

Determinacy and Large Cardinals

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Use **PageDown** or the down arrow to scroll through slides.
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Players I and II alternate playing numbers $a_n \in \omega$,

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II	a_1

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<i>I</i>	a_0	a_2	a_4
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<i>I</i>	a_0	a_2	a_4
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Players *I* and *II* alternate playing numbers $a_n \in \omega$, forming together an infinite sequence $z = \langle a_0, a_1, a_2, \dots \rangle \in \omega^\omega$.

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Players I and II alternate playing numbers $a_n \in \omega$, forming together an infinite sequence $z = \langle a_0, a_1, a_2, \dots \rangle \in \omega^\omega$.

If z belongs to A then player I wins.

If z does not belong to A then player II wins.

$G_\omega(A)$ is *determined* if one of the players has a winning strategy.

(A *strategy* is a complete recipe that instructs the player precisely how to play in each conceivable situation.)

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But determinacy for *definable* sets is: (1) true; and (2) useful.

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$\{\text{Borel sets}\} \subsetneq \{\text{analytic sets}\} \subsetneq \{\text{projective sets}\}$.

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$L(\mathbb{R})$ is the smallest model of set theory which contains all the reals and all the ordinals. It is obtained as the union $\bigcup_{\alpha \in \text{ON}} L_\alpha(\mathbb{R})$ where:

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$\{\text{projective sets}\} \subset L_1(\mathbb{R})$.

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Theorems 1 and 2 are in ZFC, the basic system of axioms for set theory.

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Theorems 1 and 2 are in ZFC, the basic system of axioms for set theory.

Theorems 3, 4, and 5 require large cardinal axioms.

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Theorem 8 (Davis 1964) *Assume $\text{det}(\Gamma)$. Let $A \in \Gamma$. Then either A is countable or else it contains a perfect set.*

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Every Σ_n^1 set A is obtained from an underlying open set using negations and existential quantifiers. A is (*lightface*) Σ_n^1 if the underlying open set is recursive. Similarly with Π_n^1 .

Γ has the *reduction property* if for any $A, B \in \Gamma$, there are $A' \subset A$, $B' \subset B$ in Γ , so that $A' \cup B' = A \cup B$ and $A' \cap B' = \emptyset$.

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Theorem 9 (Kuratowski 1936) *The pointclasses Π_1^1 and Σ_2^1 have the reduction property.*

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In fact they did more. They obtained a fundamental property, the prewellordering property, which implies reduction.

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The pointclasses in Theorem 10 are therefore precisely the pointclasses $\partial^{(n)}\Pi_1^1$, $n < \omega$.

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Theorem 12 (Steel–Van Wesep–Woodin) *Assume $AD^{L(\mathbb{R})}$. Then it is consistent (with $AD^{L(\mathbb{R})}$ and AC) that $(\omega_2)^{L(\mathbb{R})} = \omega_2$, and hence $\delta_2^1 = \omega_2$.*

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Similar arguments show κ must be inaccessible, and in fact cannot be described from below in any absolute manner.

So the existence of non-trivial $\pi: V \rightarrow M \subset V$ cannot be proved in ZFC, and the first ordinal moved by π must be very large.

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(Using an ultrapower construction, the measurability of κ is equivalent to the existence of a total, non-principal, countably complete, 2-valued measure on κ .)

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Note: if κ is the first measurable cardinal, then κ is only κ^+ -strong.

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δ is a *Woodin cardinal* if for every $D \subset \delta$ there is $\kappa < \delta$ which is $<\delta$ -strong wrt D .

Let $\pi: V \rightarrow M$. Let $\kappa = \text{crit}(\pi)$ and $\lambda \leq \pi(\kappa)$. The (κ, λ) -*extender* induced by π is the function $E: \mathcal{P}(\kappa) \rightarrow \mathcal{P}(\lambda)$ defined by $E(X) = \pi(X) \cap \lambda$.

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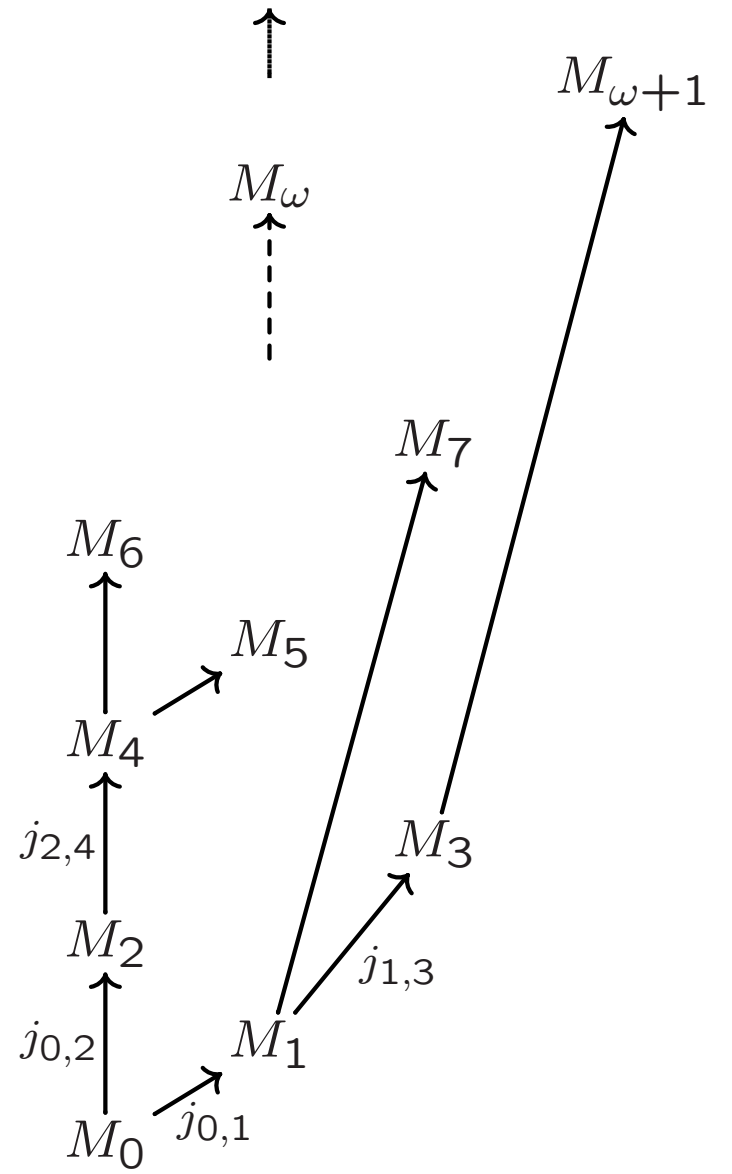
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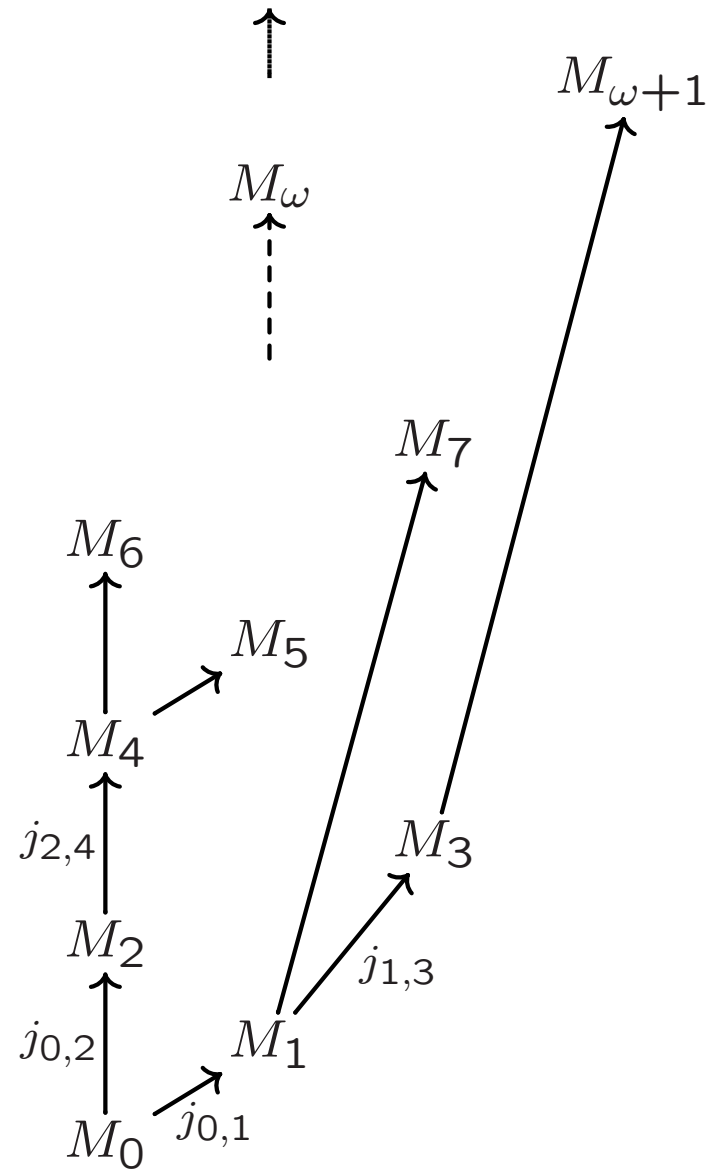
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This allows constructing iterated ultrapowers with non-linear base orders.

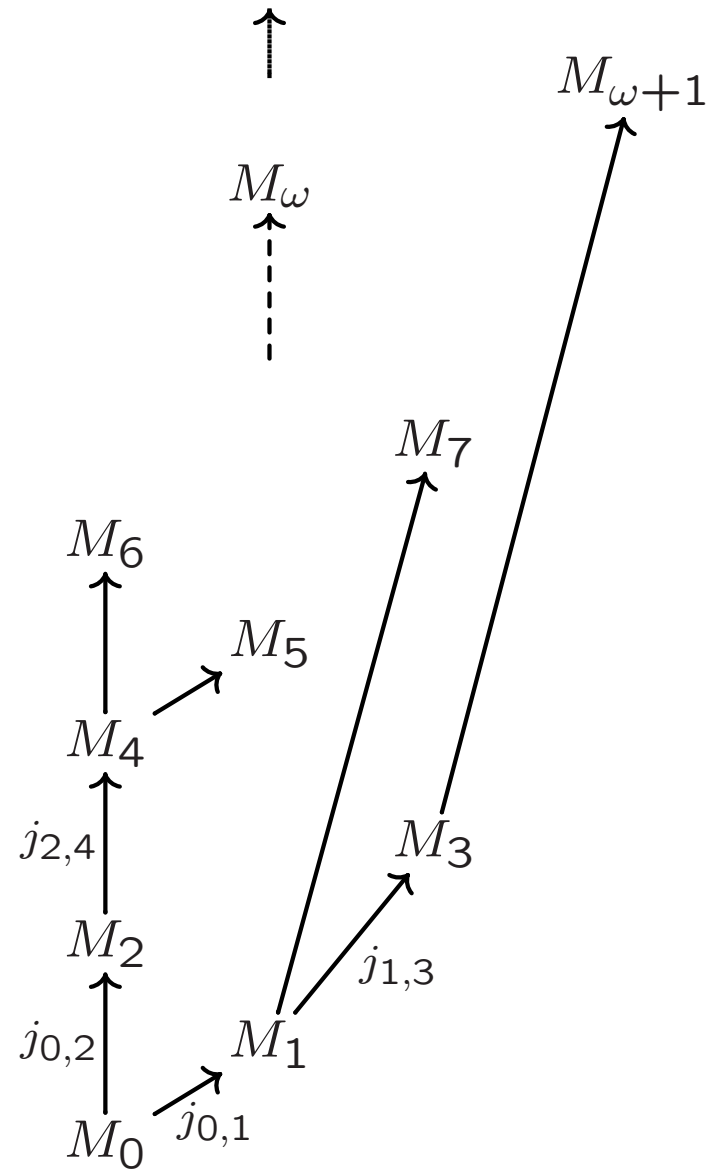


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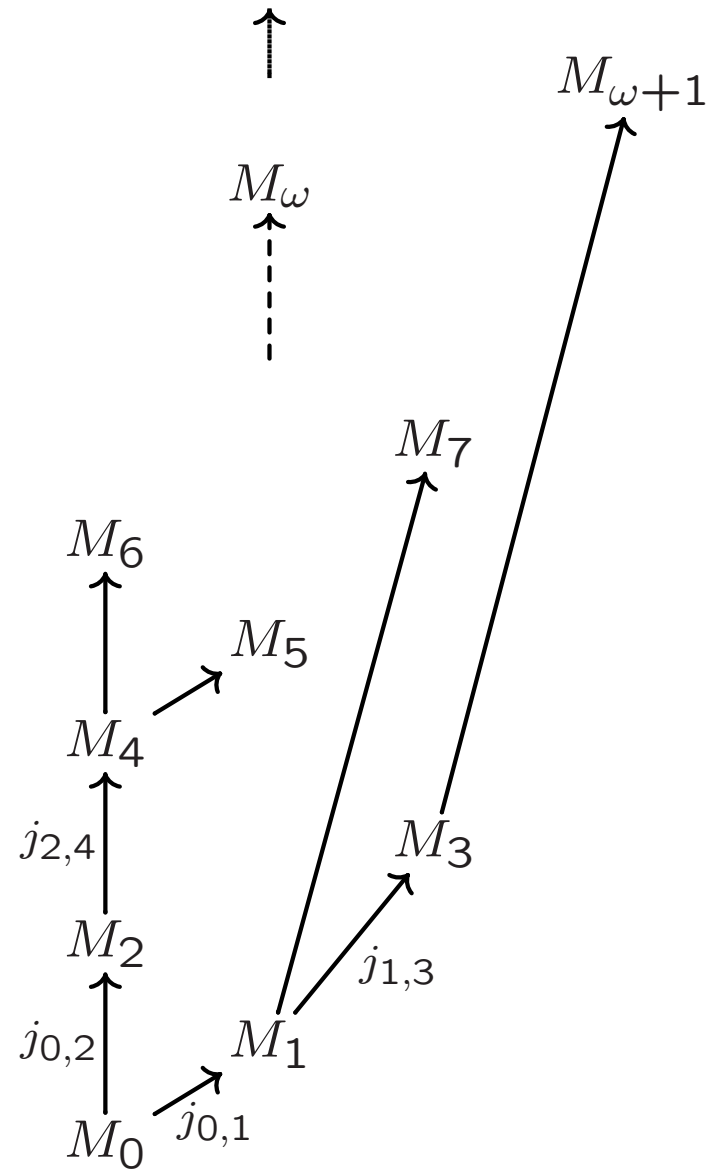
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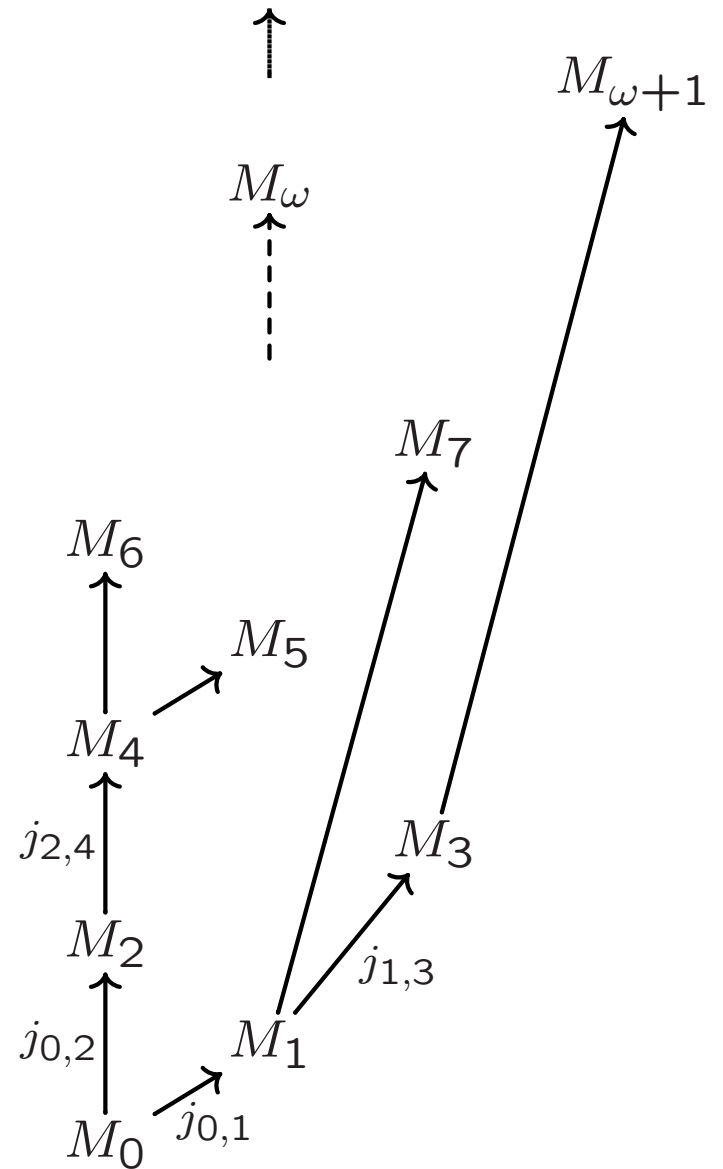
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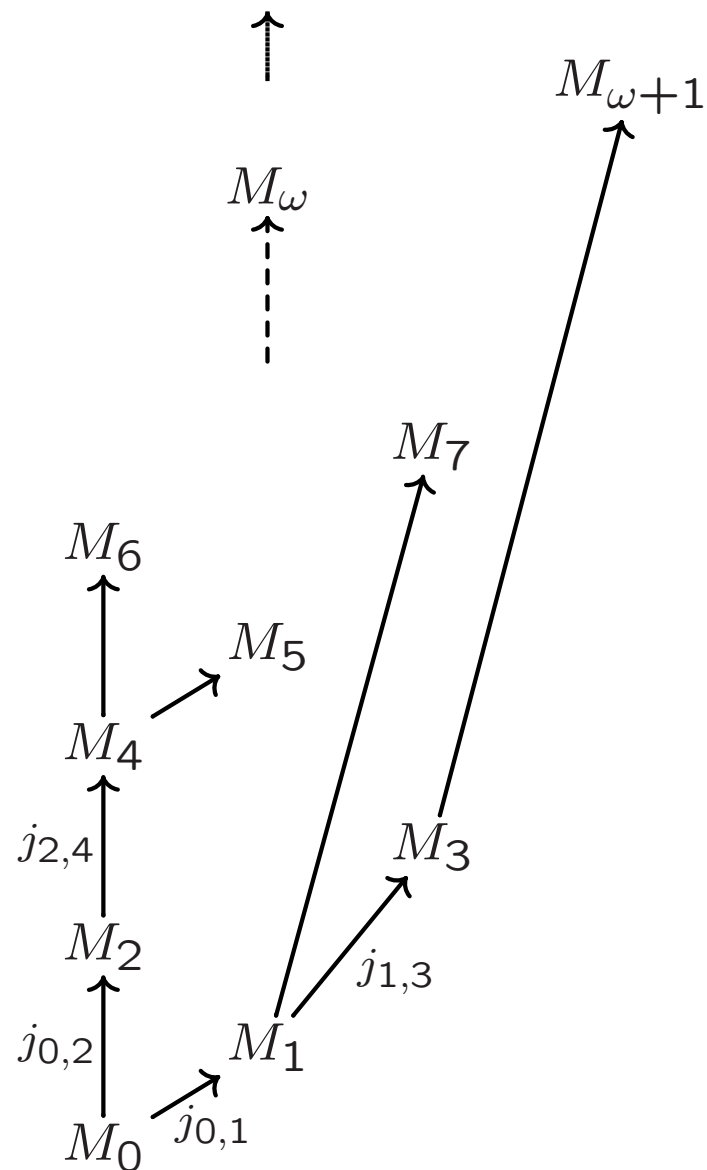
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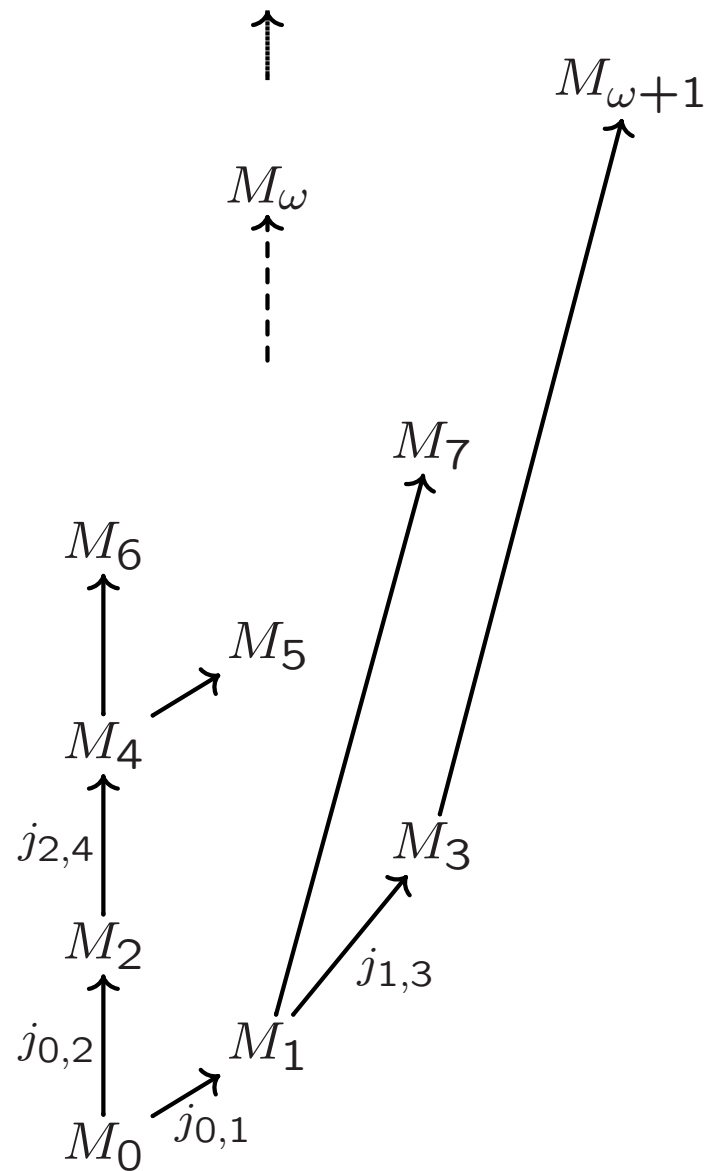
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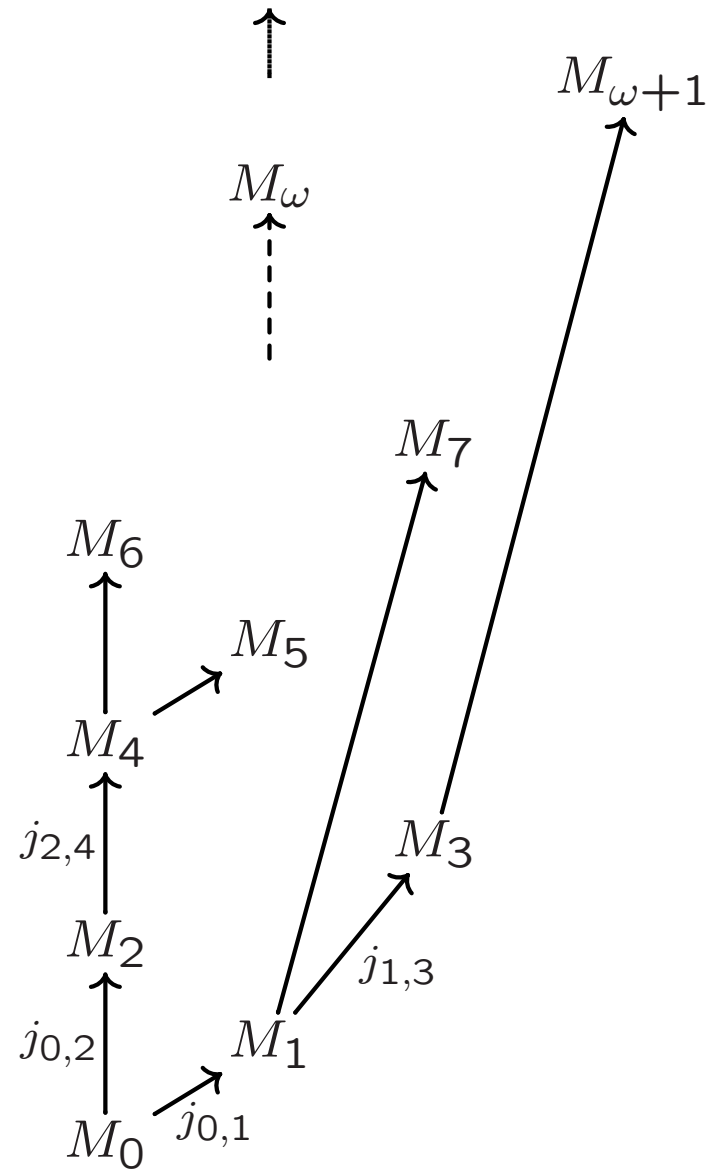
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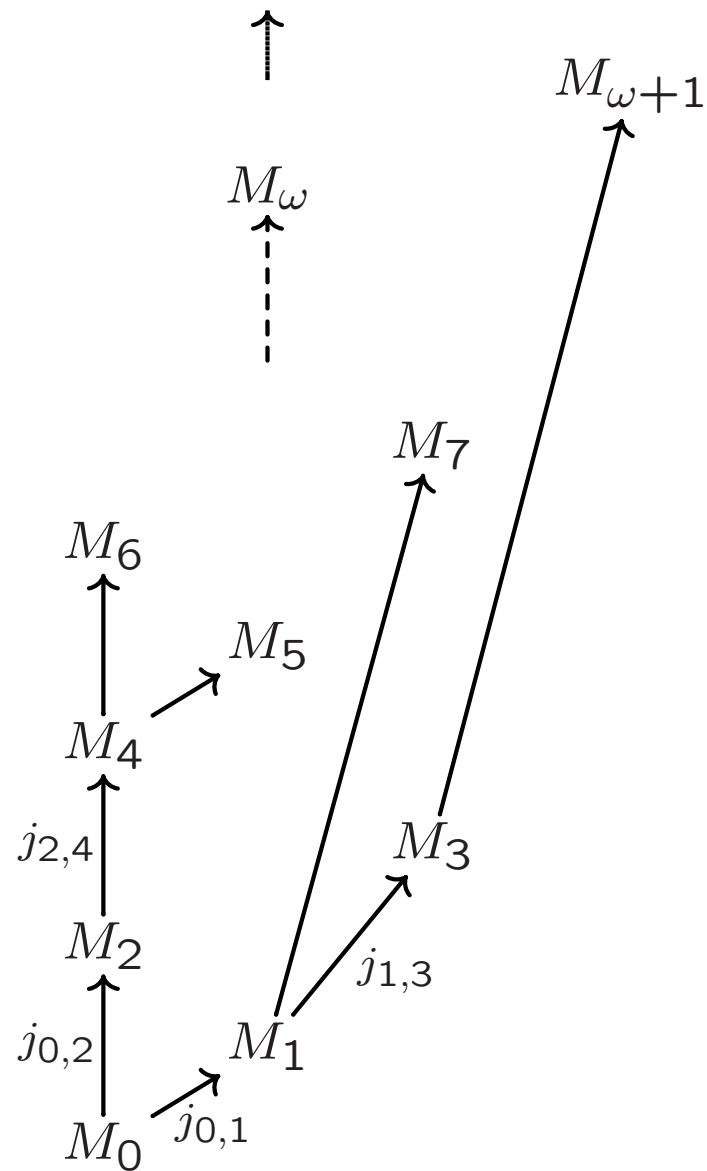
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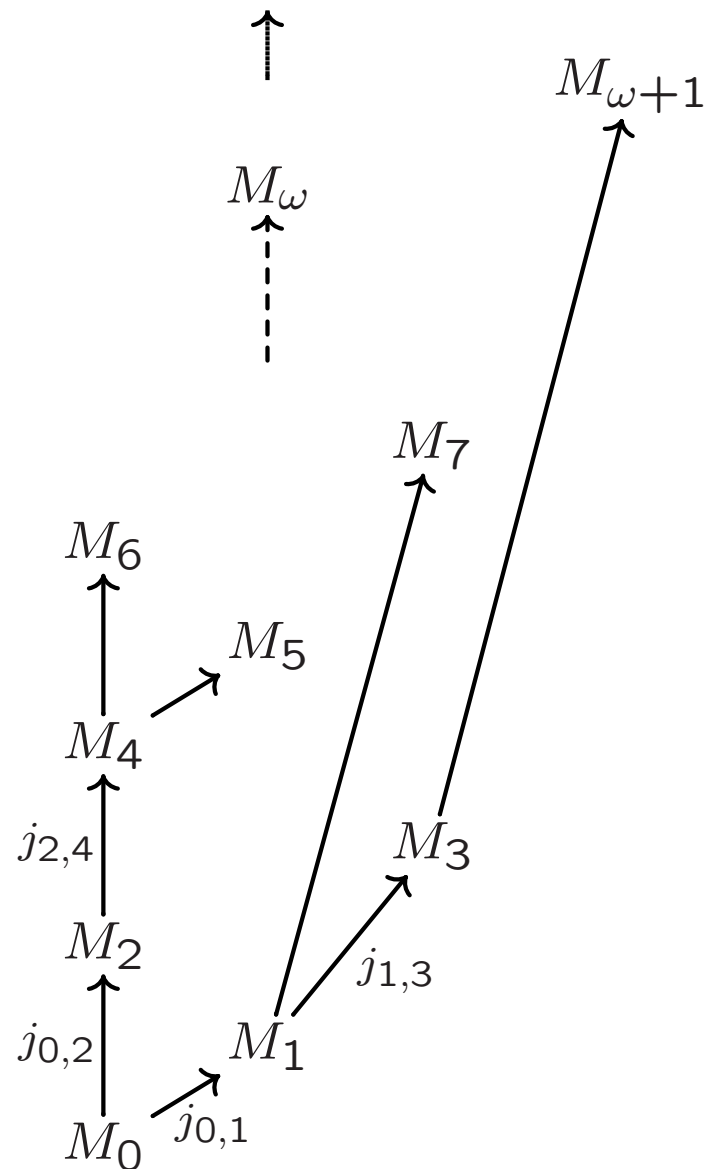


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The result is an *iteration tree* on M .



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In both cases Woodin cardinals in iterable inner models (rather than the actual universe V) are enough, and moreover *necessary*.

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But some make direct use of inner models for Woodin cardinals.

A set A is α - Π_1^1 if there is a sequence $\langle A_\xi \mid \xi < \alpha \rangle$ of Π_1^1 sets so that $x \in A$ iff the least ξ so that $x \notin A_\xi \vee \xi = \alpha$ is odd.

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(The hierarchy generated by this definition is the *difference hierarchy* on Π_1^1 sets. If $\alpha = 2$ for example, then the condition states simply that $A = A_0 - A_1$.)

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Theorem known previously for odd n , not using large cardinals (Kechris–Woodin).

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Theorem 15 (Neeman, Woodin) *Assume $\text{AD}^{L(\mathbb{R})}$. Then it is consistent (with $\text{AD}^{L(\mathbb{R})}$ and the axiom of choice) that $\delta_3^1 = \omega_2$.*

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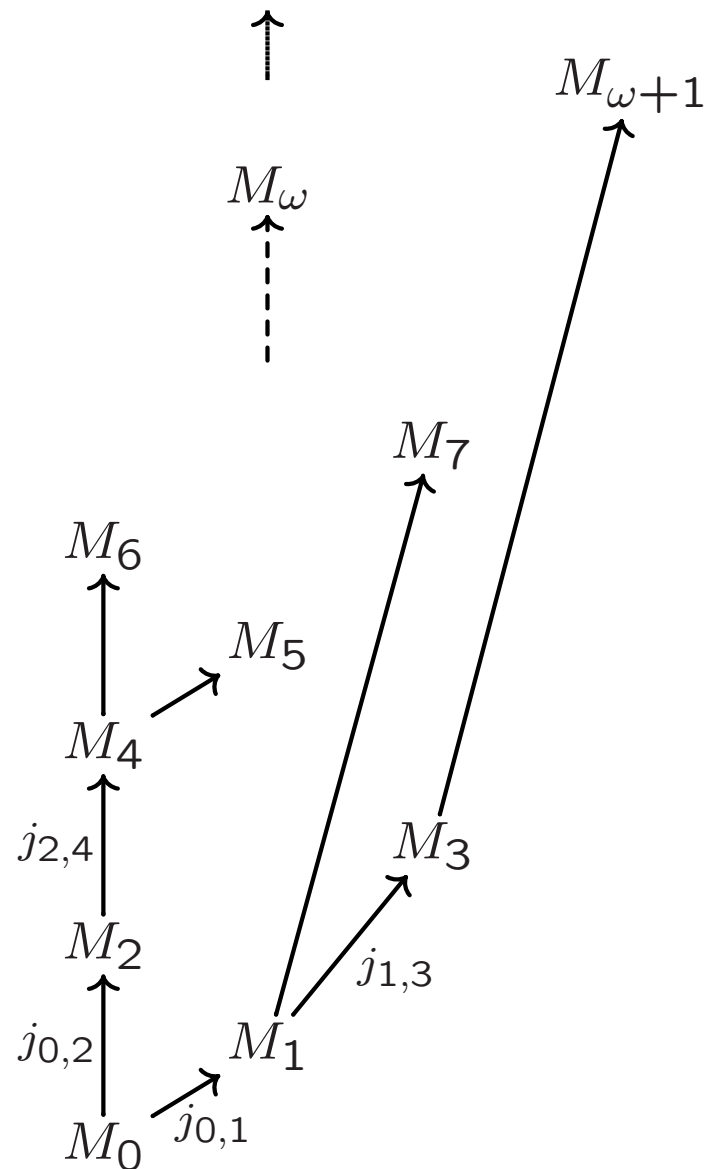
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M is *iterable* if these choices can be made in a way that secures the wellfoundedness of all the models created.

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Comparisons through iterated ultrapowers show that any two ways to witness θ are compatible.

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Theorem 17 (Neeman) *Let B_i be a recursive enumeration of the $\mathfrak{D}^{(n)}(<\omega^2 - \Pi_1^1)$ sets. Suppose a sharp for n Woodin cardinals exists. Then all $\mathfrak{D}^{(n)}(<\omega^2 - \Pi_1^1)$ games are determined, and $\{i \mid I \text{ has a w.s. in } G_\omega(B_i)\}$ is recursively isomorphic to the theory of the sharp for n Woodin cardinals.*

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The connection (with analogues for ω Woodin cardinals) is crucial for Theorems 13–15.

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Let $[\vec{S}]$ denote the set

$$\{ \langle \alpha_0, \dots, \alpha_{k-1} \rangle \in [\omega_1]^{<\omega} \mid (\forall i < k) \alpha_i \in S_{\langle \alpha_0, \dots, \alpha_{i-1} \rangle} \}.$$

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If neither condition holds then both players lose.

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The theorem establishes a precise analogue of Theorems 16 and 17, but for embeddings concentrating on Woodin cardinals and for games of length ω_1 .

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Games motivated by Theorem 18 were used by Woodin in results on Σ_2^2 absoluteness. Other games similar to those in the theorem are enough to capture the theory of superstrong cardinals. But there are no determinacy proofs for these games from large cardinals, and indeed there are some negative results (Larson).

The End

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