iff f(A) \subseteq f(B)$  for all  $A, B \subseteq S$ .

*Proof.* Let  $A, B \subseteq S$ , and suppose first that f is a closure operator. Then  $A \subseteq f(B) \implies f(A) \subseteq f(f(B))$  (since f is increasing)  $\implies f(A) \subseteq f(B)$  (since f is idempotent). Conversely,  $f(A) \subseteq f(B) \implies A \subseteq f(B)$  (since f is extensive).

On the other hand, if *f* satisfies the biconditional above, then  $f(A) \subseteq f(A) \implies A \subseteq f(A)$ , so *f* is extensive. Using extensivity, we obtain  $A \subseteq B \implies A \subseteq f(B) \implies f(A) \subseteq f(B)$ , so *f* is increasing. Finally,  $f(A) \subseteq f(A) \implies f(f(A)) \subseteq f(A)$ , where the reverse inclusion holds by extensivity, so *f* is idempotent.

The following characterization of closed sets will be used extensively, and is frequently taken as the *definition* of a closed set.

**Proposition.** A set  $A \subseteq S$  is closed *if and only if* it is equal to its own closure, i.e., cl(A) = A.

*Proof.* If *A* is closed, then A = cl(B) for some  $B \subseteq S$ , whence cl(A) = cl(cl(B)) = cl(B) = A (using idempotence). The converse follows immediately from the definition of a closed set.

**Corollary.** *S* is a closed set.

*Proof.* Clearly  $cl(S) \subseteq S$ , and the reverse inclusion follows from extensivity.

## **Moore collections**

**Definition.** Let *S* be a set. A collection  $\mathcal{C} \subseteq \mathcal{P}(S)$  of subsets of *S* is called a **Moore collection** if any intersection of elements in  $\mathcal{C}$  is also in  $\mathcal{C}$ .

The following two theorems exhibit a 'duality' between closure operators and Moore collections.

**Theorem 2.** Let cl be a closure operator on *S*. If  $\mathcal{C}$  is the collection of all closed sets of cl, then  $\mathcal{C}$  is a Moore collection, and for any  $A \subseteq S$ ,

$$cl(A) = \bigcap \{ C \in \mathcal{C} : C \supseteq A \}.$$

In other words, any intersection of closed sets is also closed, and cl(A) is the intersection of all closed sets containing *A*.

*Proof.* Let  $\{C_i\}_{i \in I} \subseteq \mathbb{C}$  be a collection of closed sets, and define  $C = \bigcap_{i \in I} C_i$ . If  $x \in cl(C)$  and  $i \in I$ , then  $C \subseteq C_i \implies cl(C) \subseteq cl(C_i) \implies x \in cl(C_i) = C_i$ , so  $cl(C) \subseteq C$ . As the reverse inclusion holds by extensivity, *C* is equal to its own closure, and hence is closed. Thus  $\mathbb{C}$  is a Moore collection.

Now if  $A \subseteq S$ , let *B* denote the set on the right-hand side above. Clearly  $A \subseteq B$  (an intersection of sets containing *A* must also contain *A*), so  $cl(A) \subseteq cl(B) = B$ , since *B* is an intersection of closed sets. On the other hand,  $cl(A) \in C$  and  $cl(A) \supseteq A$ , so  $B \subseteq cl(A)$ .

*Remark.* cl(A) is therefore the "smallest closed set containing *A*", in that it is contained in any other closed set containing *A*.

**Theorem 3.** Let  $\mathcal{C} \subseteq \mathcal{P}(S)$  be a Moore collection. Then the function  $f : \mathcal{P}(S) \to \mathcal{P}(S)$  given by

$$f(A) := \bigcap \{ C \in \mathcal{C} : C \supseteq A \}$$

is a closure operator on S, and the collection of closed sets of f is C.

*Proof.* By Theorem 1, it suffices to prove that  $A \subseteq f(B) \iff f(A) \subseteq f(B)$  for all  $A, B \subseteq S$  to show that f is a closure operator on S. Suppose first that  $A \subseteq f(B)$ . Then  $f(B) \in C$ , as it is an intersection of elements of the Moore collection C, and  $f(B) \supseteq A$ , so  $f(A) \subseteq f(B)$ . On the other hand, if  $f(A) \subseteq f(B)$ , we must also have  $A \subseteq f(B)$ , since  $A \subseteq f(A)$ . Hence f is a closure operator on S.

Now if  $C \subseteq S$  is closed, then  $C = f(C) \in \mathbb{C}$  since  $\mathbb{C}$  is a Moore collection. Conversely, if  $C \in \mathbb{C}$ , then  $C \supseteq C \implies f(C) \subseteq C$ , and the reverse inclusion holds by extensivity. Thus the closed sets of *f* are exactly the sets of  $\mathbb{C}$ .

Thus, any closure operator is determined by a Moore collection (namely, the collection of its closed sets), and any Moore collection is determined by a closure operator.

## Examples

If *V* is a vector space, it is easily verified that the collection of all subspaces of *V* is a Moore collection. The closure operator that this collection defines is none other than the *span* of a set of vectors. In other words, if  $A \subseteq V$ , then span(A) is *the intersection of all subspaces of V containing A*, or equivalently, *the smallest subspace of V containing A*.

Other examples are tabulated below:

Set	Moore collection	Closure operator	Notation
A vector space V	Subspaces of V	Span of A	span(A)
A topological space <i>X</i>	Closed sets of X	(Topological) closure of A	Ā
A group G	Subgroups of G	Subgroup generated by A	$\langle A \rangle$
A group G	Normal subgroups of G	Normal closure of A	$\langle A^G \rangle$
A ring R	Subrings of R	Subring generated by A	
A ring R	Ideals of R	Ideal generated by A	(A)
All subsets of a set <i>X</i>	$\sigma$ -algebras on X	$\sigma$ -algebra generated by $A$	$\sigma(A)$
$\mathbb{R}^n$	Convex subsets of $\mathbb{R}^n$	Convex hull of <i>A</i>	conv(A)
$\mathbb{R}^{n}$	Affine subsets of $\mathbb{R}^n$	Affine hull of A	aff(A)
$\mathbb{R}^{n}$	Convex cones in $\mathbb{R}^n$ containing the origin	Conical hull of A	cone(A)
All binary relations on a set <i>X</i>	Transitive binary relations on <i>X</i>	Transitive closure of <i>A</i>	$A^+$
All words over an alphabet $\Sigma$	Languages containing the empty string that are closed under string concatenation	Kleene star of A	A*