

# Determinacy and Large Cardinals

Itay Neeman

Department of Mathematics

University of California Los Angeles

Los Angeles, CA 90095-1555

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Players  $I$  and  $II$  alternate playing numbers  $a_n \in \omega$ , forming together an infinite sequence  $z = \langle a_0, a_1, a_2, \dots \rangle \in \omega^\omega$ .

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If  $z$  belongs to  $A$  then player  $I$  wins.

If  $z$  does not belong to  $A$  then player  $II$  wins.

$G_\omega(A)$  is *determined* if one of the players has a winning strategy.

(A *strategy* is a complete recipe that instructs the player precisely how to play in each conceivable situation.)

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But determinacy for *definable* sets is: (1) true; and (2) useful.

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Following standard abuse of notation identify it with  $\mathbb{R}$ .

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$\{\text{Borel sets}\} \subsetneq \{\text{analytic sets}\} \subsetneq \{\text{projective sets}\}$ .

**Theorem 1 (Gale–Stewart 1953)** *All open sets are determined.*

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$L(\mathbb{R})$  is the smallest model of set theory which contains all the reals and all the ordinals. It is obtained as the union  $\bigcup_{\alpha \in \text{ON}} L_\alpha(\mathbb{R})$  where:

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$\{\text{projective sets}\} \subset L_1(\mathbb{R})$ .

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Theorems 1 and 2 are in ZFC, the basic system of axioms for set theory.



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Theorems 3, 4, and 5 require large cardinal axioms.

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**Theorem 6 (Banach, Oxtoby 1957)** Assume  $\text{det}(\Gamma)$ . Then all sets in  $\Gamma$  have the Baire property.

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**Theorem 8 (Davis 1964)** Assume  $\text{det}(\Gamma)$ . Let  $A \in \Gamma$ . Then either  $A$  is countable or else it contains a perfect set.

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In fact they did more. They obtained a fundamental property, the prewellordering property, which implies reduction.

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For a pointclass  $\Gamma$  set  $\partial \Gamma = \{\partial B \mid B \in \Gamma\}$ .

Easy to check  $\partial \Pi_n^1 = \Sigma_{n+1}^1$ , and (using determinacy)  $\partial \Sigma_n^1 = \Pi_{n+1}^1$ .

The pointclasses in Theorem 10 are therefore precisely the pointclasses  $\partial^{(n)} \Pi_1^1$ ,  $n < \omega$ .

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**Theorem 12 (Steel–Van Wesep–Woodin)** Assume  $AD^{L(\mathbb{R})}$ . Then it is consistent (with  $AD^{L(\mathbb{R})}$  and AC) that  $(\omega_2)^{L(\mathbb{R})} = \omega_2$ , and hence  $\delta_2^1 = \omega_2$ .

**Large cardinals:**

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$\kappa$  must be a limit cardinal. Otherwise have  $\tau < \kappa$  so that  $\kappa = \tau^+$ . But then by elementarity  $\pi(\kappa) = (\pi(\tau)^+)^M$ . Yet  $\pi(\tau) = \tau$ , so  $\pi(\kappa) = (\tau^+)^M = \kappa$ , contradiction.

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## Large cardinals:

Large cardinal axioms state the existence of (non-trivial) elementary embeddings  $\pi: V \rightarrow M \subset V$ .

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So the existence of non-trivial  $\pi: V \rightarrow M \subset V$  cannot be proved in ZFC, and the first ordinal moved by  $\pi$  must be very large.

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(Using an ultrapower construction, the measurability of  $\kappa$  is equivalent to the existence of a total, non-principal, countably complete, 2-valued measure on  $\kappa$ .)

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Note: if  $\kappa$  is the first measurable cardinal, then  $\kappa$  is only  $\kappa^+$ -strong.



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$\delta$  is a *Woodin cardinal* if for every  $D \subset \delta$  there is  $\kappa < \delta$  which is  $<\delta$ -strong wrt  $D$ .

Let  $\pi: V \rightarrow M$ . Let  $\kappa = \text{crit}(\pi)$  and  $\lambda \leq \pi(\kappa)$ . The  $(\kappa, \lambda)$ -*extender* induced by  $\pi$  is the function  $E: \mathcal{P}(\kappa) \rightarrow \mathcal{P}(\lambda)$  defined by  $E(X) = \pi(X) \cap \lambda$ .

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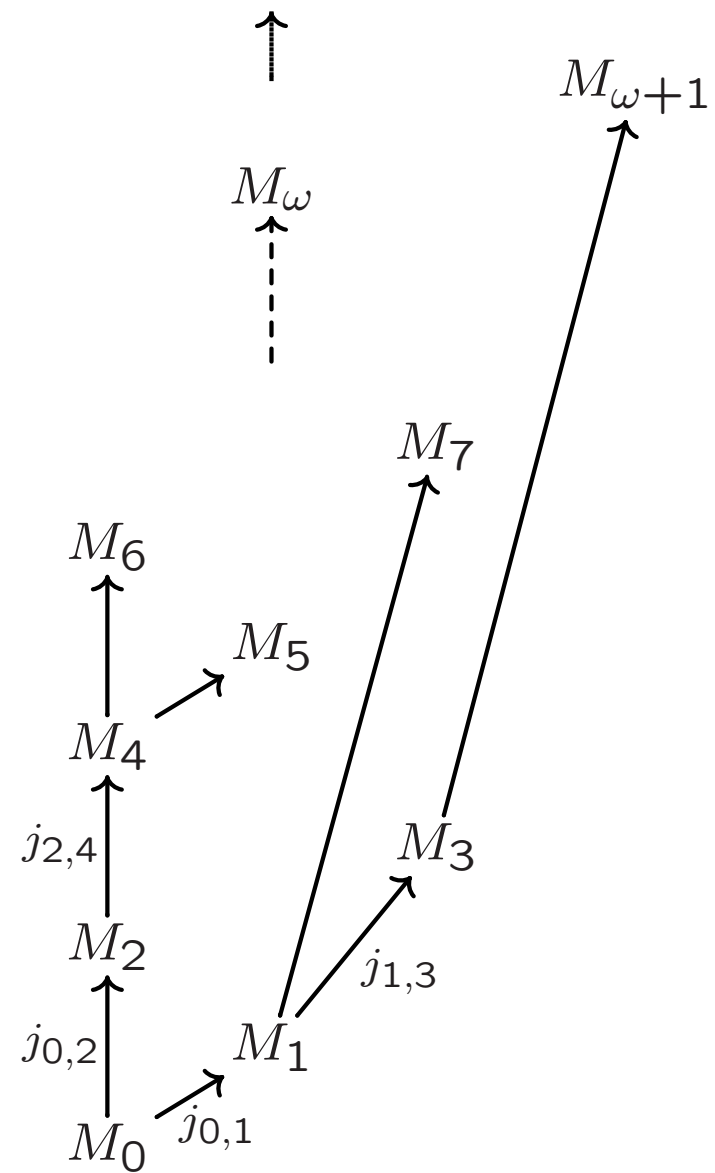
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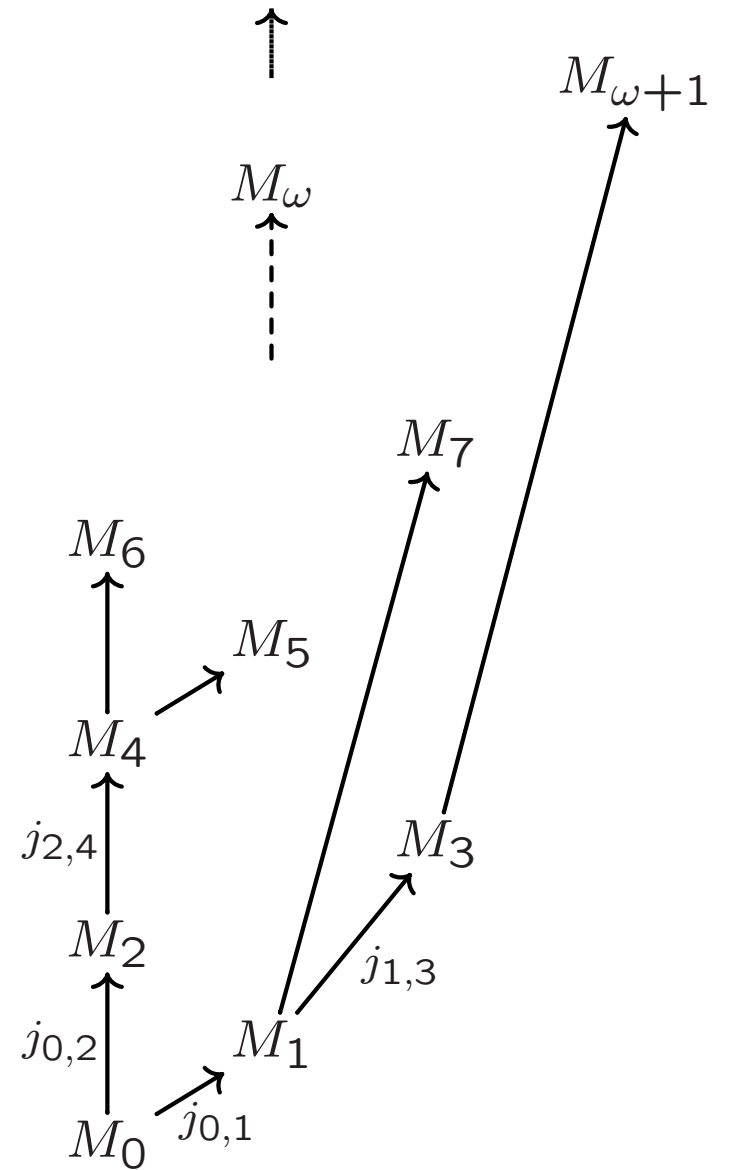
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This allows constructing iterated ultrapowers with non-linear base orders.

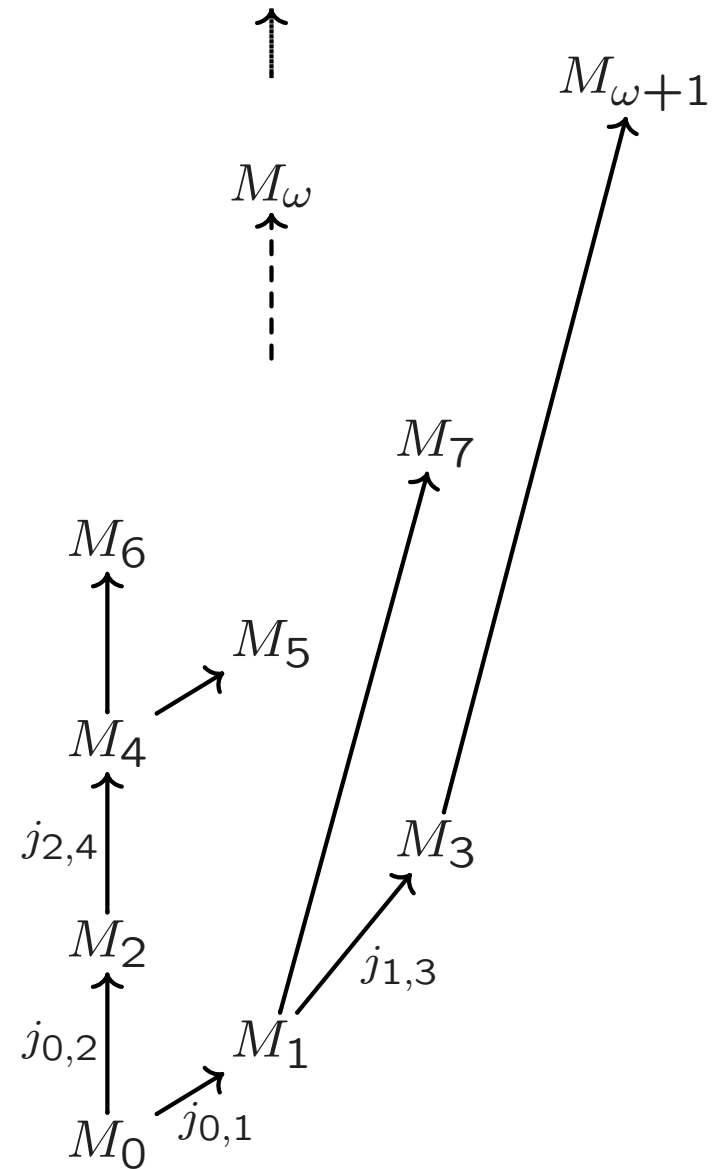


Constructed in stages, starting from a base model  $M = M_0$ .



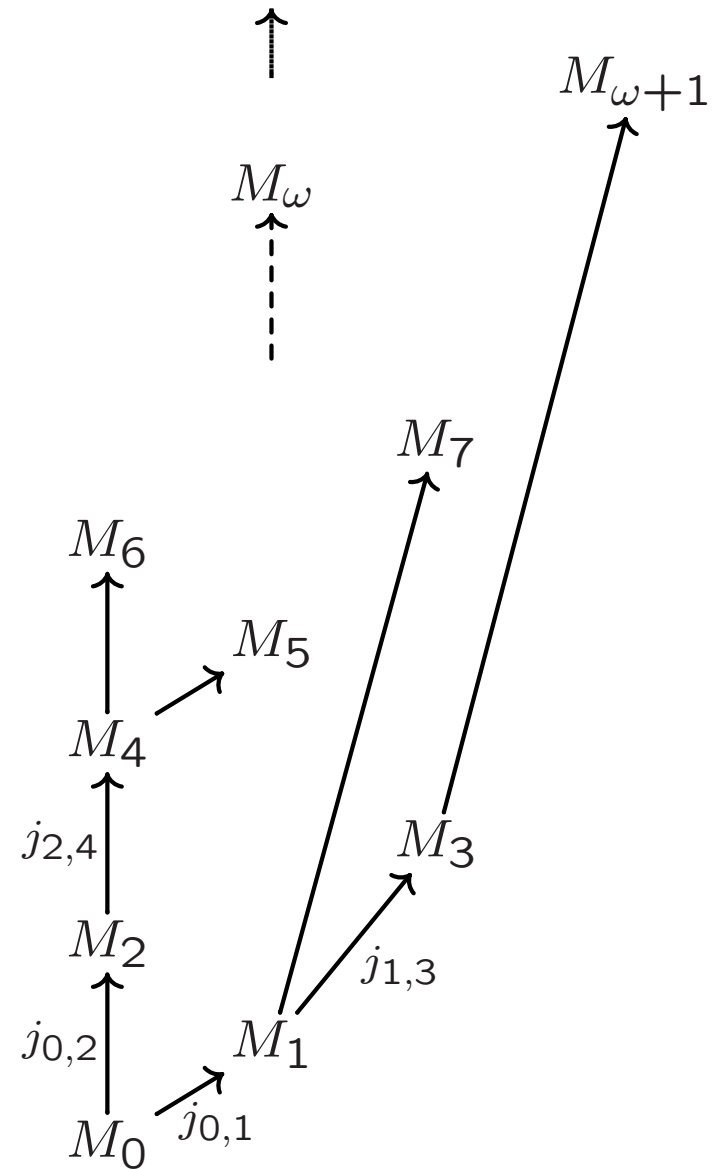
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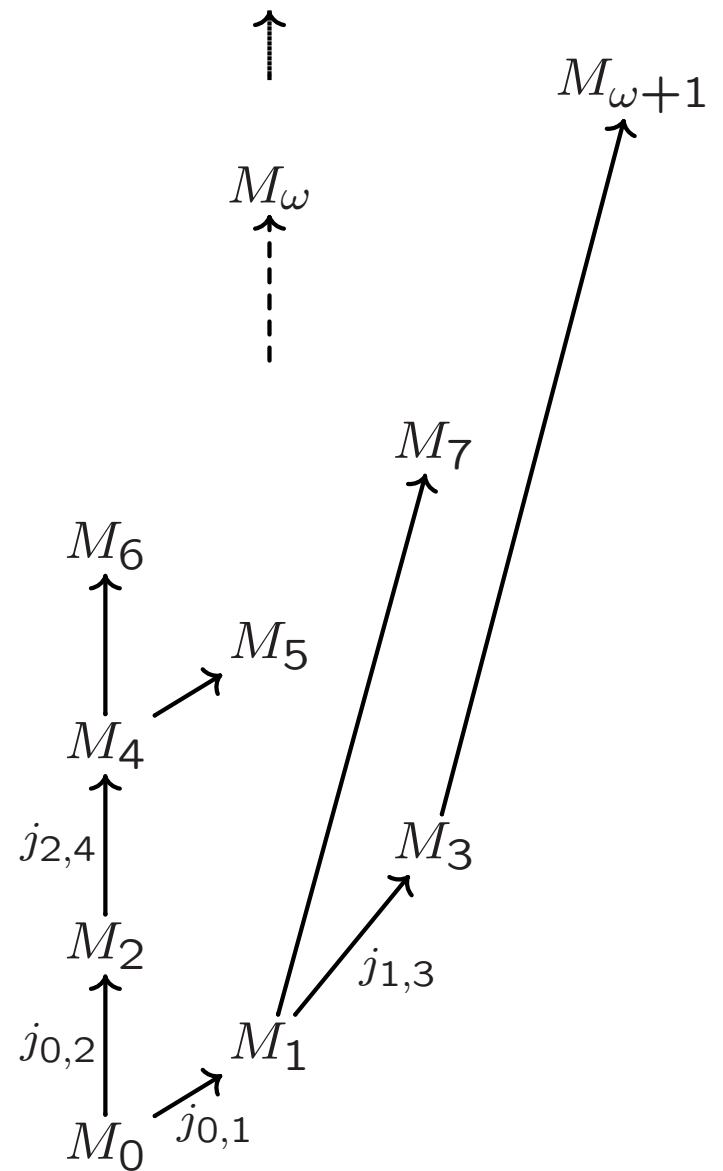
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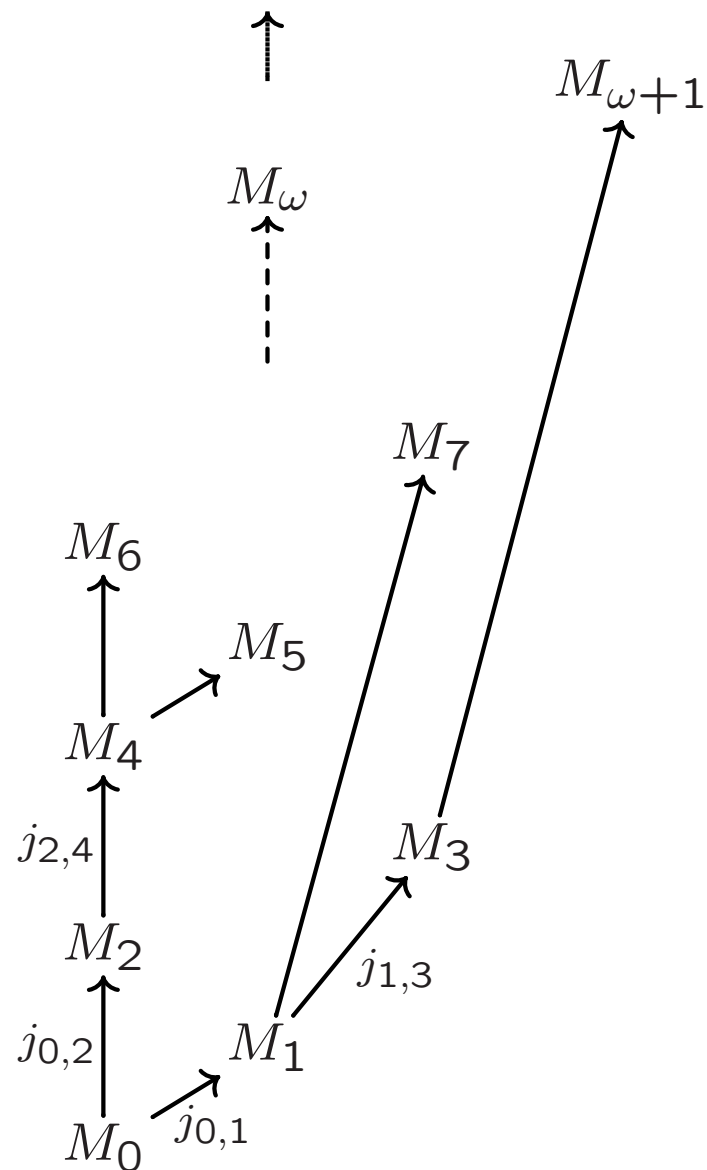
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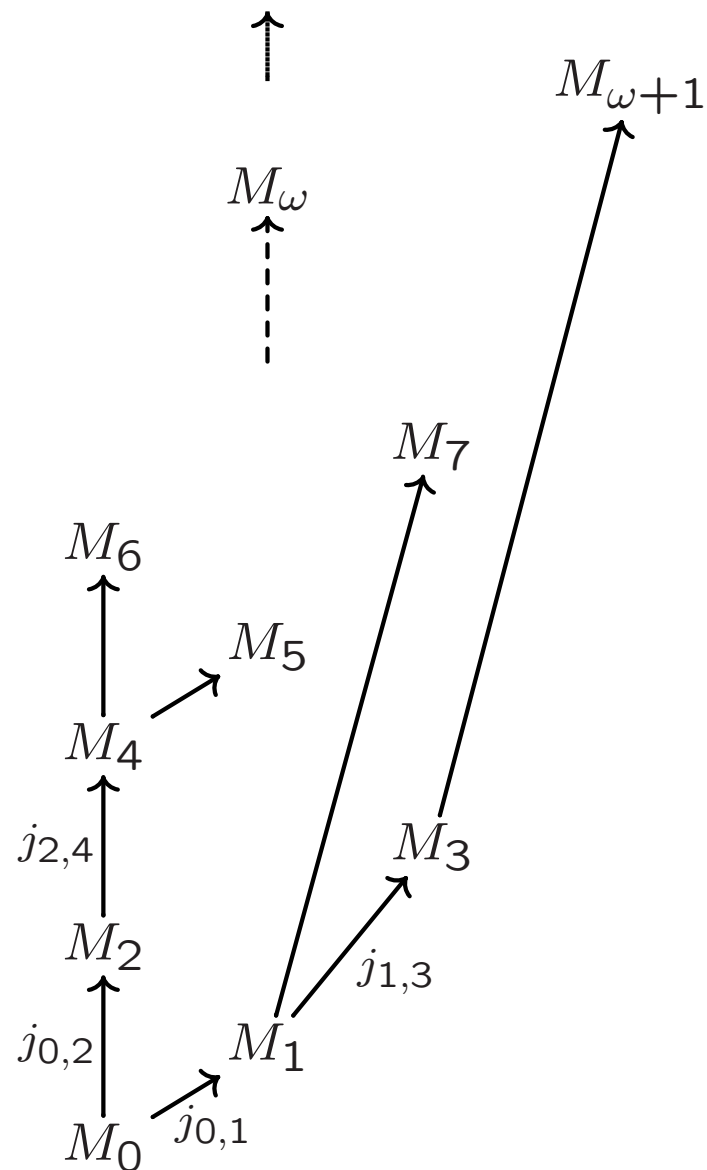
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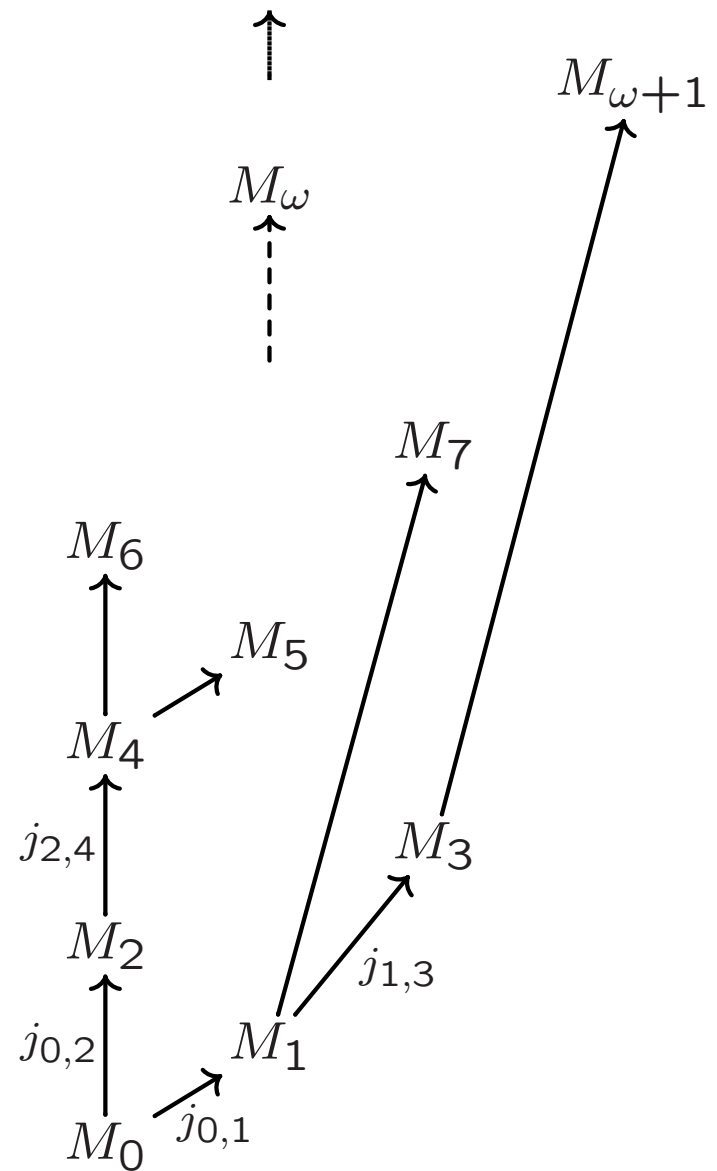




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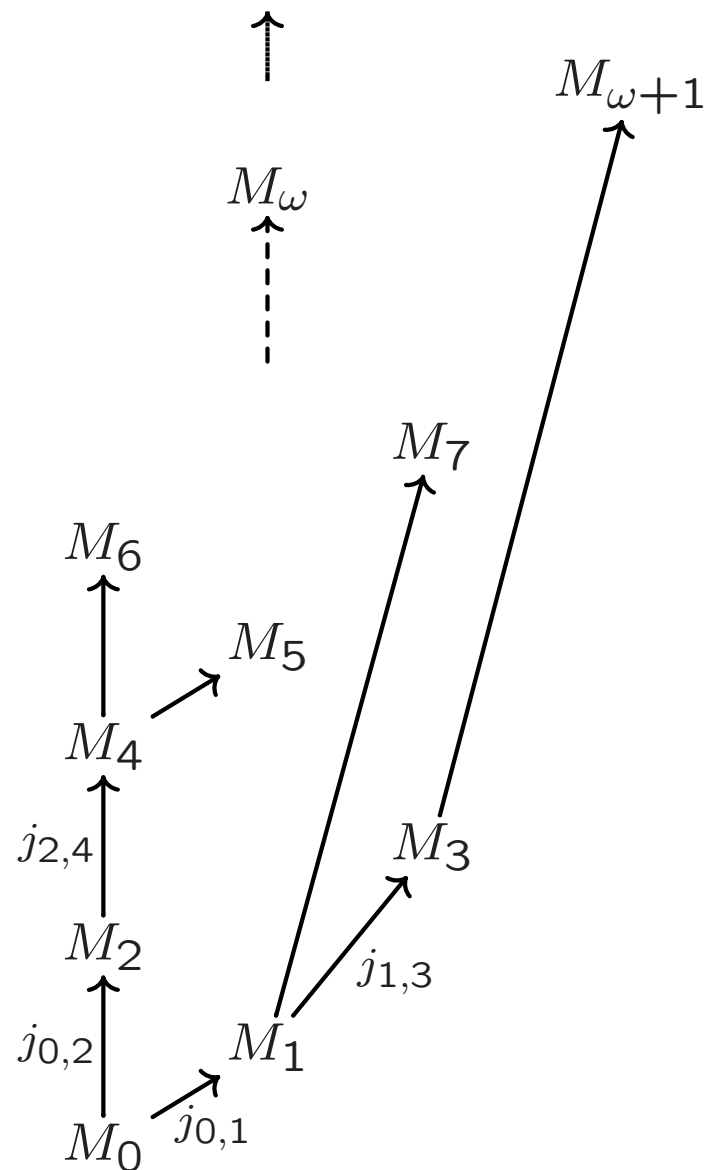
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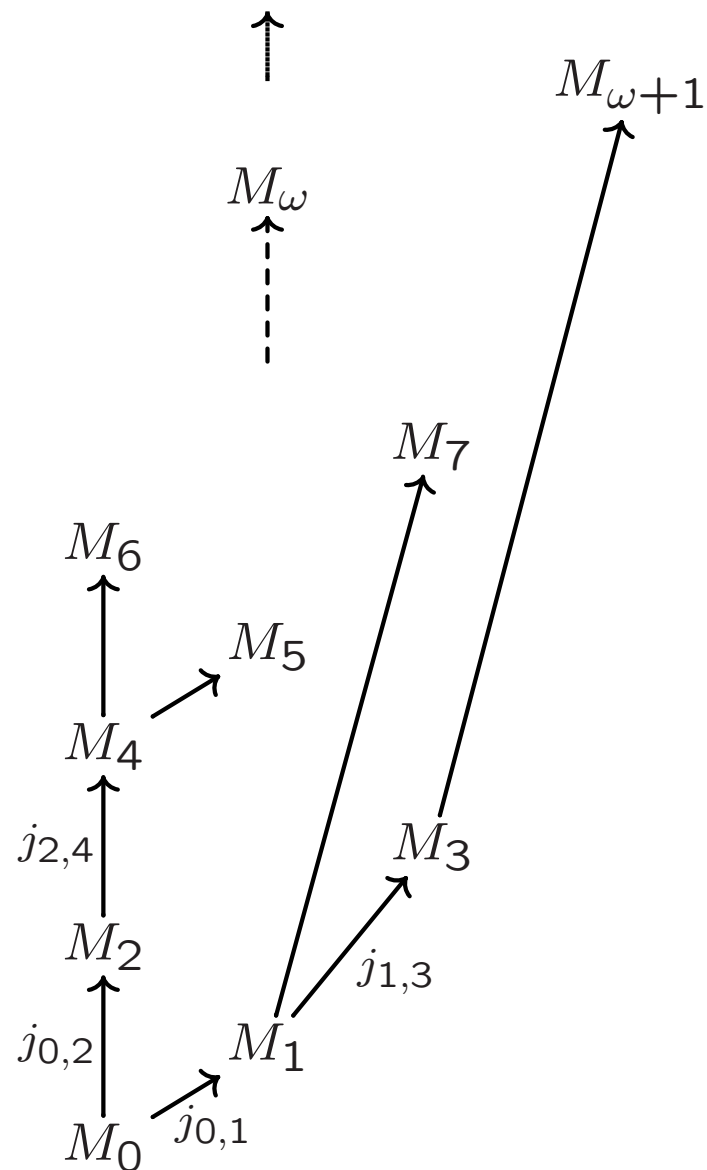


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The result is an *iteration tree* on  $M$ .



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In both cases Woodin cardinals in iterable inner models (rather than the actual universe  $V$ ) are enough, and moreover *necessary*.

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But some make direct use of inner models for Woodin cardinals.

A set  $A$  is  $\alpha$ - $\Pi_1^1$  if there is a sequence  $\langle A_\xi \mid \xi < \alpha \rangle$  of  $\Pi_1^1$  sets so that  $x \in A$  iff the least  $\xi$  so that  $x \notin A_\xi \vee \xi = \alpha$  is odd.

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(The hierarchy generated by this definition is the *difference hierarchy* on  $\Pi_1^1$  sets. If  $\alpha = 2$  for example, then the condition states simply that  $A = A_0 - A_1$ .)

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For  $n = 0$ :  $\Pi_1^1$  determinacy gives a non-trivial  $\pi: L \rightarrow L$  (Harrington), which in turn gives  $<\omega^2\text{-}\Pi_1^1$  determinacy (Martin).

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Theorem known previously for odd  $n$ , not using large cardinals (Kechris–Woodin).

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**Theorem 15 (Neeman, Woodin)** Assume  $AD^L(\mathbb{R})$ . Then it is consistent (with  $AD^L(\mathbb{R})$  and the axiom of choice) that  $\delta_3^1 = \omega_2$ .

Let  $\theta(v)$  be a formula. A *sharp* for  $\theta$  is a non-trivial embedding  $\pi: M \rightarrow M$  where  $M$  is the minimal iterable class model admitting a non-trivial embedding  $\pi$  and satisfying  $\theta[\text{crit}(\pi)]$ .

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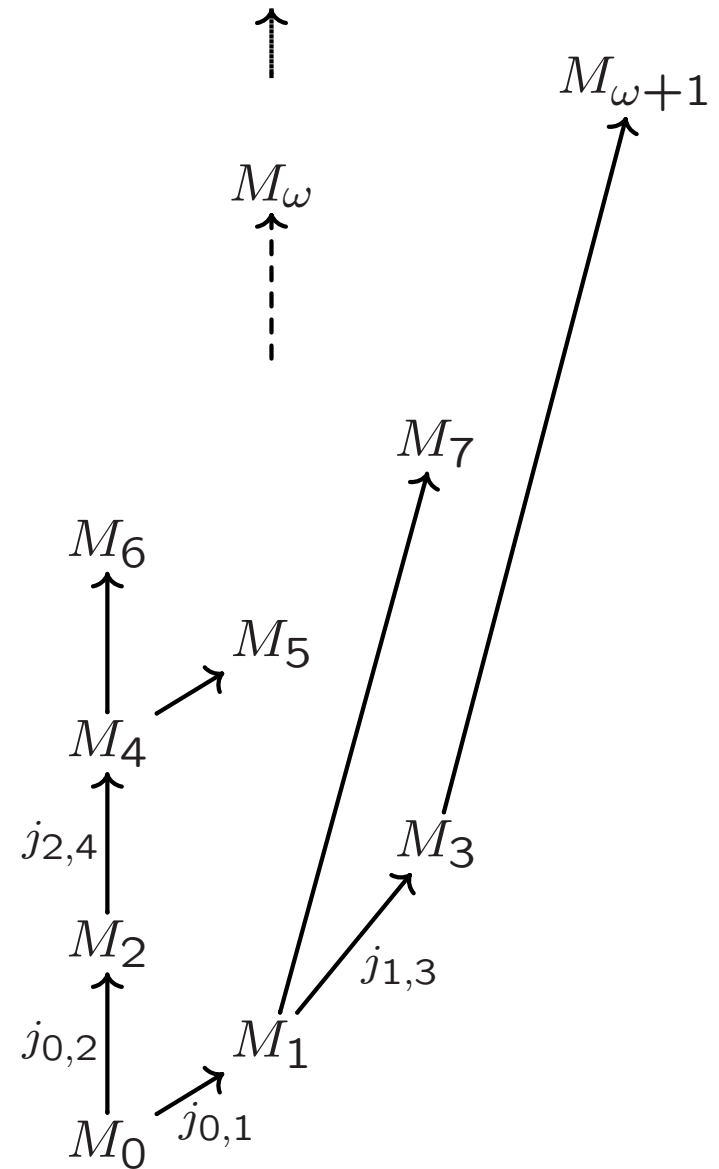
The creation of iteration trees requires some choice at limits.

Constructed in stages, starting from a base model  $M = M_0$ .

E.g., having constructed  $M_1, \dots, M_6$ : pick an extender  $E_6 \in M_6$ , apply it to  $M_1$ , setting  $M_7 = \text{Ult}(M_1, E_6)$  and letting  $j_{1,7}: M_1 \rightarrow M_7$  be the ultrapower embedding.

At limit  $\lambda$ : **pick** a branch through the tree, cofinal in  $\lambda$ . Set  $M_\lambda$  equal to the direct limit of models and embeddings along this branch.

The result is an *iteration tree* on  $M$ .





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$M$  is *iterable* if these choices can be made in a way that secures the wellfoundedness of all the models created.

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Comparisons through iterated ultrapowers show that any two ways to witness  $\theta$  are compatible.



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The connection (with analogues for  $\omega$  Woodin cardinals) is crucial for Theorems 13–15.



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Let  $[\vec{S}]$  denote the set

$$\{ \langle \alpha_0, \dots, \alpha_{k-1} \rangle \in [\omega_1]^{<\omega} \mid (\forall i < k) \alpha_i \in S_{\langle \alpha_0, \dots, \alpha_{i-1} \rangle} \}.$$

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If neither condition holds then both players lose.

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The theorem establishes a precise analogue of Theorems 16 and 17, but for embeddings concentrating on Woodin cardinals and for games of length  $\omega_1$ .



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Games motivated by Theorem 18 were used by Woodin in results on  $\Sigma_2^2$  absoluteness. Other games similar to those in the theorem are enough to capture the theory of superstrong cardinals. But there are no determinacy proofs for these games from large cardinals, and indeed there are some negative results (Larson).

The End

Press **Esc.**