

## Solutions to Exercises 6.6, 6.7, and 6.8

**Exercise 6.6.** The new deductive system would be complete but not sound. It is complete because the old system is complete and every deduction in the old system is a deduction in the new system. To see that the new system is not sound, note that  $(v_1 = v_2 \rightarrow v_1 = v_2)$  is a tautology, and so it is deducible from the empty set in both old and new systems. The new Quantifier Rule allows us to deduce  $(v_1 = v_2 \rightarrow \forall v_2 v_1 = v_2)$  from the empty set. This is clearly not a valid formula.

**Exercise 6.7.** Let  $\mathcal{L}^*$  come from  $\mathcal{L}$  by adding infinitely many new constants. Let  $C$  be the set of all the new constants. Let  $\Sigma$  be

$$\text{Th}(\mathfrak{A}) \cup \{(c \neq c' \wedge c \sim c') \mid c \in C \wedge c' \in C \wedge c \neq c'\}.$$

Let  $\Delta$  be any finite subset of  $\Sigma$ . Extend  $\mathfrak{A}$  to a model for  $\mathcal{L}^*$ . There is an equivalence class of  $\mathfrak{A}$  that has at least as many members as there are members of  $C$  that occur in  $\Delta$ . Extend  $\mathfrak{A}$  to a model for  $\mathcal{L}^*$  by assigning the members of  $C$  that occur in  $\Delta$  to distinct members of this equivalence class and assigning the other members of  $C$  arbitrarily. Clearly  $\Delta$  is true in this model. By Compactness, there is a model  $\mathfrak{B}^*$  in which  $\Sigma$  is true. Let  $\mathfrak{B}$  be the reduct of  $\mathfrak{B}^*$  to  $\mathcal{L}$ .

**Exercise 6.8.** Let  $\mathcal{L}'$  be gotten from  $\mathcal{L}$  by removing the symbol  $P$ . By completeness for our deductive system for  $\mathcal{L}'$ , it is enough to show that  $\Sigma \models \tau$  in  $\mathcal{L}'$ . To show this, let  $\mathfrak{A}'$  be a model for  $\mathcal{L}'$  in which  $\Sigma$  is true. Extend  $\mathfrak{A}'$  to a model  $\mathfrak{A}$  for  $\mathcal{L}$  by assigning some relation as  $P_{\mathfrak{A}}$ . Since  $P$  does not occur in  $\Sigma$ ,  $\Sigma$  is true in  $\mathfrak{A}$ . Since By soundness for our deductive system for  $\mathcal{L}$ ,  $\Sigma \models \tau$  in  $\mathcal{L}$ . Hence  $\tau$  is true in  $\mathfrak{A}$ . Since  $P$  does not occur in  $\tau$ ,  $\tau$  is true in  $\mathfrak{A}'$ .