

# MORE ON $SOP_1$ AND $SOP_2$

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ABSTRACT. This paper continues [Sh500] and [DzSh692]. We present a rank function for  $NSOP_1$  theories and give an example of a theory which is  $NSOP_1$  but not simple. We also investigate the connection between maximality in the ordering  $\triangleleft^*$  among complete first order theories and the  $(N)SOP_2$  property. We prove that  $\triangleleft^*$ -maximality implies  $SOP_2$  and obtain certain results in the other direction. The paper provides a step toward the classification of unstable theories without the strict order property.

## 1. INTRODUCTION AND PRELIMINARIES

We continue the work started by Mirna Džamonja and the first author in [Sh500] and [DzSh692]. The main goal of this project is to throw more light on first order theories with the tree property (that is, non-simple) and without the strict order property (more specifically, without the  $SOP_3$ , see Definition 1.1 below). We pursue and finalize certain directions started in [DzSh692] and answer several questions asked there, providing a more general and complete picture.

The reader may be familiar with a former version of this paper that has been available as a preprint on Shelah's archive (under the number

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“ShUs:E32”) and on Usvyatsov’s webpage, and is referred to in the most recent version of [DžSh692].

Some connections between the work of Džamonja and Shelah and this article are already explained in the introduction of [DžSh692]. In particular, our results provide a generalization of the main theorem of section 1 in [DžSh692], expand the results of section 2 there and answer certain questions which were left open in section 3. One of the answers leads to a complete proof of a theorem which had been the original motivation of section 3 of [DžSh692], Corollary 3.15 here (see also Discussion 3.12 here, Theorem 0.5 in [DžSh692] and discussion preceding it). We give more details below.

Before describing the background and the results obtained in this paper, let us recall the definitions of  $SOP_n$  hierarchy, starting with the more classical concepts introduced in [Sh500].

Let  $T$  be a complete first order theory,  $\mathfrak{C}$  - the monster model of  $T$  (a  $\kappa^*$  - saturated model for  $\kappa^*$  big enough).

**Definition 1.1.** (1) Let  $n \geq 3$ . We say  $\varphi(\bar{x}, \bar{y})$  (with  $\text{lg}(\bar{x}) = \text{lg}(\bar{y})$ ) exemplifies the strong order property of order  $n$  ( $SOP_n$ ) in  $T$  if it defines on  $\mathfrak{C}$  a directed graph with infinite indiscernible chains and no cycles of length  $n$ .

(2) We say  $\varphi(\bar{x}, \bar{y})$  (with  $\text{lg}(\bar{x}) = \text{lg}(\bar{y})$ ) exemplifies the *strict order property* in  $T$  if it defines on  $\mathfrak{C}$  a partial order with infinite indiscernible chains.

**Fact 1.2.** For a theory  $T$ , *strict order property*  $\implies SOP_{n+1} \implies SOP_n$  for all  $n \geq 3$ .

*Proof.* The first implication is trivial, for the other one see [Sh500], Claim (2.6).  $\square$

We also remind the reader the following equivalent definition of  $SOP_3$ :

**Fact 1.3.** *T has  $SOP_3$  if and only if there is an indiscernible sequence  $\langle \bar{a}_i : i < \omega \rangle$  and formulae  $\varphi(\bar{x}, \bar{y})$ ,  $\psi(\bar{x}, \bar{y})$  such that*

- (a)  $\{\varphi(\bar{x}, \bar{y}), \psi(\bar{x}, \bar{y})\}$  is contradictory,
- (b) for some sequence  $\langle \bar{b}_j : j < \omega \rangle$  we have

$$i \leq j \implies \models \varphi[\bar{b}_j, \bar{a}_i] \text{ and } i > j \implies \models \psi[\bar{b}_j, \bar{a}_i].$$

- (c) for  $i < j$ , the set  $\{\varphi(\bar{x}, \bar{a}_j), \psi(\bar{x}, \bar{a}_i)\}$  is contradictory.

*Proof.* Easy, or see [Sh500], Claim (2.20).  $\square$

*Remark 1.4.* Note that if in the previous definition  $\psi = \neg\varphi$ , we get the strict order property.

Now we recall the definitions of  $SOP_1$ ,  $SOP_2$  and related properties:

**Definition 1.5.** (1) *T has  $SOP_2$  if there is a formula  $\varphi(\bar{x}, \bar{y})$  which exemplifies this property in  $\mathfrak{C}$ , and this means:*

There are  $\bar{a}_\eta \in \mathfrak{C}$  for  $\eta \in {}^\omega 2$  such that

- (a) For every  $\eta \in {}^\omega 2$ , the set  $\{\varphi(\bar{x}, \bar{a}_{\eta \upharpoonright l}) : l < \omega\}$  is consistent.
- (b) If  $\eta, \nu \in {}^\omega 2$  are incomparable,  $\{\varphi(\bar{x}, \bar{a}_\eta), \varphi(\bar{x}, \bar{a}_\nu)\}$  is inconsistent.

(2) *T has  $SOP_1$  if there is a formula  $\varphi(\bar{x}, \bar{y})$  which exemplifies this in  $\mathfrak{C}$ , which means:*

There are  $\bar{a}_\eta \in \mathfrak{C}$ , for  $\eta \in {}^\omega 2$  such that:

- (a) for  $\rho \in {}^\omega 2$  the set  $\{\varphi(\bar{x}, \bar{a}_{\rho \upharpoonright n}) : n < \omega\}$  is consistent.
- (b) if  $\nu \frown \langle 0 \rangle \trianglelefteq \eta \in {}^\omega 2$ , then  $\{\varphi(\bar{x}, \bar{a}_\eta), \varphi(\bar{x}, \bar{a}_{\nu \frown \langle 1 \rangle})\}$  is inconsistent.

(3)  $NSOP_2$  and  $NSOP_1$  are the negations of  $SOP_2$  and  $SOP_1$  respectively.

(4)  $T$  has  $SOP'_1$  if there is a formula  $\varphi(\bar{x}, \bar{y})$  which exemplifies this property in  $\mathfrak{C}$ , and this means:

there are  $\langle \bar{a}_\eta : \eta \in {}^\omega 2 \rangle$  in  $\mathfrak{C}_T$  such that

- (a)  $\{\varphi(\bar{x}, \bar{a}_{\eta \upharpoonright n})^{\eta(n)} : n < \omega\}$  is consistent for every  $\eta \in {}^\omega 2$ , where we use the notation

$$\varphi^l = \begin{cases} \varphi & \text{if } l = 1, \\ \neg\varphi & \text{if } l = 0 \end{cases}$$

for  $l < 2$ .

- (b) If  $\nu \frown \langle 0 \rangle \trianglelefteq \eta \in {}^\omega 2$ , then  $\{\varphi(\bar{x}, \bar{a}_\eta), \varphi(\bar{x}, \bar{a}_\nu)\}$  is inconsistent.

(5)  $T$  has  $SOP''_2$  if there is a formula  $\varphi(\bar{x}, \bar{y})$  which exemplifies this property in  $\mathfrak{C}$ , and this means:

there is  $n < \omega$  and a sequence

$$\langle \bar{a}_{\bar{\eta}} : \bar{\eta} = \langle \eta_0, \dots, \eta_{n-1} \rangle, \eta_0 \triangleleft \eta_1 \triangleleft \dots \triangleleft \eta_{n-1} \in {}^{\lambda > 2} \text{ and } \text{lg}(\eta_i) \text{ successor} \rangle$$

such that

(a) for each  $\eta \in {}^\lambda 2$ , the set

$$\left\{ \begin{array}{l} \varphi(\bar{x}, \bar{a}_{\bar{\eta}}) : \bar{\eta} = \langle \eta \upharpoonright (\alpha_0 + 1), \eta \upharpoonright (\alpha_1 + 1), \dots, \eta \upharpoonright (\alpha_{n-1} + 1) \rangle \\ \text{and } \alpha_0 < \alpha_1 < \dots < \alpha_{n-1} < \lambda \end{array} \right\}$$

is consistent

(b) for every large enough  $m$ , if  $h$  is a 1-to-1 function from  ${}^{n \geq m}$  into  ${}^{\lambda > 2}$  preserving  $\eta \triangleleft \nu$  and  $\eta \perp \nu$  (incomparability) then  $\{\varphi(\bar{x}, \bar{a}_{\bar{\nu}}) : \text{for some } \eta \in {}^n m \text{ we have } \bar{\nu} = \langle h(\eta \upharpoonright \ell) : \ell \leq n \rangle\}$  is inconsistent.

**Fact 1.6.** (1) For a theory  $T$ ,  $SOP_3 \implies SOP_2 \implies SOP_1$   
 (2)  $T$  has  $SOP_1$  if and only if it has  $SOP'_1$

*Proof.* See [DžSh692]. □

It is still not known whether the implications in 1.6(1) are strict, but for now we investigate each one of these order properties on its own.

In the second section we expand our knowledge on  $SOP_1$ . We present a rank function measuring type-definable “squares”, i.e. pairs of types of the form  $(p(\bar{x}), q(\bar{y}))$  and show the rank is finite for every such a pair if and only if  $T$  does not have  $SOP'_1$  (if and only if  $T$  does not have  $SOP_1$ ). In fact, if one calls a tree of parameters  $\{\bar{a}_\eta : \eta \in {}^{>2}\}$  showing that  $\varphi(\bar{x}, \bar{y})$  exemplifies  $SOP'_1$  in  $\mathfrak{C}$  (as in the definition of  $SOP'_1$ ) a  $\varphi - SOP'_1$  tree, the rank measures exactly the maximal depth of a tree like this that can be built in  $\mathfrak{C}$ . We also show a small application of the rank.

It is easy to see (see [DžSh692]) that if  $\varphi(\bar{x}, \bar{y})$  exemplifies  $SOP_1$  in  $\mathfrak{C}$  then it also exemplifies the tree property, so  $T$  has  $SOP_1 \implies T$  is not simple. We show that the implication is proper, i.e. find an example of a theory  $T$  which is not simple, but is  $NSOP_1$ . This theory which we call  $T_{\text{feq}}^*$ , was first defined in [Sh457], and is used in [Sh500] as an example of an  $NSOP_3$  non-simple theory. Here we use a slightly different definition of the same theory, as given in [DžSh692].

**Definition 1.7.** (1)  $T_{\text{feq}}$  is the following theory in the language

$\{Q, P, E, R, F\}$

- (a) Predicates  $P$  and  $Q$  are unary and disjoint, and  $(\forall x) [P(x) \vee Q(x)]$ ,
- (b)  $E$  is an equivalence relation on  $Q$ ,
- (c)  $R$  is a binary relation on  $Q \times P$  such that

$$[x R z \ \& \ y R z \ \& \ x E y] \implies x = y.$$

(so  $R$  picks for each  $z \in Q$  (at most one) representative of any  $E$ -equivalence class).

- (d)  $F$  is a (partial) binary function from  $Q \times P$  to  $Q$ , which satisfies

$$F(x, z) \in Q \ \& \ (F(x, z)) R z \ \& \ x E (F(x, z)).$$

(so for  $x \in Q$  and  $z \in P$ , the function  $F$  picks the representative of the  $E$ -equivalence class of  $x$  which is in the relation  $R$  with  $z$ ).

- (2)  $T_{\text{feq}}^*$  is the model completion of  $T_{\text{feq}}$ .

If the reader thinks about the definition above, they will find out that  $T_{\text{feq}}^*$  is just the model completion of the theory of infinitely many (independent) parameterized equivalence relations. The reader can also compare between the definition of  $T_{\text{feq}}^*$  here and in [Sh457]. As we have already mentioned, it was shown in [Sh500] that this theory does not have  $SOP_3$  (but is not simple). Here we prove an (a priori) stronger result:  $T_{\text{feq}}^*$  does not have  $SOP_1$ .

In the third section we deal with  $\triangleleft_\lambda^*$ -maximality (see the beginning of the section for definitions). For a theory  $T$  to be  $\triangleleft_\lambda^*$ -maximal means to be complicated. In a way, it means that it is hard to make its models  $\lambda$ -saturated.

The motivation for considering this property comes from Classification Theory and the search for “dividing lines”. The authors believe that a “good” property of a theory  $T$  should have several characterizations of different types, both “internal” (something happens in the monster model of  $T$ , such as  $SOP_n$ ) and “external” (how  $T$  compares to other theories, such as maximality in a certain order). Although all the approximations to the strict order property (including  $SOP_n$  and the strict order property itself) seem very natural syntactic internal definitions, no external property is known to characterize any of them. There are natural conjectures, though. The following question was partly guiding our current work:

*Question 1.8.* Does  $\triangleleft_\lambda^*$ -maximality characterize either  $SOP_3$  or  $SOP_2$ , maybe both?

There are several indications that the answer should be positive. It had been already known before our work that  $\triangleleft_\lambda^*$ -maximality lies strictly “above” the tree property (non-simplicity): Džamonja and Shelah showed in [DžSh692] that  $T_{feq}^*$  (which is not simple) is not  $\triangleleft_\lambda^*$ -maximal. The question where exactly above non-simplicity this property lies is still open, but we narrow the possibilities down significantly. It follows from our results here that  $SOP_3 \implies \triangleleft_\lambda^*$ -maximality  $\implies SOP_2$ . We also obtain a local version of the reversed direction of the second implication.

Our analysis also provides an alternative proof of the fact that  $T_{feq}^*$  is not  $\triangleleft_\lambda^*$ -maximal, Theorem 1.17 in [DžSh692]: no  $NSOP_2$  theory is  $\triangleleft_\lambda^*$ -maximal, and  $T_{feq}^*$  is  $NSOP_1$ , therefore  $NSOP_2$ . So by bringing the “internal” and the “external” dividing lines close together, we also give many examples of non-simple theories which are not  $\triangleleft_\lambda^*$ -maximal,  $T_{feq}^*$  being a particular case. See also Discussion 3.18.

Let us now give more details concerning some results in the paper and explain how exactly they fit in the general picture. In [Sh500] it was stated that  $SOP_3$  implies  $\triangleleft_\lambda^*$ -maximality, but the proof there is not complete: it is shown that every theory with  $SOP_3$  is  $\triangleleft_\lambda^*$ -above  $T_{tr}^*$ , the model completion of the theory of trees. The first theorem in section 3, Theorem 3.5, fills the missing part, showing explicitly that  $T_{tr}^*$  is  $\triangleleft_\lambda^*$ -maximal for every  $\lambda > \aleph_0$ . This also continues [Sh:c], chapter VI, where Keisler’s order, a relative of  $\triangleleft_\lambda^*$ , is studied.

One of the reason for giving an explicit proof for Theorem 3.5 here was to provide more tools for strengthening the result above to  $SOP_2$

theories, i.e. showing that  $SOP_2 \implies \triangleleft_\lambda^*$ -maximality. A step in this direction is Theorem 3.11 where we show a “local” version: if a formula  $\vartheta$  exemplifies  $SOP_2$  in  $T$ , then the pair  $(T, \vartheta)$  is  $\triangleleft_\lambda^*$ -above the pair  $(T_{tr}^*, y < x)$  for every regular  $\lambda > |T|$  (again, see the beginning of section 3 for precise definitions). This result, although interesting on its own, is insufficient for “global”  $\triangleleft_\lambda^*$ -maximality of  $SOP_2$  theories, as explained in Discussion 3.17. Nevertheless, combined with Theorem 3.5 and its proof, it gives more information on the behavior of  $SOP_2$  theories and  $\triangleleft_\lambda^*$ -order altogether.

As for the other direction ( $\triangleleft_\lambda^*$ -maximality  $\implies SOP_2$ ), we provide a complete proof, based on several related results achieved by Džamonja and the first author, who showed in [DžSh692] that a property similar to  $\triangleleft_\lambda^*$ -maximality (which also follows from  $\triangleleft_\lambda^*$ -maximality for some  $\lambda$  under certain set theoretic conditions) implies  $SOP_2''$ . One of the questions left open in [DžSh692] is the connection between  $SOP_2''$  and the  $SOP_n$  hierarchy. Of course, it would be natural to connect between  $SOP_2''$  and  $SOP_2$ , and indeed we prove here that these two properties are equivalent for a theory  $T$  (not necessarily for a formula), Theorem 3.13.

So we can conclude  $SOP_3 \implies \triangleleft_\lambda^*$ -maximality  $\implies SOP_2$ , while very little is known at this point concerning implications in the other directions.

The following definitions and facts are going to be very useful.

In [DžSh692] two notions of “tree indiscernibility” were defined. We recall the definitions:

**Definition 1.9.** (1) Given an ordinal  $\alpha$  and sequences  $\bar{\eta}_l = \langle \eta_0^l, \eta_1^l, \dots, \eta_{n_l}^l \rangle$

for  $l = 0, 1$  of members of  ${}^{\alpha>2}$ , we say that  $\bar{\eta}_0 \approx_1 \bar{\eta}_1$  iff

- (a)  $n_0 = n_1$ ,
- (b) the truth values of

$$\eta_{k_3}^l \trianglelefteq \eta_{k_1}^l \cap \eta_{k_2}^l, \quad \eta_{k_1}^l \cap \eta_{k_2}^l \triangleleft \eta_{k_3}^l, \quad (\eta_{k_1}^l \cap \eta_{k_2}^l) \frown \langle 0 \rangle \trianglelefteq \eta_{k_3}^l,$$

for  $k_1, k_2, k_3 \leq n_0$ , do not depend on  $l$ .

- (2) We say that the sequence  $\langle \bar{a}_\eta : \eta \in {}^{\alpha>2} \rangle$  of  $\mathfrak{C}$  (for an ordinal  $\alpha$ ) are *1-fully binary tree indiscernible (1-fbti)* iff whenever  $\bar{\eta}_0 \approx_1 \bar{\eta}_1$  are sequences of elements of  ${}^{\alpha>2}$ , then

$$\bar{a}_{\bar{\eta}_0} =: \bar{a}_{\eta_0^0} \frown \dots \frown \bar{a}_{\eta_{n_0}^0}$$

and the similarly defined  $\bar{a}_{\bar{\eta}_1}$ , realize the same type in  $\mathfrak{C}$ .

- (3) We replace 1 by 2 in the above definition iff  $(\eta_{k_1}^l \cap \eta_{k_2}^l) \frown \langle 0 \rangle \trianglelefteq \eta_{k_3}^l$  is omitted from clause (b) above.

We will need the following fact proved in [DžSh692], (2.11):

**Fact 1.10.** *If  $t \in \{1, 2\}$  and  $\langle \bar{b}_\eta : \eta \in {}^{\omega>2} \rangle$  are given, and  $\delta \geq \omega$ , then we can find  $\langle \bar{a}_\eta : \eta \in {}^{\delta>2} \rangle$  such that*

- (a)  $\langle \bar{a}_\eta : \eta \in {}^{\delta>2} \rangle$  is *t-fbti*,
- (b) if  $\bar{\eta} = \langle \eta_m : m < n \rangle$ , where each  $\eta_m \in {}^{\delta>2}$  is given, and  $\Delta$  is a finite set of formulae of  $T$ , then we can find  $\nu_m \in {}^{\omega>2}$  ( $m < n$ ) such that with  $\bar{\nu} =: \langle \nu_m : m < n \rangle$ , we have  $\bar{\nu} \approx_t \bar{\eta}$  and the sequences  $\bar{a}_{\bar{\eta}}$  and  $\bar{b}_{\bar{\nu}}$ , realize the same  $\Delta$ -types.

**Convention 1.11.** We work with a complete first order theory  $T$ , let  $\mathfrak{C}$  be its “monster” model (saturated in some very big  $\kappa^*$ ). Let  $\mathcal{L} = \mathcal{L}(T)$  (the language of  $T$ ). Every formula we mention is an  $\mathcal{L}$ -formula, maybe with parameters from  $\mathfrak{C}$ .

## 2. MORE ON $SOP_1$

$SOP_1$  was introduced by Džamonja and Shelah as an intermediate property between simplicity and  $SOP_3$ . A natural question is: does the class  $NSOP_1$  coincide with either simple theories or  $NSOP_3$ ? Here we give a negative answer to the first question above. The answer for the second one is still not known.

**Theorem 2.1.**  $T_{feq}^*$  does not have  $SOP_1$ .

*Proof.* Suppose there exists  $\varphi(\bar{x}, \bar{y})$  with  $lg(\bar{x}) = n$ ,  $lg(\bar{y}) = m$ , and  $\langle \bar{a}_\eta : \eta \in {}^\omega > 2 \rangle$  in  ${}^m \mathfrak{C}$  which exemplify  $SOP_1$  in  $\mathfrak{C}$  ( $\mathfrak{C}$  is the monster model of  $T_{feq}^*$ ). Without loss of generality, (by Fact 1.10)  $\langle \bar{a}_\eta : \eta \in {}^\omega > 2 \rangle$  is 1-full tree indiscernible. Also, by elimination of quantifiers, we may assume that  $\varphi(\bar{x}, \bar{y})$  is quantifier free. As the only function symbol in the language is  $F$  and  $F^\mathfrak{C}$  has the property  $F^\mathfrak{C}(F^\mathfrak{C}(x, z), y) = F^\mathfrak{C}(x, y)$  for all  $z$ , we will also assume wlog that  $\bar{x}$  and  $\bar{y}$  in  $\varphi(\bar{x}, \bar{y})$  are closed under  $F$  and  $\varphi(\bar{x}, \bar{y})$  gives the full diagram of  $\bar{x} \frown \bar{y}$ . We shall regard  $\bar{x}$  as  $\langle x^0, \dots, x^{n-1} \rangle$ ,  $\bar{y}$  as  $\langle y^0, \dots, y^{m-1} \rangle$ ,  $\bar{a}_\eta$  as  $\langle a_\eta^0, \dots, a_\eta^{m-1} \rangle$ .

By the definition of  $SOP_1$ , there exist  $\bar{e} = \langle e^0, \dots, e^{n-1} \rangle$ ,  $\bar{d} = \langle d^0, \dots, d^{n-1} \rangle$  in  ${}^n \mathfrak{C}$  s.t.

$$\mathfrak{C} \models \varphi(\bar{e}, \bar{a}_{(\cdot)}) \wedge \varphi(\bar{e}, \bar{a}_{(0)}) \wedge \varphi(\bar{e}, \bar{a}_{(00)})$$

and

$$\mathfrak{C} \models \varphi(\bar{d}, \bar{a}_{\langle \cdot \rangle}) \wedge \varphi(\bar{d}, \bar{a}_{\langle 1 \rangle})$$

Denote  $\eta = \langle 00 \rangle$ . Let  $B = \mathfrak{C} \upharpoonright \bar{a}_\eta \frown \bar{a}_{\langle 1 \rangle}$ . By our assumptions, there exists a model  $N_0$  whose universe is  $\bar{x} \frown \bar{a}_\eta$ , extending  $\mathfrak{C} \upharpoonright \bar{a}_\eta$ , whose basic diagram is  $\varphi(\bar{x}, \bar{a}_\eta)$ . Similarly, there exists a model  $N_1$  with universe  $\bar{x} \frown \bar{a}_{\langle 1 \rangle}$  and basic diagram  $\varphi(\bar{x}, \bar{a}_{\langle 1 \rangle})$ . We shall amalgamate  $B, N_0$  and  $N_1$  into a model of  $T_{feq}, N$ . This will immediately give a contradiction: first, extend  $N$  to  $N^* \models T_{feq}^*$ , then amalgamate  $N^*$  and  $\mathfrak{C}$  over  $B$  into some  $\mathfrak{C}^+ \models T_{feq}^*$ . By model completeness of  $T_{feq}^*$ ,  $\mathfrak{C} \prec \mathfrak{C}^+$ , but  $\mathfrak{C}^+ \models \exists \bar{x}(\varphi(\bar{x}, \bar{a}_\eta) \wedge \varphi(\bar{x}, \bar{a}_{\langle 1 \rangle}))$ , which is a contradiction to the definition of  $SOP_1$ .

It is left, therefore, to show that we can define on  $|N_0| \cup |N_1|$  a structure which will be a model of  $T_{feq}$ , extending  $B$ .

We define  $N$  as follows:

$$|N| = |N_1| \cup |N_2|, \quad P^N = P^{N_1} \cup P^{N_2}, \quad Q^N = Q^{N_1} \cup Q^{N_2}.$$

Note that the diagram of  $\bar{x}$  in  $N_0$  is the same as the diagram of  $\bar{x}$  in  $N_1$  (both implied by  $\varphi(\bar{x}, \bar{y})$ ), and the diagrams of  $\bar{a}_\eta, \bar{a}_{\langle 1 \rangle}$  in  $N_i$  are the same as in  $\mathfrak{C}$ , hence the same as in  $B$ . Therefore,  $P^N$  and  $Q^N$  are well defined and give a partition of  $|N|$ . Also, so far  $N$  extends  $B$  (as a structure).

Considering  $E$  and  $R$ , we define

$$R^N = R^{N_1} \cup R^{N_2} \cup R^B$$

$$E^N = E^{N_1} \cup E^{N_2} \cup E^B$$

Once we have proven the following claims, we will be able to define  $F^N$  in a natural way, and in fact will be done.

*Claim 2.1.1.*  $E^N$  is an equivalence relation on  $Q^N$ , extending  $E^B$ .

*Claim 2.1.2.*  $R^N$  is a two-place relation on  $N$ ,  $R^N \subseteq P^N \times Q^N$ , satisfying:

for every  $y \in P^N$  and every equivalence class  $C$  of  $E^N$ , there exists a unique  $z \in C$  such that  $(y, z) \in R^N$ .

*Proof of 2.1.1.* The only nonobvious thing is transitivity. We check two main cases, all the rest are either similar or trivial.

- (1) Assume  $x^i E^N a_\eta^j$ ,  $x^i E^N a_{\langle 1 \rangle}^k$  for some  $i, j, k$ . We want to show  $a_\eta^j E^N a_{\langle 1 \rangle}^k$ . It is enough to see  $a_\eta^j E^\mathfrak{C} a_{\langle 1 \rangle}^k$ . We will write  $E$  instead of  $E^\mathfrak{C}$ .

$N \models x^i E a_\eta^j \Rightarrow N_0 \models x^i E a_\eta^j \Rightarrow \varphi(\bar{x}, \bar{y}) \vdash x^i E y^j$ . Similarly,  $\varphi(\bar{x}, \bar{y}) \vdash x^i E y^k$ , and we get (by the choice of  $\bar{e}, \bar{d} \in {}^n \mathfrak{C}$ )  $e^i E a_\eta^j$ ,  $e^i E a_{\langle 1 \rangle}^j$ ,  $e^i E a_\eta^k$ ,  $e^i E a_{\langle 1 \rangle}^k$ ,  $d^i E a_{\langle 1 \rangle}^j$ ,  $d^i E a_{\langle 1 \rangle}^j$ ,  $d^i E a_{\langle 1 \rangle}^k$ ,  $d^i E a_{\langle 1 \rangle}^k$ . Now it is easy to see that all the above elements are  $E$ -equivalent in  $\mathfrak{C}$ , in particular  $a_\eta^j$  and  $a_{\langle 1 \rangle}^k$ , as required.

- (2) Assume  $x^i E^N a_\eta^n$ ,  $a_{\langle 1 \rangle}^k E^N a_\eta^n$ , and we show  $x^i E^N a_{\langle 1 \rangle}^k$ , i.e.  $\varphi(\bar{x}, \bar{y}) \vdash x^i E y^k$ . As  $\varphi(\bar{d}, \bar{a}_{\langle 1 \rangle})$  holds in  $\mathfrak{C}$  and  $\varphi(\bar{x}, \bar{y})$  gives a full diagram, it will be enough to see  $d^i E a_{\langle 1 \rangle}^k$ .

We know that  $\varphi(\bar{x}, \bar{y}) \vdash x^i E y^j$  therefore  $e^i E a_\eta^j$ ,  $e^i E a_{\langle 1 \rangle}^j$ ,  $d^i E a_{\langle 1 \rangle}^j$ ,  $d^i E a_{\langle 1 \rangle}^j$ . In particular,  $d^i E a_\eta^j$ , but, by our assumption,  $a_\eta^j E a_{\langle 1 \rangle}^k$ , so we are done.

□<sub>1</sub>

*Proof of 2.1.2.* Like in the previous lemma, the only nontrivial thing to prove is the last part, and we will deal with two main cases.

(1)  $N \models (a_\eta^i R a_{\langle 1 \rangle}^j) \wedge (a_\eta^i R x^k) \wedge (x^k E a_{\langle 1 \rangle}^j)$ . We aim to show  $N \models (x^k = a_{\langle 1 \rangle}^j)$ . We know:

$$(*)_1 \mathfrak{C} \models a_\eta^i R a_{\langle 1 \rangle}^j$$

$$(*)_2 N_0 \models a_\eta^i R x^k, \text{ therefore } \varphi(\bar{x}, \bar{y}) \vdash y^i R x^k$$

$$(*)_3 N_1 \models x^k E a_{\langle 1 \rangle}^j, \text{ therefore } \varphi(\bar{x}, \bar{y}) \vdash x^k E y^j.$$

So we can conclude:

$$(*)_2 \Rightarrow a_{\langle \rangle}^i R e^k, \quad a_{\langle \rangle}^i R d^k$$

$$(*)_3 \Rightarrow e^k E a_{\langle \rangle}^j, \quad d^k E a_{\langle \rangle}^j \Rightarrow e^k E d^k.$$

As the above two relations hold in  $\mathfrak{C}$ , which is a model of  $T_{feq}$ , we get  $\mathfrak{C} \models e^k = d^k$ . Denote  $e^* = e^k = d^k$ .

$$(*)_1 \Rightarrow a_\eta^i R a_{\langle 1 \rangle}^j$$

$$(*)_2 \Rightarrow a_\eta^i R e^*$$

$$(*)_1 \Rightarrow e^* E a_{\langle 1 \rangle}^j$$

Together (once again,  $\mathfrak{C} \models T_{feq}$ ) we get  $e^* = a_{\langle 1 \rangle}^j$ , therefore  $\varphi(\bar{x}, \bar{y}) \vdash x^k = y^j$ , so  $N_1 \models x^k = a_{\langle 1 \rangle}^j$ , and we are done.

(2)  $N \models (x^i R a_{\langle 1 \rangle}^j) \wedge (x^i R a_\eta^k) \wedge (a_\eta^k E a_{\langle 1 \rangle}^j)$  and we aim to show  $N \models (a_\eta^k = a_{\langle 1 \rangle}^j)$ .

We know:

$$(*)_1 N_1 \models x^i R a_{\langle 1 \rangle}^j, \text{ so } \varphi(\bar{x}, \bar{y}) \vdash x^i R y^j$$

$$(*)_2 N_0 \models x^i R a_\eta^k, \text{ so } \varphi(\bar{x}, \bar{y}) \vdash x^i R y^k$$

$$(*)_3 \mathfrak{C} \models a_\eta^k E a_{\langle 1 \rangle}^j$$

Note that by indiscernibility of  $\langle \bar{a}_r : r \in {}^{w>2} \rangle$  and  $(*)_3$  we get  $a_{\langle 0 \rangle}^k E a_{\langle 1 \rangle}^j$ , therefore  $a_{\langle 0 \rangle}^k E a_\eta^k$ . Now, by  $(*)_2$ ,  $e^i R a_\eta^k$  &  $e^i R a_{\langle 0 \rangle}^k$ . Therefore, by

$\mathfrak{C} \models T_{feq}$ ,  $a_{\langle 0 \rangle}^k = a_\eta^k$ . Now by indiscernibility

$$a_{\langle 0 \rangle}^k = a_{\langle \cdot \rangle}^k, \quad a_{\langle 1 \rangle}^k = a_{\langle \cdot \rangle}^k$$

So we get that all of the above are equal (and in fact  $a_{r_1}^k = a_{r_2}^k$  for all  $r_1, r_2 \in {}^{w>2}$ ).

Now:

$$(*)_1 \Rightarrow d^i R a_{\langle 1 \rangle}^j$$

$$(*)_2 \Rightarrow d^i R a_{\langle 1 \rangle}^k \Rightarrow d^i R a_\eta^k \text{ (as } a_{\langle 1 \rangle}^k = a_\eta^k)$$

$$(*)_3 \Rightarrow a_\eta^k E a_{\langle 1 \rangle}^j.$$

By  $\mathfrak{C} \models T_{feq}$ , we conclude  $a_\eta^k = a_{\langle 1 \rangle}^k$ , which finishes the proof of the lemma, and therefore the proof of the theorem.  $\square_2$

$\square$

Our next goal is to show that there is a rank function closely related to being (N)SOP<sub>1</sub>. Let  $\varphi(\bar{x}, \bar{y})$  be a formula.

**Definition 2.2.** Given (partial) types  $p(\bar{x}), q(\bar{y})$ . By induction on  $n < \omega$  we define when

$$\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p(\bar{x}), q(\bar{y})) \geq n :$$

If  $\underline{n = 0}$ , this happens if both  $p(\bar{x}), q(\bar{y})$  are consistent

For  $\underline{n + 1}$ , the rank is  $\geq n + 1$  if for some  $\bar{c} \models q(\bar{y})$ , both

$$\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p(\bar{x}) \cup \{\varphi(\bar{x}, \bar{c})\}, q(\bar{y})) \geq n$$

and

$$\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p(\bar{x}), q(\bar{y}) \cup \{\neg(\exists \bar{x})(\varphi(\bar{x}, \bar{y}) \wedge \varphi(\bar{x}, \bar{c}))\}) \geq n.$$

We say  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p(\bar{x}), q(\bar{y})) = \infty$  iff  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p(\bar{x}), q(\bar{y})) \geq n$  for all  $n$ .

We say the rank is  $-1$  if it is not bigger or equal to 0.

*Remark 2.3.* (1) (Definability) Given formulae  $\theta_1, \theta_2$  and  $n < \omega$ , the statement  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(\theta_1(\bar{x}; \bar{a}), \theta_2(\bar{x}; \bar{b})) \geq n$  is a first order formula with parameters  $\bar{a}, \bar{b}$ .

(2) (Finite Character) If  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p(\bar{x}), q(\bar{y})) = n$ , then for some finite  $p_0(\bar{x}) \subseteq p(\bar{x})$  and  $q_0(\bar{y}) \subseteq q(\bar{y})$  we have  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0(\bar{x}), q_0(\bar{y})) = n$ .

(3) (Monotonicity) If  $p' \vdash p''$  and  $q' \vdash q''$ , then  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p', q') \leq \text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p'', q'')$ .

(4) We can continue to define when  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p(\bar{x}), q(\bar{y})) \geq \alpha$  for any ordinal  $\alpha$ , but by the compactness theorem, part (1) (Definability) and part (2) (the Finite Character) it follows that  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p, q) \geq \alpha$  for some  $\alpha \geq \omega$  iff  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p, q) \geq \omega$  iff  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p, q) = \infty$ .

(5) If  $p' \equiv p''$ , and  $q' \equiv q''$ , then  $\text{rk}_{\varphi}^1(p', q') = \text{rk}_{\varphi}^1(p'', q'')$ .

We aim to show that  $\text{rk}_{\varphi}^1(p(\bar{x}), q(\bar{y}))$  is finite for every  $p(\bar{x}), q(\bar{y})$  (or, equivalently,  $\text{rk}_{\varphi}^1(\bar{x} = \bar{x}, \bar{y} = \bar{y})$  is finite) if and only if  $\varphi(\bar{x}, \bar{y})$  does not exemplify  $SOP'_1$  in  $T$ . For this purpose we shall need another definition and several easy claims.

**Definition 2.4.** Given (partial) types  $p(\bar{x})$  and  $q(\bar{y})$ , we say that  $\{\bar{a}_\eta : \eta \in {}^{n \geq 2}\}$  is a  $\varphi$ - $SOP'_1$  tree for  $p(\bar{x})$  and  $q(\bar{y})$  (of depth  $n$ ) if

- (a)  $p(\bar{x}) \cup \{\varphi^{\eta(i)}(\bar{x}, \bar{a}_{\eta \upharpoonright i}) : i < n\}$  is consistent for every  $\eta \in {}^{n \geq 2}$ .
- (b)  $\bar{a}_\eta \models q(\bar{y})$  for all  $\eta \in {}^{n \geq 2}$
- (c) If  $\eta, \nu$  are in  ${}^{n \geq 2}$  satisfying  $\eta \frown \langle 0 \rangle \trianglelefteq \nu$ , then the set  $\{\varphi(\bar{x}, \bar{a}_\eta), \varphi(\bar{x}, \bar{a}_\nu)\}$  is inconsistent.

**Proposition 2.5.** *Suppose  $\{\bar{a}_\eta : \eta \in {}^{n \geq 2}\}$  is a  $\varphi$ - $SOP'_1$  tree for  $p(\bar{x})$  and  $q(\bar{y})$  of depth  $n$ , and denote  $A^0 = \{\bar{a}_\eta : \langle 0 \rangle \leq \eta\}$ ,  $A^1 = \{\bar{a}_\eta : \langle 1 \rangle \leq \eta\}$ .*

*Then*

- (1)  $A^1$  is a  $\varphi$ - $SOP'_1$  tree for  $p(\bar{x}) \cup \{\varphi(\bar{x}, \bar{a}_{\langle 0 \rangle})\}$  and  $q(\bar{y})$
- (2)  $A^0$  is a  $\varphi$ - $SOP'_1$  tree for  $p(\bar{x})$  and  $q(\bar{y}) \cup \{\neg(\exists \bar{x})(\varphi(\bar{x}, \bar{y}) \wedge \varphi(\bar{x}, \bar{a}_{\langle 0 \rangle}))\}$ .

*Proof.* The clauses (a) and (c) of the definition easily hold both for  $A^1$  and  $A^0$ , so we should only check (b), which is also obvious for  $A^1$ .

Therefore, we're left to show that for every  $\eta \in A^0$ ,  $\bar{a}_\eta \models \neg(\exists \bar{x})(\varphi(\bar{x}, \bar{y}) \wedge \varphi(\bar{x}, \bar{a}_{\langle 0 \rangle}))$ , and this is clear by clause (c) of the definition ( $\{\bar{a}_\eta : \eta \in {}^{n \geq 2}\}$  is a  $\varphi$ - $SOP'_1$  tree, and  $\langle \rangle \frown 0 \leq \eta$ ).  $\square$

Now we show the connection between the rank and  $SOP'_1$  trees.

**Proposition 2.6.**  $\text{rk}_\varphi^1(p(\bar{x}), q(\bar{y})) \geq n \iff$  *there exists a  $\varphi$ - $SOP'_1$  tree for  $p(\bar{x})$  and  $q(\bar{y})$  of depth  $n$ .*

*Proof.* Both directions are proved by induction on  $n$ . The case  $n = 0$  is obvious. For  $n = m + 1$ , the right-to-left direction follows immediately by the induction hypothesis and 2.5. So we will elaborate more only about the other direction, although it is also straightforward.

Suppose  $n = m + 1$  and  $\text{rk}_\varphi^1(p(\bar{x}), q(\bar{y})) \geq n$ . By the definition of the rank and the induction hypothesis, for some  $\bar{c} \models q(\bar{y})$ , there are

- (1) a  $\varphi$ - $SOP'_1$  tree  $A^1 = \{\bar{a}_\eta^1 : \eta \in {}^{m \geq 2}\}$  for  $p(\bar{x}) \cup \{\varphi(\bar{x}, \bar{c})\}$  and  $q(\bar{y})$
- (2) a  $\varphi$ - $SOP'_1$  tree  $A^0 = \{\bar{a}_\eta^0 : \eta \in {}^{m \geq 2}\}$  for  $p(\bar{x})$  and  $q(\bar{y}) \cup \{\neg(\exists \bar{x})(\varphi(\bar{x}, \bar{y}) \wedge \varphi(\bar{x}, \bar{c}))\}$

(both of depth  $m$ ). We define a tree  $\{\bar{a}_\eta : \eta \in {}^{n \geq 2}\}$  by

$$\bar{a}_\langle \rangle = \bar{c}$$

$$\bar{a}_{\langle \ell \rangle \smallfrown \eta} = \bar{a}_\eta^\ell \text{ for } \ell \in \{0, 1\}$$

which is as required, i.e. a  $\varphi$ - $SOP'_1$  tree for  $p(\bar{x})$  and  $q(\bar{y})$ , since:

- (a) of the definition obviously holds by (1) above.
- (b) holds as  $\bar{c} \models q(\bar{y})$ .
- (c) obviously holds by (2) above.

□

The following remark is obvious:

*Remark 2.7.*  $\varphi(\bar{x}, \bar{y})$  exemplifies  $SOP'_1$  in  $T \iff$  there exists a  $\varphi$ - $SOP'_1$  tree for  $\bar{x} = \bar{x}$  and  $\bar{y} = \bar{y}$  of any depth.

So we can conclude the following

**Theorem 2.8.** *A formula  $\varphi(\bar{x}, \bar{y})$  does not exemplify  $SOP'_1$  in  $T \iff \text{rk}_\varphi^1(\bar{x} = \bar{x}, \bar{y} = \bar{y}) < \omega \iff \text{rk}_\varphi^1(p(\bar{x}), q(\bar{y})) < \omega$  for every two (partial) types  $p(\bar{x})$  and  $q(\bar{y})$ . Moreover,  $\text{rk}_\varphi^1(\bar{x} = \bar{x}, \bar{y} = \bar{y})$  is exactly the maximal depth of a  $\varphi$ - $SOP'_1$  tree that can be built in  $\mathfrak{C}$ .*

**Corollary 2.9.**  *$T$  does not have  $SOP_1 \iff T$  does not have  $SOP'_1 \iff \text{rk}_\varphi^1(\bar{x} = \bar{x}, \bar{y} = \bar{y})$  is finite for every formula  $\varphi(\bar{x}, \bar{y})$ .*

Now we show an application of the rank.

**Theorem 2.10.** *Suppose that  $T$  satisfies  $NSOP_1$ . Assume that*

- (a)  $M_1 \prec M_2 \prec \mathfrak{C}$ .

- (b)  $p$  is a (not necessarily complete) type over  $M_2$ , containing the formula  $\varphi(\bar{x}, \bar{b}^*)$  for some  $\bar{b}^* \in M^2 \setminus M^1$ .

Then for some finite  $q' \subseteq \text{tp}(\bar{b}^*/M_1)$  at least one of the following holds:

- (i) If  $\bar{b} \in M_1$  realises  $q'(\bar{y})$  then  $\varphi(\bar{x}, \bar{b}) \notin p$ , or  
(ii) If  $\bar{b} \in M_1$  realises  $q'(\bar{y})$  then  $\{\varphi(\bar{x}, \bar{b}), \varphi(\bar{x}, \bar{b}^*)\}$  is consistent.

In fact, all we need to assume for this Claim is that  $\varphi(\bar{x}, \bar{y})$  does not exemplify SOP<sub>1</sub> in  $\mathfrak{C}$ .

*Proof.* Denote  $q = \text{tp}(b^*/M_1)$ . As  $T$  is NSOP<sub>1</sub>, we have that  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p \upharpoonright M_1, q) = n^* < \omega$  (certainly  $n^* \geq 0$ ). By the finite character of the rank, we have that for some finite  $p_0 \subseteq p \upharpoonright M_1$  and  $q_0 \subseteq q$ ,

$$\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0, q_0) = n^*.$$

Hence for no  $\bar{c} \models q_0(\bar{y})$  do we have that both  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0 \cup \{\varphi(\bar{x}, \bar{c})\}, q_0) \geq n^*$  and  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0, q_0 \cup \{(\neg \exists \bar{x})[\varphi(\bar{x}, \bar{y}) \wedge \varphi(\bar{x}, \bar{c})]\}) \geq n^*$ . In particular, this holds for  $\bar{c} = \bar{b}^*$  (remember that  $\bar{b}^* \models q$  and therefore certainly  $\bar{b}^* \models q_0$ ). So

⊗ 2.10.1.

If  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0 \cup \{\varphi(\bar{x}, \bar{b}^*)\}, q_0) \geq n^*$ , then  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0, q_0 \cup \{(\neg \exists \bar{x})[\varphi(\bar{x}, \bar{y}) \wedge \varphi(\bar{x}, \bar{b}^*)]\}) < n^*$ .

By Remark 2.3(1), there is a finite  $q' \subseteq q$  such that

⊗ 2.10.2.

$\bar{b}$  realises  $q' \implies \text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0 \cup \{\varphi(\bar{x}, \bar{b})\}, q_0) = \text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0 \cup \{\varphi(\bar{x}, \bar{b}^*)\}, q_0)$ .

We aim to show that  $q'$  is as required.

Case 1.  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0 \cup \{\varphi(\bar{x}, \bar{b}^*)\}, q_0) = n < n^*$ .

We note that possibility (i) holds.

Namely, suppose  $\bar{b}$  realises  $q'$ , then  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0 \cup \{\varphi(\bar{x}, \bar{b})\}, q_0) = n < n^*$ , so if  $\varphi(\bar{x}, \bar{b}) \in p$ , we obtain a contradiction with monotonicity of the rank.

Case 2.  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0 \cup \{\varphi(\bar{x}, \bar{b}^*)\}, q_0) = n^*$ .

We shall show that (ii) holds.

Suppose otherwise, so let  $\bar{b} \in M_1$  realise  $q'$  and  $\{\varphi(\bar{x}, \bar{b}), \varphi(\bar{x}, \bar{b}^*)\}$  is contradictory. By 2.10.2,

$$\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0 \cup \{\varphi(\bar{x}, \bar{b})\}, q_0) = n^*$$

and by 2.10.1,

$$\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p_0, q_0 \cup \{(\neg \exists \bar{x})(\varphi(\bar{x}, \bar{y}) \wedge \varphi(\bar{x}, \bar{b}))\}) < n^*.$$

We have that  $(\neg \exists \bar{x})[\varphi(\bar{x}, \bar{y}) \wedge \varphi(\bar{x}, \bar{b})] \in q$ , hence  $q_0 \cup \{(\neg \exists \bar{x})[\varphi(\bar{x}, \bar{y}) \wedge \varphi(\bar{x}, \bar{b})]\} \subseteq q$ , in contradiction with monotonicity and  $\text{rk}_{\varphi(\bar{x}, \bar{y})}^1(p \upharpoonright M_1, q) = n^*$ .  $\square$

### 3. MORE ON $SOP_2$ , $SOP_3$ AND $\triangleleft_\lambda^*$ -ORDER

We try to find a connection between the syntactic properties  $SOP_2, SOP_3$  and the semantic property of being  $\triangleleft_\lambda^*$ -maximal. Our guess is that  $\triangleleft_\lambda^*$ -maximality should be equivalent to one of the above order properties (maybe both), but all we prove here is  $SOP_3 \implies \triangleleft_\lambda^*$ -maximality  $\implies SOP_2$ . We also give a weaker “local” result in the other direction.

First we generalize the definitions from [DžSh692], of  $\triangleleft_{\lambda}^*$ -maximality, making them local as well as global.

**Definition 3.1.** (1) For given (complete first order theories)  $T_1, T_2$

and cardinals  $\lambda \geq \mu > \kappa, \mu \geq \theta > |T_1| + |T_2| + \aleph_0$

(a)  $T_1 \triangleleft_{<\lambda, <\mu, \kappa, <\theta}^* T_2$  means that there is a (complete first order theory)  $T^*$  and interpretations  $\bar{\varphi}_1, \bar{\varphi}_2$  of  $T_1, T_2$  in  $T^*$  respectively,  $|T^*| < \theta$  such that:

–  $\boxtimes_{T^*, \bar{\varphi}_1, \bar{\varphi}_2}^{<\lambda, <\mu, \kappa}$  if  $M$  is a  $\kappa$ -saturated model of  $T^*$  and  $M_\ell = M^{[\bar{\varphi}^\ell]}$  for  $\ell = 1, 2$  and  $M_2$  is  $\lambda$ -saturated (model of  $T_2$ ), then  $M_1$  is  $\mu$ -saturated

(b)  $(T_1, \vartheta_1(\bar{x}, \bar{y})) \triangleleft_{<\lambda, <\mu, <\kappa}^* (T_2, \vartheta_2(\bar{x}, \bar{y}))$  means that  $\vartheta_\ell(\bar{x}, \bar{y}) \in L(\tau_{T_\ell})$  and that there is a  $T^*$  and interpretations  $\bar{\varphi}_1, \bar{\varphi}_2$  of  $T_1, T_2$  in  $T^*$  respectively,  $|T^*| < \mu$  such that  $\boxtimes_{T^*, \vartheta_1, \vartheta_2, \bar{\varphi}_1, \bar{\varphi}_2}^{<\lambda, <\mu, \kappa}$  if  $M$  is a  $\kappa$ -saturated model of  $T^*$  and  $M_\ell = M^{[\bar{\varphi}^\ell]}$  for  $\ell = 1, 2$  and  $M_2$  is  $(\lambda, \vartheta_1(\bar{x}, \bar{y}))$ -saturated (see 3 below), then  $M_i$  is  $(\mu, \vartheta_2)$ -saturated.

(2) Instead of “ $< \lambda^+$ ” we may write “ $\lambda$ ”, and instead of “ $< \mu^+$ ” we may write  $\mu$ , instead of  $< \theta^+$  we may write  $\theta$ . If we omit  $\mu$  we mean  $\mu = \lambda$ , and if we write  $\kappa = 0$  then “ $\kappa$ -saturated” becomes the empty demand, if we omit  $\theta$  we mean  $|T_1| + |T_2| + \aleph_0$  and if we omit  $\kappa$  and  $\theta$  then we mean that  $\mu = \lambda, \theta = |T_1| + |T_2| + \aleph_0$ .

(3) We say  $M$  is  $(\lambda, \Delta)$ -saturated when: if  $p \subseteq \{\vartheta(\bar{x}; \bar{a}) : \vartheta(\bar{x}; \bar{y}) \in \Delta, \bar{a} \in {}^{\ell g(\bar{y})} M\}$  is finitely satisfiable of cardinality  $< \lambda$  then  $p$  is realized in  $M$ . If  $\Delta = \{\vartheta(\bar{x}, \bar{y})\}$  we may write  $\vartheta(\bar{x}, \bar{y})$  instead of  $\Delta$ .

- (4) If  $T_1, T_2$  are not necessarily complete, then above  $T^*$  is not necessarily complete and we demand: if  $M_1 \models T_1, M_2 \models T_2$  then there is  $M \models T^*$  such that  $M^{[\bar{\varphi}_\ell]} \models Th(M_\ell)$  for  $\ell = 1, 2$ .
- (5) We say  $T$  is  $\triangleleft_{\lambda, \kappa}^*$ -maximal if  $|T'| < \lambda \Rightarrow T' \triangleleft_{\lambda, \kappa}^* T$ . We say  $(T, \vartheta(\bar{x}; \bar{y}))$  is  $\triangleleft_{\lambda, \kappa}^*$ -maximal if  $|T'| < \lambda \& \vartheta'(\bar{x}'; \bar{y}') \in L(\mathcal{T}_{T'}) \Rightarrow (T', \vartheta'(\bar{x}'; \bar{y}')) \triangleleft_{\lambda, \kappa}^* (T, \vartheta(\bar{x}; \bar{y}))$ .

**Definition 3.2.** (1)  $T_{tr}$  is the theory of trees (i.e. the vocabulary is  $\{<\}$  and the axioms state that  $<$  is a partial order and  $\{y : y < x\}$  is a linear order for every  $x$ ), so  $T_{tr}$  is not complete, and let  $\vartheta_{tr}(x, y) = (y < x)$ .

- (2)  $T_{tr}^*$  is the model completion of  $T_{tr}$ .
- (3)  $T_{ord}$  is the theory of linear orders,  $T_{ord}^*$  is its model completion (i.e. the theory of dense linear order without endpoints).

We note connection to previous works and obvious properties

- Proposition 3.3.** (1)  $T_1 \triangleleft_{\lambda, \mu, 0}^* T_2$  is  $T_1 \triangleleft_{\lambda, \mu}^* T_2$  of [DžSh692].
- (2)  $T_1 \triangleleft_{\lambda, \lambda; < \kappa}^* T_2$  implies  $T_1 \triangleleft_{\lambda, \kappa}^* T_2$  of [Sh500].
- (3)  $\triangleleft_{\lambda, \mu; \kappa, \theta}^*$  has the obvious monotonicity properties: if  $T_1 \triangleleft_{< \lambda_1, < \mu'_1; < \kappa_1, < \theta_1}^* T_2$  and  $\lambda_2 \geq \lambda_1, \mu_2 \leq \mu_1, \kappa_2 \geq \kappa_1, \theta_2 \geq \theta_1$  then  $T_1 \triangleleft_{< \lambda_2, < \mu_2; < \kappa_2, < \theta_2}^* T_2$ .
- (4)  $T \triangleleft_{\lambda, \mu; \kappa, \theta}^* T$  if  $|T| < \theta, \lambda \geq \mu > \kappa, \mu \geq \theta$ .
- (5) If  $\mu$  is a limit cardinal, then  $T_1 \triangleleft_{< \lambda, < \mu; < \kappa, < \theta}^* T_2$  iff for every  $\mu_1 < \mu, \mu_1 \geq \kappa$  we have

$$T_1 \triangleleft_{< \lambda, < \mu_1; < \kappa, < \theta}^* T_2.$$

(6) *Similar results hold for  $(T_\ell, \vartheta_\ell(\bar{x}; \bar{y}))$ .*

*Proof.* Easy. □

**Proposition 3.4.** (1) *Assume  $T_1 \triangleleft_{<\lambda, <\mu; <\kappa, <\theta}^* T_2$ . Then for any theory  $T^*$ , we can find  $T^{**} \supseteq T^*$  complete  $|T^{**}| < (|T^*|^{\tau(T_1)} + |T^*|^{\tau(T_2)})^{++}$   $\theta$  such that: for any interpretations  $\bar{\varphi}_1, \bar{\varphi}_2$  of  $T_1, T_2$  in  $T^{**}$  respectively the definition 3.1(1) of  $T_1 \triangleleft_{<\lambda, <\mu; <\kappa, <\theta}^* T_2$  holds.*

(2) *Assume  $\tau(T_1), \tau(T_2)$  are disjoint. Then  $T_1 \triangleleft_{<\lambda, <\mu; <\kappa, <\theta}^* T_2$  if for any  $T \supseteq T_1 \cup T_2$  there is  $T^* \supseteq T$  as demanded in Definition 3.1(1) for the trivial interpretations  $M^{[\bar{\varphi}^\ell]}$  is the  $\tau(T_\ell)$ -reduct.*

*Proof.* Easy, or see [DžSh692], Observation 1.4. □

Now we will show that  $T_{tr}^*$  is  $\triangleleft_\lambda^*$ -maximal for every  $\lambda$  big enough, and conclude that  $SOP_3 \implies \triangleleft_\lambda^*$ -maximality. The last result appears already in [Sh500], Theorem 2.9, but the proof is not complete - in fact, the proof shows the following theorem:

**Theorem 3.5.** *Any theory  $T$ ,  $|T| < \lambda$ , with  $SOP_3$  is  $\triangleleft_\lambda^*$ -above  $T_{ord}^*$ .*

*Proof.* See [Sh500], (2.12). □

Here we prove explicitly that  $T_{tr}^*$ , and therefore  $T_{ord}^*$  are maximal.

**Theorem 3.6.**  *$T_{tr}^*$  is  $\triangleleft_\lambda^*$ -maximal for any  $\lambda > \aleph_0$ ; the witness  $T^*$  does not depend on  $\lambda$ .*

*Proof.* Let  $T$  be any complete theory,  $|T| < \lambda$  and  $M_1$  a model of  $T$ .

Let  $\Phi = \{\varphi(x, \bar{a}) : \varphi(x, \bar{y}) \in L(\tau_T), \bar{a} \in {}^{\ell g(\bar{y})}(M_1)\}$ , so  $|\Phi| = \|M_1\|$ . So  $M = (\omega \triangleright \Phi, \triangleleft)$  is a model of  $T_{tr}$  and there is a model  $M_2$  of  $T_{tr}^*$  of

cardinality  $\|M_1\|$  extending  $M$  such that every member of  $M_2$  is below some member of  $M$ .

Let  $\chi$  be large enough such that  $M_1, M_2 \in \mathcal{H}(\chi)$  and we define  $\mathcal{B}^*$  expanding  $(\mathcal{H}(\chi), \in)$  by  $P_1 = |M_1|, P_2 = |M_2|, P = |M|, Q_0 = \Phi, <_1 = <^{M_2}, < = <_1 \upharpoonright P, m$  a constant symbol for a set  $M_1, R^{\mathcal{B}^*} = R^{M_1}$  for  $R \in \tau_T$  (wlog  $\tau(T)$  does not contain any other predicate mentioned here)

$$Q = \{(\langle \varphi_\ell(x, \bar{a}_\ell) : \ell < n \rangle : M_1 \models \exists x[\wedge \varphi_\ell(x, \bar{a}_\ell)]\}.$$

$H$  is a partial unary function with domain  $Q$  and range  $P_1, H(\langle \varphi_\ell(x, \bar{a}_\ell) : \ell < n \rangle)$  satisfies  $\{\varphi_\ell(x, \bar{a}_\ell) : \ell < n\}$ , i.e.  $\mathcal{B}^*$  satisfies the formula “ $m \models (\exists x) \bigwedge_{\ell < n} \varphi_\ell(x, \bar{a}_\ell)$ ”.

Let  $T^* = Th(\mathcal{B}^*)$ , let  $\bar{\varphi}_1$  be the trivial interpretation of  $T$  in  $T^*$  (the restriction + reduct) and  $\bar{\varphi}_2 = \langle P_2(x), x_0 <_1 x_1 \rangle$  is an interpretation of  $T_{tr}^*$ . So  $T^*, \bar{\varphi}_1, \bar{\varphi}_2$  does not depend on  $\lambda$ .

Now we assume  $\mathcal{B}$  is a model of  $T^*, N_1 = \mathcal{B}^{[\bar{\varphi}_1]}, N_2 = \mathcal{B}^{[\bar{\varphi}_2]}, N_3 = (P^{\mathcal{B}}, <^{\mathcal{B}})$  and we aim to show that (i) below implies (iii). We will first show that (i)  $\Rightarrow$  (ii) and use this fact in the proof.

- (i)  $N_2$  is  $\lambda$ -saturated
- (ii) in  $N_3$  every branch has cofinality  $\geq \lambda$ , equivalently: every increasing sequence of length  $< \lambda$  has an upper bound
- (iii)  $N_1$  is  $\lambda$ -saturated.

Let us first show (i)  $\Rightarrow$  (ii). If  $\langle a_i : i < \delta \rangle$  is  $<^{N_3}$ -increasing,  $\delta < \lambda$  then it is  $<^{N_2}$ -increasing hence has a  $<^{N_2}$ -upper bound  $a$  but  $(\forall x \in$

$P_2)(\exists y)(x <_1 y \& P(y))$  belongs to  $T^*$  so there is  $b, a <^{N_2} b \in P^N = N_3$  so  $b$  is as required.

So we can assume clause (i) and we shall prove (iii).

Before we proceed, let us note several easy but important properties of  $\mathcal{B}$ .

- (a) We can talk inside  $\mathcal{B}$  about a set being a model, (standard coding of) a formula, a proof, etc. In particular, we can speak about  $m$  (as a model) satisfying or not satisfying certain sentences. Also, given a formula with free variables we can speak about substitution of other variables or parameters into the formula. Given  $s \in \mathcal{B}$  which is a formula with free variables  $\bar{x}$ , we will allow ourselves to write  $s = s(\bar{x})$ , and if  $\mathcal{B}$  thinks that substitution of  $\bar{a} \in P_1$  into  $s$  will turn it into a true sentence in  $m$  as a model, we will write  $m \models s(\bar{a})$  or just  $s(\bar{a})$ .
- (b)  $\mathcal{B} \models \forall z Q_0(z) \iff$  “ $z$  is a formula with one free variable with parameters from  $P_1$ ”. Moreover, suppose  $\varphi(x, \bar{a})$  is a formula in  $L(\tau_T)$  s.t.  $\bar{a} \in P_1^{\mathcal{B}}$ .  $\mathcal{B}^*$  and therefore  $\mathcal{B}$  satisfy  $(\forall \bar{y} \in P_1)(\exists! s \in Q_0)$  such that  $(\forall x \in P_1)\varphi(x, \bar{y}) \iff “m \models s(x, \bar{y})”$ . Let us denote by  $\ulcorner \varphi(x, \bar{a}) \urcorner$  this “canonical encoding” of  $\varphi(x, \bar{a})$  in  $Q_0^{\mathcal{B}}$ .
- (c)  $\mathcal{B} \models \forall s P(s) \iff$  “ $s$  is a finite sequence of members of  $Q_0$ , i.e.  $(\exists n \in \omega)(s : n \rightarrow Q_0)$ ”.
- (d) For simplicity of notation, given  $s \in P^{\mathcal{B}}$ , we will write “ $z \in s$ ” instead of “ $z \in Im(s^{\mathcal{B}})$ ”.

- (e) For  $z \in P^{\mathcal{B}}$ ,  $c \in P_1^{\mathcal{B}}$ , we write  $z(c)$  meaning  $(\forall s \in z)s(c)$ .
- (f) For every  $\varphi(x, \bar{a}) \in L(\tau_T)$  for  $a \in P_1^{\mathcal{B}}$ , there exists an element of  $P^{\mathcal{B}}$  corresponding to the finite sequence  $\langle \varphi(x, \bar{a}) \rangle$ . We denote this element by  $\langle \ulcorner \varphi(x, \bar{a}) \urcorner \rangle$ . Moreover,  $\mathcal{B} \models \exists x(P_1(x) \wedge \varphi(x, \bar{a})) \rightarrow Q(\langle \ulcorner \varphi(x, \bar{a}) \urcorner \rangle)$ .

*Subclaim 3.6.1.* (1) Suppose  $\mathcal{B} \models Q(z)$ . Then  $\mathcal{B} \models \forall w(Q(w) \wedge z < w) \rightarrow z(H(w))$ .

- (2) Let  $\varphi(x, \bar{a}) \in L(\tau_T)$  and suppose  $\mathcal{B} \models \exists x P_1(x) \wedge \varphi(x, \bar{a})$ . Then  $\mathcal{B} \models \forall z(Q(z) \wedge \langle \ulcorner \varphi(x, \bar{a}) \urcorner \rangle < z) \rightarrow \varphi(H(z), \bar{a})$ .

*Proof.* (1) Trivial as  $\mathcal{B}^*$  satisfies it.

- (2) Let  $z^* = \langle \ulcorner \varphi(x, \bar{a}) \urcorner \rangle$ . First,  $Q(z^*)$  holds by (f) above. By (1),  $z^*(H(z))$  holds for each  $z \in Q^{\mathcal{B}}$ ,  $z^* < z$ . Now by (b) and (f) above,  $\mathcal{B} \models \forall x P_1(x) \rightarrow (z^*(x) \iff \varphi(x, \bar{a}))$ . As  $\mathcal{B} \models \text{Range}(H) \subseteq P_1$ , we are done. □

We now proceed with the proof (i)  $\implies$  (iii). So let  $p$  be a 1-type in  $N_1$  of cardinality  $< \lambda$ , so let  $p = \{\varphi_\beta(x, \bar{a}_\beta) : \beta < \alpha\}$  with  $\alpha < \lambda$ ,  $\bar{a}_\beta \in N_1 \forall \beta$ . Without loss of generality  $p$  is closed under conjunction, i.e. for every  $\varepsilon, \zeta < \alpha$  for some  $\xi < \alpha$  we have  $\varphi_\xi(x, \bar{a}_\xi) = \varphi_\varepsilon(x, \bar{a}_\varepsilon) \wedge \varphi_\zeta(x, \bar{a}_\zeta)$ . We shall now choose by induction on  $\beta \leq \alpha$  an element  $b_\beta$  of  $N$  such that

- (A)  $b_\beta \in P^{\mathcal{B}} = N_3$  moreover  $b_\beta \in Q^{\mathcal{B}}$  and  $\gamma < \beta \implies b_\gamma <^{N_3} b_\beta$
- (B) if  $\gamma < \beta$  then  $\mathcal{B} \models (\forall z)(Q(z) \wedge (b_\beta \leq z) \rightarrow \varphi_\gamma(H(z), \bar{a}_\gamma))$

(C) if  $\gamma < \alpha$  (but not necessarily  $\gamma < \beta$ ) then  $\mathcal{B} \models (\exists z)[Q(z) \wedge (b_\beta \leq z) \wedge (\forall y)(Q(y) \wedge z \leq y \rightarrow \varphi_\gamma(H(y), \bar{a}_\gamma))]$ .

If we succeed then  $H^{\mathcal{B}}(b_\alpha)$  is as required.

Case 1:  $\beta = 0$ .

Define  $b_0 = \langle \rangle$  (the element of  $P^{\mathcal{B}}$  corresponding to the empty sequence). Clearly  $\mathcal{B} \models Q(b_0)$ , i.e. the demand (A) holds. (B) holds trivially. Why does (C) hold? Let  $\gamma < \alpha$ .  $\mathcal{B} \models \exists x \varphi_\gamma(x, \bar{a}_\gamma)$  therefore denoting  $z_\gamma^* = \langle \ulcorner \varphi_\gamma(x, \bar{a}_\gamma) \urcorner \rangle$ , we have  $\mathcal{B} \models Q(z_\gamma^*) \wedge b_0 < z_\gamma^*$ . Now we finish by part (2) of the subclaim.

Case 2:  $\beta = v + 1$ .

$\mathcal{B}$  satisfies the sentence saying that for every  $\eta \in Q$  and  $\bar{y} \in P_1$  there exists an element of  $P$  that we denote by  $Conc_v(\eta, \bar{y})$  corresponding to  $\eta \hat{\ } \langle \ulcorner \varphi_v(x, \bar{y}) \urcorner \rangle$ . We define  $b_\beta = Conc_v(b_v, \bar{a}_v)$ . Now we have to check (A) - (C).

(A) By the induction hypothesis, clause (C) holds for  $b_v$  and  $v$  (standing for  $b_\beta$  and  $\gamma$  there). Therefore  $\mathcal{B} \models \exists z \in Q (b_v \leq z) \wedge \varphi_v(H(z), \bar{a}_v)$ . But  $\mathcal{B}^*$  (and so  $\mathcal{B}$ ) satisfies that  $\forall \bar{y} \in P_1$  if there exists  $z \in Q$  s.t.  $\varphi_v(H(z), \bar{y})$  holds, then  $Conc_v(z, \bar{y})$  is an element of  $Q$  (as in  $\mathcal{B}^*$  the assumption means that there exists an element of  $m$  satisfying all the formulae in  $z$  plus  $\varphi_v(x, \bar{y})$ ). So we get the required.

(B) is clear as by the induction hypothesis,  $\varphi_\zeta(H(z), \bar{a}_\zeta)$  holds for every  $\zeta < v$ ,  $b_\beta \leq z$  (recall that  $b_v \leq b_\beta$ ). As for  $\varphi_v(x, \bar{a}_v)$ ,  $\mathcal{B}^*$  clearly satisfies that for every  $z \in Q, \bar{y} \in P_1$ , if  $b = Conc_v(z, \bar{y})$  is in  $Q$  then  $\varphi_v(H(z), \bar{y})$  holds  $\forall z \in Q, b \leq z$ .

(C) Let  $\zeta < \alpha$ . As  $p$  is closed under conjunctions, for some  $\xi$ ,  $\varphi_\gamma(x, \bar{a}_\gamma) \wedge \varphi_\zeta(x, \bar{a}_\zeta) = \varphi_\xi(x, \bar{a}_\xi)$ . Now we apply clause (C) holding for  $b_v$  to  $\gamma = \xi$  and get  $z \in Q, b_v \leq z$  with  $H(z)$  satisfying both  $\varphi_v(x, \bar{a}_v)$  and  $\varphi_\zeta(x, \bar{a}_\zeta)$ . Once again using the satisfaction by  $\mathcal{B}$  of natural sentences, we show that  $b = \text{Conc}_\zeta(b_\beta, \bar{a}_\zeta)$  is in  $Q, b_\beta \leq b$  and  $\forall z \in Q$  which is above  $b, \varphi_\zeta(x, \bar{a}_\zeta)$  holds, i.e.  $b$  is as required.

Case 3:  $\beta = \delta$  limit.

By our present assumption, clause (i), and therefore clause (ii), hold. Hence there is  $b \in P^{\mathcal{B}}$  which is an upper bound to  $\{b_\gamma : \gamma < \beta\}$ . Now  $\mathcal{B}$  satisfies “for every element  $z$  of  $P$  there is a  $y \leq z$  which is in  $Q$  and  $x \leq z \wedge Q(x) \rightarrow x \leq y$ ”. Apply this to  $b$  for  $z$  and get  $b'_\delta$  for  $y$ . So  $b'_\delta \in Q$  and  $\gamma < \delta \Rightarrow b_\gamma \leq b'_\delta$ , as required in clauses (A) + (B) but not necessarily (C).

Define for each  $\zeta < \alpha$  a formula  $\psi_\zeta(w, \bar{a}_\zeta) = (\exists z)(w \leq z \wedge Q(z) \wedge (\forall y)(z \leq y \wedge Q(y) \rightarrow \varphi_\zeta(H(y), \bar{a}_\zeta))$  Now we find  $c_\zeta$  (for  $\zeta < \alpha$ ) such that:

- (a)  $c_\zeta \in Q^{\mathcal{B}}, c_\zeta \leq b$
- (b)  $\psi_\zeta(c_\zeta, \bar{a}_\zeta)$  holds.
- (c) under (a) + (b), the element  $c_\zeta$  is maximal.

Why do  $c_\zeta$  exist?  $\mathcal{B}$  satisfies “for every element  $s$  of  $P$  there is a  $w \leq s$  which satisfies  $\psi_\zeta(w, \bar{a}_\zeta)$ , is in  $Q$  and  $(x \leq s \wedge \psi_\zeta(x, \bar{a}_\zeta) \wedge Q(x)) \rightarrow (x \leq w)$ ”.

By the induction hypothesis we have:

$$\gamma < \delta, \zeta < \alpha \Rightarrow b_\gamma <^{N_3} c_\zeta.$$

Clearly it suffices to find  $b_\delta$  satisfying  $Q(b_\delta)$  and  $b_\gamma <^{N_3} b_\delta <^{N_3} c_\zeta$  for  $\gamma < \delta, \zeta < \alpha$ . As  $N_3 \upharpoonright \{c : c \leq b\}$  is linearly ordered, this follows from  $N_2$  being  $\lambda$ -saturated.  $\square$

**Proposition 3.7.** (1) For every  $T^*$ , there is  $T^{**} \supseteq T^*, |T^{**}| = |T^*| + \aleph_0$  such that for every model  $\mathcal{B}$  of  $T^{**}$  we have

(a) for any  $\lambda$ , the following are equivalent

- ( $\alpha$ ) if  $\bar{\varphi}_1$  is an interpretation of  $T_{tr}^*$  in  $\mathcal{B}$  (possibly with parameters) then  $\mathcal{B}^{[\bar{\varphi}_1]}$  is  $\lambda_{tr}$ -saturated
- ( $\beta$ ) if  $\bar{\varphi}_2$  is an interpretation of  $T_{ord}$  in  $\mathcal{B}$  (possibly with parameters) then  $\mathcal{B}^{[\bar{\varphi}_2]}$  is  $\lambda$ -saturated

(b) for any  $\lambda$ , the following are equivalent

- ( $\alpha$ ) if  $\bar{\varphi}_1$  is an interpretation of  $T_{tr}$  in  $\mathcal{B}$  (possibly with parameters) then in  $\mathcal{B}^{[\bar{\varphi}_1]}$ , every branch with no last element has cofinality  $\geq \lambda$
- ( $\beta$ ) if  $\bar{\varphi}_2^*$  is an interpretation of  $T_{ord}$  in  $\mathcal{B}$  (possibly with parameters) then in  $\mathcal{B}^{[\bar{\varphi}_2]}$  there is no Dedekind cut  $(I_1, I_2)$  with both cofinalities  $< \lambda$  and at least one  $\geq \aleph_0$ .

*Proof.* Easy.  $\square$

**Corollary 3.8.** (1)  $T_{ord}^*$  is  $\triangleleft_\lambda^*$ -maximal.

(2) If  $|T| < \lambda$  and  $T$  has  $SOP_3$  then  $T$  is  $\triangleleft_\lambda^*$ -maximal.

*Proof.* (1) Follows from 3.7

(2) By (1) and 3.5.

□

*Question 3.9.* Is the other direction of 3.8 (2) true?

*Remark 3.10.* See Theorem 3.13 and Corollary 3.15 below for a proof of a weaker version of the other direction: we get  $SOP_2$  instead of  $SOP_3$ .

We would like to know whether it is possible to weaken the assumptions of Corollary 3.8(2) to  $SOP_2$ . The following theorem is a step in this direction, showing a local version. See also Discussion 3.17.

**Theorem 3.11.** *If  $T$  has  $SOP_2$  as exemplified by  $\vartheta(\bar{x}; \bar{y})$ , then  $(T_{tr}^*, \vartheta_{tr}(x; y)) \triangleleft_{\lambda}^* (T, \vartheta(\bar{x}; \bar{y}))$  for any  $\lambda \geq |T| + \aleph_0$  regular.*

*Proof.* We can find a model  $M_1$  of  $T_{tr}^*$  and model  $M_2$  of  $T$  and  $\bar{a}_b \in {}^{\ell g(\bar{y})}M_2$  for  $b \in M_1$  such that:

- ( $\alpha$ ) if  $M_1 \models b_0 < \dots < b_{n-1}$  then  $\{\vartheta(\bar{x}, \bar{a}_{b_\ell}) : \ell < n\}$  is satisfiable in  $M_2$
- ( $\beta$ ) if  $b_1, b_2$  are incomparable in  $M_1$  then

$$M_2 \models \neg(\exists \bar{x})(\vartheta(\bar{x}, \bar{a}_{b_1}) \& \vartheta(\bar{x}, \bar{a}_{b_2}))$$

- ( $\gamma$ ) for no  $\bar{d} \in {}^{\ell g(\bar{x})}(M_2)$  is  $\{b \in M_1 : M_2 \models \vartheta(\bar{d}, \bar{a}_b)\}$  unbounded in  $M_1$  (note that by ( $\beta$ ) it is always linearly ordered in  $M_1$ , therefore ( $\gamma$ ) means that for each  $\bar{d} \in {}^{\ell g(\bar{x})}(M_2)$ , there exists an element of  $M_1$  which is above every  $b$  satisfying  $\vartheta(\bar{d}, \bar{a}_b)$ ).

[The construction of  $M_1$  and  $M_2$  is as follows: choose by induction on  $n$ ,  $(M_{1,n}, M_{2,n}, \langle \bar{a}_b : b \in M_{1,n} \rangle : n < \omega)$  such that:

- (a)  $M_{1,n}$  is a model of  $T_{tr}^*$
- (b)  $M_{2,n}$  is a model of  $T$
- (c)  $M_{1,n} \prec M_{1,n+1}$  moreover, every branch of  $M_{1,n}$  has an upper bound in  $M_{1,n+1}$
- (d)  $M_{2,n} \prec M_{2,n+1}$
- (e)  $\bar{a}_b \in {}^{\ell g(\bar{y})}(M_{2,n})$  for  $b \in M_{1,n}$
- (f) clauses  $(\alpha), (\beta)$  hold
- (g) if  $b \in M_{1,n+1}$  and  $[b' \in M_{1,n} \Rightarrow M_{1,n+1} \models \neg(b < b')]$  then  $\vartheta(\bar{x}, \bar{a}_b)$  is not satisfied by any sequence from  $M_{1,n}$ .

There is no problem to carry the definition.

Now  $M_1 = \bigcup_n M_{1,n}$ ,  $M_2 = \bigcup_n M_{2,n}$  and  $\langle \bar{a}_b : b \in M_1 \rangle$  are as required above.]

Now let  $\chi$  be such that  $M_1, M_2 \in \mathcal{H}(\chi)$ , wlog  $\tau_T = \tau(M_2)$ ,  $\{<\} = \tau(T_{tr}) = \tau(M_1)$  and  $\{\in\}$  are pairwise disjoint. Now we define a model  $\mathcal{B}_0$ .

Its universe is  $\mathcal{H}(\chi)$  relation  $\in$  (membership)

$$P_1 = |M_1|,$$

$$P_2 = |M_2|$$

$R = R^{M_\ell}$  if  $R \in \tau(M_\ell)$ ,  $\ell \in \{1, 2\}$   $F_\ell$  (for  $\ell < \ell g(\bar{y})$ ) a partial unary function such that:  $b \in M_1 \Rightarrow \langle F_\ell(b) : \ell < \ell g(\bar{y}) \rangle = \bar{a}_b$ .

Let  $T^* = Th(\mathcal{B}_0)$ . For the obvious  $\bar{\varphi}$  and  $\bar{\psi}$ ,  $T^*$  is  $(T, T_{tr})$ -superior and  $|T^*| = |T| + \aleph_0$ . Assume  $\lambda = \text{cf}(\lambda) > |T^*|$ .

So let  $\mathcal{B}$  be a model of  $T^*$  such that  $M'_2 = \mathcal{B}^{[\bar{\varphi}]}$ , the model of  $T$  interpreted in it, is  $\lambda^+$ -saturated. It will be enough to prove that  $M'_1 = \mathcal{B}^{[\bar{\psi}]}$  satisfies: for every branch of cofinality  $\theta \leq \lambda$  there exists an

upper bound. So let  $\{b_i : i < \theta\}$  be  $<^{M_1}$ -increasing let  $\bar{c}_i = \langle F_\ell^{\mathcal{B}}(b_\ell) : \ell < \ell g(\bar{y}) \rangle$ . Hence for any  $n < \omega, i_0 < \dots < i_{n-1} < \theta$  we have  $M'_2 \models (\exists \bar{x})[\bigwedge_{m < n} \vartheta(\bar{x}, \bar{c}_i)]$  because  $\mathcal{B}_0 \models (\forall z_0, \dots, z_{n-1})[\bigwedge_{k < n} P_1(z_k) \Rightarrow z_0 < z_1 < \dots < z_{n-1} \rightarrow (\exists \bar{x}) \bigwedge_{m < n} \vartheta(\bar{x}, \langle F_\ell(z_m) : \ell < \ell g(\bar{y}) \rangle)]$ .

So  $\{\vartheta(\bar{x}, \bar{c}_i) : i < \theta\}$  is finitely satisfiable in  $M'_2$  hence some  $\bar{d} \in {}^{\ell g(\bar{x})}(M'_2)$  realizes it. Now we claim that  $\{b \in M'_1 : \mathcal{B} \models \vartheta(\bar{d}, \bar{a}_b)\}$  is bounded in  $M'_1$ : recall that by clause  $(\gamma)$   $\mathcal{B}_0$  satisfies: for every  $\bar{x} \in {}^{\ell g(y)}P_2$  there exists  $z \in P_1$  such that  $z$  is  $<^{\mathcal{B}}$ -above all the elements  $w \in P_1$  satisfying  $\vartheta(\bar{x}, \bar{a}_w)$ . Therefore  $\mathcal{B}$  satisfies this sentence, and applying it to  $\bar{d} \in {}^{\ell g(\bar{x})}(M'_2)$ , we get  $b^* \in M'_1$  - the required bound. As for each  $i < \theta$ ,  $\vartheta(\bar{d}, \bar{a}_{b_i})$  holds, clearly  $\mathcal{B} \models b_i < b^*$  for all  $i$ , and we are done.  $\square$

The next goal is to complete the proof started in [DžSh692] of the fact that  $\triangleleft^*$ -maximality implies  $SOP_2$ . In [DžSh692] a property “ $\triangleleft_{\lambda}^{**}$ -maximality”, which is closely related to “ $\triangleleft_{\lambda}^*$ -maximality” was defined, and it was shown (Theorem 3.6 there) that every  $T$  which is  $\triangleleft_{\lambda}^{**}$ -maximal for some (every) big enough regular  $\lambda$ , has an order property similar to  $SOP_2$ , that we call  $SOP_2''$  (see Definition 1.5). We will show that  $SOP_2''$  is equivalent to  $SOP_2$  (for a theory). This answers Question (3.8)(3) from the original version of [DžSh692] (version 1 on the arXiv).

*Discussion 3.12.* In particular, Theorem 3.13 will lead to the following conclusion: assuming that  $T$  is  $\triangleleft_{\lambda^+}^*$ -maximal for some regular  $\lambda$  satisfying  $2^\lambda = \lambda^+$ , we get by [DžSh692], Claim (3.2) that  $T$  is  $\triangleleft_{\lambda}^{**}$ -maximal, so it has  $SOP_2''$ , and therefore  $SOP_2$ . So we will obtain  $\triangleleft^*$ -maximality implies  $SOP_2$ , see Corollary 3.15.

**Theorem 3.13.** *Let  $T$  be a theory.*

- (1) *Suppose  $\vartheta(\bar{x}, \bar{y})$  exemplifies  $SOP_2$  in  $T$ . Then  $\vartheta(\bar{x}; \bar{y})$  exemplifies  $SOP_2'$  in  $T$  as well.*
- (2) *Suppose  $\vartheta(\bar{x}, \bar{y})$  exemplifies  $SOP_2'$  in  $T$ . Then for some  $k$ ,  $\vartheta^{<k>}(\bar{x}; \bar{y})$  exemplifies  $SOP_2$  in  $T$  (where  $\vartheta^{<k>}(\bar{x}; \bar{y}^{<k>}) = \bigwedge_{\ell < k} \vartheta(\bar{x}; \bar{y}_\ell)$ ).*

*Proof.* (1) is easy.

- (2) Denote  $\mathcal{J}_\lambda^n = \{\bar{\eta} : \bar{\eta} = \langle \eta_\ell : \ell \leq n \rangle, \eta_\ell \triangleleft \eta_{\ell+1}; \text{ and } \eta_\ell \in {}^{\lambda>}2\}$ . So assume  $\vartheta(\bar{x}; \bar{y})$  has  $SOP_2''$  as exemplified by  $n$ ,  $\bar{\mathbf{a}} = \langle a_{\bar{\eta}} : \bar{\eta} \in \mathcal{J}_\omega^n \rangle$ . Without loss of generality  $\langle \bar{a}_{\bar{\eta}} : \bar{\eta} \in \mathcal{J}_\omega^n \rangle$  is tree indiscernible in the relevant sense:  $\eta \frown \langle 0 \rangle, \eta \frown \langle 1 \rangle$  look the same over  $\eta$  (2 – *fbti* from 1.9). We can assume this by 1.10 (for more details, see [DžSh692], Claim (2.14)).

For  $\nu \in {}^{\omega \geq 2}$  let  $p_\nu = \{\vartheta(\bar{x}, \bar{a}_{\bar{\eta}}) : \bar{\eta} = \langle \eta_\ell : \ell < n \rangle, \eta_\ell < \eta_{\ell+1} \trianglelefteq \nu\}$  so

$\otimes_1$   $p_\eta$  for  $\eta \in {}^\omega 2$  is consistent (in  $\mathfrak{C}_T$ ).

Let

$\Xi = \{(h, \Upsilon) : h \text{ is a one-to-one mapping from } {}^{n \geq m} \text{ to } {}^{\omega > 2}$

preserving  $\triangleleft, \perp$  and  $\Upsilon \subseteq {}^n m$  and there is

$\langle \nu_\eta^* : \eta \in \Upsilon \rangle, h(\eta) \triangleleft \nu_\eta^* \in {}^\omega 2 \text{ for } \eta \in {}^n m$

such that  $\cup \{p_{\nu_\eta^*} : \eta \in \Upsilon\}$  is inconsistent

Now

$\otimes_2$   $\Xi$  is nonempty

[By the definition of  $SOP_2''$ , clause (b), choose  $\Upsilon = {}^n m$  ]

Choose  $(h^*, \Upsilon^*) \in \Xi$  with  $|\Upsilon^*|$  of minimal cardinality and  $\langle \nu_\eta^* : \eta \in \Upsilon^* \rangle$  as there. By  $\otimes_1$  clearly  $|\Upsilon^*| \geq 2$ . So choose  $\eta_0 \neq \eta_1$  from  $\Upsilon^*$  with  $\nu^* = \nu_{\eta_0}^* \cap \nu_{\eta_1}^*$  ( $= h(\eta_0) \cap h(\eta_1)$ ) being of maximal length and let  $k^* = \ell g(\nu^*)$ . We can find  $\ell^* < \omega$  sufficiently large such that  $\cup \{p_{\nu_\eta^* \upharpoonright \ell^*} : \eta \in \Upsilon^*\}$  is inconsistent. We choose by induction on  $i < \omega$  for every  $\rho \in {}^\omega 2$ , a sequence  $\nu_\rho \in {}^{\omega > 2}$  by  $\nu_{\langle \rangle} = \langle \rangle, \nu_{\rho \hat{< j}} = \nu_\rho \hat{<} (\nu_{\eta_j}^* \upharpoonright \ell^*)$ .

Lastly for  $\rho \in {}^{\omega > 2} \notin \{< >\}$  let  $\vartheta^*(\bar{x}, \bar{b}_\rho^*)$  be the conjunction of

$$\bigcup \{p_{\nu_\eta^* \upharpoonright \ell^*} : \eta \in \Upsilon^* \setminus \{\eta_0, \eta_1\}\} \cup \{\vartheta(\bar{x}, \bar{a}_{\bar{\eta}}) : \bar{\eta} = \langle \eta_\ell : \ell \leq n \rangle,$$

$$\eta_\ell \triangleleft \eta_{\ell+1} \trianglelefteq \nu_\rho \text{ and } (\forall \ell \leq n)[\ell g(\eta_\ell) \notin [k^*, \ell g(\nu_\rho) - \ell^* + k^*]]$$

(the last condition is empty if  $\ell g(\rho) = 1$ )

In other words, we are taking the “upper part” of  $\nu_\rho$  that “looks like”  $\eta_0$  or  $\eta_1$  after they split.

Now if  $\rho^* \in {}^\omega 2$  then  $\{\vartheta^*(x, \bar{b}_\rho) : \rho \triangleleft \rho^*\}$  is consistent as all its members are conjunctions of formulas from

$$\cup \{p_{\nu_\eta^*} : \eta \in \Upsilon^* \setminus \{\eta_0^*, \eta_1^*\}\} \cup p_{\rho^*}$$

and this is consistent as otherwise  $(h^* \upharpoonright (\Upsilon^* \setminus \{\eta_0^*, \eta_1^*\})) \cup \{\langle \eta_0^*, \rho^* \upharpoonright \ell^{**} \rangle, \Upsilon^* \setminus \{\eta_1^*\}\}$  belongs to  $\Xi$  for some  $\ell^{**}$ , thus contradicting the choice of  $(h^*, \Upsilon^*)$ , i.e. with minimal  $|\Upsilon^*|$ .

Lastly if  $\rho_0, \rho_1 \in {}^{\omega > 2}$  are  $\triangleleft$ -incomparable then  $\{\vartheta^*(\bar{x}; \bar{b}_{\rho_0}), \vartheta^*(\bar{x}; \bar{b}_{\rho_1})\}$  is inconsistent: we know that

⊗ 3.13.1.

$$\bigcup \{p_{\nu_{\bar{\eta}}^* \upharpoonright \ell^*} : \eta \in \Upsilon^* \setminus \{\eta_0, \eta_1\}\} \cup \{\vartheta(\bar{x}, \bar{a}_{\bar{\eta}}) : \bar{\eta} = \langle \eta_\ell : \ell \leq n \rangle, \\ \eta_\ell \triangleleft \eta_{\ell+1} \trianglelefteq \nu_{\eta_0}^* \upharpoonright \ell^*\} \cup \{\vartheta(\bar{x}, \bar{a}_{\bar{\eta}}) : \bar{\eta} = \langle \eta_\ell : \ell \leq n \rangle, \eta_\ell \triangleleft \eta_{\ell+1} \trianglelefteq \nu_{\eta_1}^* \upharpoonright \ell^*\}.$$

is inconsistent (by the choice of  $(h^*, \Upsilon^*) \in \Xi$  and the choice of  $\ell^*$ ). Now, by the fact that  $\nu^* = \nu_{\eta_0}^* \cap \nu_{\eta_1}^*$  was chosen to be maximal among other pairs in  $\Upsilon^*$ , we see that if

$$\bar{\eta}_0 = \langle \eta_\ell^0 : \ell \leq n \rangle, \text{ where for each } \ell, \eta_\ell^0 \triangleleft \eta_{\ell+1}^0 \trianglelefteq \nu_{\eta_0}^* \upharpoonright \ell^*$$

and

$$\bar{\eta}_1 = \langle \eta_\ell^1 : \ell \leq n \rangle, \text{ where for each } \ell, \eta_\ell^1 \triangleleft \eta_{\ell+1}^1 \trianglelefteq \nu_{\eta_1}^* \upharpoonright \ell^*$$

while

$$\bar{\eta}_3 = \langle \eta_\ell^3 : \ell \leq n \rangle, \text{ where for each } \ell, \eta_\ell^3 \triangleleft \nu_{\eta^*}^* \text{ for some } \eta^* \in \Upsilon^* \setminus \{\eta_0^*, \eta_1^*\}$$

then

⊗ 3.13.2.

$$\bar{\eta}_1 \frown \bar{\eta}_2 \frown \bar{\eta}_3 \equiv \bar{\varsigma}_1 \frown \bar{\varsigma}_2 \frown \bar{\eta}_3$$

where  $\bar{\varsigma}_j = \langle \varsigma_\ell^j : \ell \leq n \rangle$  and

$$\varsigma_\ell^j = \eta_\ell^j, \text{ if } \text{lg}(\eta_\ell^j) \leq k^*$$

$$\varsigma_\ell^j = \nu_{\rho_j} \upharpoonright [\text{lg}(\nu_{\rho_j}) - (\ell^* - \text{lg}(\eta_\ell^j))], \text{ otherwise}$$

In simpler words: we replace every  $\eta_\ell^j$  (an initial segment of  $\nu_{\eta_j} \upharpoonright \ell^*$ ) whose length is bigger than  $k^*$  (in particular, it is not below any element in the image of  $\Upsilon^*$  other than  $\nu_{\eta_j}$  itself) by an appropriate initial segment of  $\nu_{\rho_j}$ , and get a similar sequence over the image of  $\Upsilon^* \setminus \{\eta_0^*, \eta_1^*\}$ .

Now, by indiscernibility of  $\langle \bar{a}_{\bar{e}t\alpha} \rangle$ , the definition of  $\vartheta^*(\bar{x}, \bar{b}_\rho^*)$ , 3.13.1 and 3.13.2, we conclude  $\{\vartheta^*(\bar{x}; \bar{b}_{\rho_0}), \vartheta^*(\bar{x}; \bar{b}_{\rho_1})\}$  is also inconsistent.

□

Let us summarize the main results of this section.

- Definition 3.14.** (1) We call a theory  $T$   $\triangleleft^*$ -maximal if it is  $\triangleleft_\lambda^*$ -maximal for every regular  $\lambda > |T| + \aleph_0$ .
- (2) We call a pair of theory and formula  $(T, \vartheta)$   $\triangleleft^*$ -maximal if it is  $\triangleleft_\lambda^*$ -maximal for every regular  $\lambda > |T| + \aleph_0$ .

- Corollary 3.15.** (1) If  $T$  has  $SOP_3$  then it is  $\triangleleft^*$ -maximal.
- (2) If  $T$  is  $\triangleleft^*$ -maximal then it has  $SOP_2$ .

*Proof.* (1) Corollary 3.8.

- (2) By [DžSh692] Claim 3.2, [DžSh692] Theorem 3.6 and Theorem 3.13 above.

□

So we have shown  $SOP_3 \implies \triangleleft^*$ -maximality  $\implies SOP_2$ , and for the second implication we also have a weak (local) “converse”, Theorem 3.11. See Discussion 3.17 below.

*Question 3.16.* Is any of the two implications above reversible?

*Discussion 3.17.* Note that Theorem 3.11 is a step in (possibly) reversing the second implication above: we show that if  $T$  has  $SOP_2$  exemplified by a formula  $\vartheta$ , then the pair  $(T_{tr}^*, \vartheta_{tr})$  is  $\triangleleft^*$ -below the pair  $(T, \vartheta)$ . By Theorem 3.6 (and quantifier elimination), in order to obtain  $SOP_2 \implies \triangleleft^*$ -maximality it is enough to show that  $(T_{tr}^*, \{\vartheta_{tr}, \neg\vartheta_{tr}\})$  is  $\triangleleft^*$ -below  $(T, \Delta)$  where  $\Delta$  is some fragment of the language of  $T$ . This was our original motivation for proving Theorem 3.11, which is in a sense a “local” or “positive” version of what we are interested in, but right now it is unclear to us whether similar techniques will lead to the desired “global” result.

One should remark that Theorem 3.11 is not weaker than the global version since  $\Delta = \{\vartheta\}$ , so it is really localized to a single formula, with no use of negation (hence “positive”). Therefore, although it does not quite do what we one would hope for, we find Theorem 3.11 interesting on its own.

*Discussion 3.18.* We would also like to point out that our analysis provides an alternative (in fact, in a sense a more conceptual) proof of Theorem 1.17 and Conclusion 1.18 in [DžSh692]. Theorem 3.5 here shows that  $T_{tr}^*$ , and therefore  $T_{ord}^*$  is maximal in  $\triangleleft_\lambda^*$ , and therefore is  $\triangleleft_\lambda^*$ -above  $T_{feq}^*$ . By Theorem 2.1,  $T_{feq}^*$  does not have  $SOP_1$  (in particular, does not have  $SOP_2$ ), and so by Corollary 3.15, can not be  $\triangleleft_\lambda^*$ -maximal. So  $T_{feq}^*$  is *strictly below*  $T_{tr}^*$  (and  $T_{ord}^*$ ) in  $\triangleleft_\lambda^*$ -ordering, which is precisely the statement of Conclusion 1.18 in [DžSh692]. There is no surprise here: in [DžSh692] it was shown that  $T_{feq}^*$  is on the “good”

side of an “external” dividing line. Here we showed that it is also on the “good” side of an “internal syntactic” dividing line, and brought the “internal” and “external” lines close together. So our paper also provides a generalization of section 1 of [DžSh692].

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