Stochastic Iterative Greedy Algorithms for Sparse Reconstruction

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Support from Alfred P. Sloan Foundation and NSF CAREER #1348721



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• Matrix recovery:

$$y_i = \langle A_i, X^\star \rangle + e_i$$

 $\min_{X \in \mathbb{R}^{n_1 \times n_2}} \frac{1}{2m} \|y - \mathcal{A}(X)\|_2^2 \quad \text{subject to} \quad \operatorname{rank}(X) \le k,$

• Generalized sparse recovery:

$$x = \sum_{i=1}^{k} \alpha_i d_i, \qquad d_i \in \mathcal{D},$$

$$||x||_{0,\mathcal{D}} = \min_{k} \{k : x = \sum_{i \in T} \alpha_i d_i \quad \text{with} \quad |T| = k\}$$



Methods

Compressed sensing / matrix recovery

- L1-minimization & nuclear norm minimization
- Iterative methods (IHT, CoSaMP, OMP, …)
- Optimization
 - (Stochastic) gradient descent
 - (Stochastic) coordinate descent
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Assumptions

Definition 1 (\mathcal{D} -restricted strong convexity (\mathcal{D} -RSC)). The function F(x) satisfies the \mathcal{D} -RSC if there exists a positive constant ρ_k^- such that

$$F(x') - F(x) - \left\langle \nabla F(x), x' - x \right\rangle \ge \frac{\rho_k}{2} \left\| x' - x \right\|_2^2, \tag{16}$$

for all vectors x and x' of size n such that $|\operatorname{supp}_{\mathcal{D}}(x) \cup \operatorname{supp}_{\mathcal{D}}(x')| \leq k$.

Definition 2 (\mathcal{D} -restricted strong smoothness (\mathcal{D} -RSS)). The function $f_i(x)$ satisfies the \mathcal{D} -RSS if there exists a positive constant $\rho_k^+(i)$ such that

$$\left\|\nabla f_{i}(x') - \nabla f_{i}(x)\right\|_{2} \le \rho_{k}^{+}(i) \left\|x' - x\right\|_{2}$$
(17)

for all vectors x and x' of size n such that $|\operatorname{supp}_{\mathcal{D}}(x) \cup \operatorname{supp}_{\mathcal{D}}(x')| \leq k$.

We may wish to consider blocks of the matrix A and break F(x) into functionals corresponding to each block i. Call the number of blocks b.

Restricted Isometry Property

 $(1-\delta) \|x\|_2^2 \le \|Ax\|_2^2 \le (1+\delta) \|x\|_2^2$ for all k-sparse vectors x,

In particular, we require that the design matrix A satisfies

$$\frac{1}{m} \|Ax\|_2^2 \ge (1 - \delta_k) \|x\|_2^2 \quad \frac{1}{b} \|A_{b_i}x\|_2^2 \le (1 + \delta_k) \|x\|_2^2$$

Here, $(1 + \delta_k)$ and $(1 - \delta_k)$ with $\delta_k \in (0, 1]$ play the role of $\rho_k^+(i)$ and ρ_k^-

Stochastic greedy methods

Define $\operatorname{approx}_k(x,\eta)$ as the operator that constructs a set Γ of cardinality k such that

$$\left\|\mathcal{P}_{\Gamma}x - x\right\|_{2} \le \eta \left\|x - x_{k}\right\|_{2},$$

Algorithm 1 StoIHT algorithm

 $\begin{array}{ll} \textbf{input:} \ k, \ \gamma, \ \eta, \ p(i), \ \text{and stopping criterion} \\ \textbf{initialize:} \ x^0 \ \text{and} \ t = 0 \\ \textbf{repeat} \\ \textbf{randomize:} \ \ \text{select an index} \ i_t \ \text{from} \ [M] \ \text{with probability} \ p(i_t) \\ \textbf{proxy:} \ \ b^t = x^t - \frac{\gamma}{Mp(i_t)} \nabla f_{i_t}(x^t) \\ \textbf{identify:} \ \ \Gamma^t = \text{approx}_k(b^t, \eta) \\ \textbf{estimate:} \ \ x^{t+1} = \mathcal{P}_{\Gamma^t}(b^t) \\ \ t = t + 1 \\ \textbf{until halting criterion} \ true \\ \textbf{output:} \ \hat{x} = x^t \end{array}$

Stochastic greedy methods

Define $\operatorname{approx}_k(x,\eta)$ as the operator that constructs a set Γ of cardinality k such that

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Algorithm 2 StoGradMP algorithm **input:** $k, \eta_1, \eta_2, p(i)$, and stopping criterion initialize: x^0 , $\Lambda = 0$, and t = 0repeat **randomize:** select an index i_t from [M] with probability $p(i_t)$ $r^t = \nabla f_{i_t}(x^t)$ proxy: identify: $\Gamma = \operatorname{approx}_{2k}(r^t, \eta_1)$ merge: $\widehat{\Gamma} = \overline{\Gamma} \cup \Lambda$ estimate: $b^t = \operatorname{argmin}_w F(x) \quad w \in \operatorname{span}(D_{\widehat{\Gamma}})$ **prune:** $\Lambda = \operatorname{approx}_k(b^t, \eta_2)$ update: $x^{t+1} = \mathcal{P}_{\Lambda}(b^t)$ t = t+1until halting criterion true output: $\hat{x} = x^t$

Theoretical guarantees : StoIHT

Theorem 1. Let x^* be a feasible solution of (5) and x^0 be the initial solution. At the (t+1)-th iteration of Algorithm 1, the expectation of the recovery error is bounded by

$$\mathbb{E} \|x^{t+1} - x^{\star}\|_{2} \le \kappa^{t+1} \|x^{0} - x^{\star}\|_{2} + \frac{\sigma_{x^{\star}}}{(1-\kappa)}$$
(21)

Sparse signal recovery:

$$\mathbb{E} \left\| x^{t+1} - x^{\star} \right\|_{2} \le (3/4)^{t+1} \left\| x^{\star} \right\|_{2} + c \sqrt{\frac{\sigma^{2} k_{0} \log n}{b}}.$$

Low-rank matrix recovery:

$$\mathbb{E} \left\| X^{t+1} - X^{\star} \right\|_{F} \le (3/4)^{t+1} \left\| X^{\star} \right\|_{F} + c \left(\sqrt{\frac{\sigma^{2} k n}{b}} + \sqrt{\frac{(\eta^{2} - 1)\sigma^{2} n^{2}}{b}} \right).$$



Figure 1: Sparse Vector Recovery: Percent recovery as a function of the number of measurements for IHT (left) and StoIHT (right) for various sparsity levels k_0 .



Figure 2: Sparse Vector Recovery: Percent recovery as a function of the number of measurements for GradMP (left) and StoGradMP (right) for various sparsity levels k_0 .



Figure 3: Sparse Vector Recovery: Recovery error as a function of epochs and various block sizes b for IHT methods (left) and GradMP methods (right).



Figure 4: Sparse Vector Recovery: Number of measurements required for signal recovery as a function of block size (blue marker) for StoIHT (left) and StoGradMP (right). Number of measurements required for deterministic method shown as red solid line.



Figure 5: Sparse Vector Recovery: A comparison of IHT and StoIHT in the presence of noise. Recovery error versus epoch (left) and measurements required versus block size (right).



Figure 6: Sparse Vector Recovery: A comparison of GradMP and StoGradMP in the presence of noise. Recovery error versus epoch (left) and measurements required versus block size (right).



Figure 7: Sparse Vector Recovery: (Left) A comparison of StoIHT for various values of the step size γ (shown in the colorbar). (Right) A comparison of the mean recovery error as a function of runtime for IHT, StoIHT, GradMP, and StoGradMP.



Figure 8: Low-Rank Matrix Recovery: Percent recovery as a function of the number of measurements for IHT (left) and StoIHT (right) for various rank levels k_0 .



Figure 9: Low-Rank Matrix Recovery: Percent recovery as a function of the number of measurements for GradMP (left) and StoGradMP (right) for various rank levels k_0 .



Figure 10: Low-Rank Matrix Recovery: Recovery error as a function of the number of epochs for the IHT methods using m = 90 (left) and m = 140 (right), shown for various block sizes b used in StoIHT.



Figure 11: Low-Rank Matrix Recovery: Recovery error as a function of the number of epochs for the GradMP methods using m = 90 (left) and m = 140 (right), shown for various block sizes b used in StoGradMP.



Figure 12: Low-Rank Matrix Recovery: Number of measurements required for signal recovery as a function of block size (blue marker) for StoIHT (left) and StoGradMP (right). Number of measurements required for deterministic method shown as red solid line.



Figure 13: Low-Rank Matrix Recovery with Approximations: Mean recovery error as a function of epochs (left) and runtime (right) for various over-sampling factors d using the StoIHT algorithm. Performance using full SVD computation shown as dashed line.



Figure 14: Low-Rank Matrix Recovery with Approximations: Mean recovery error as a function of epochs (left) and runtime (right) for various over-sampling factors d using the StoGradMP algorithm. Performance using full SVD computation shown as dashed line.



For more info:

- dneedell@cmc.edu
- www.cmc.edu/pages/faculty/DNeedell
- Linear Convergence of Stochastic Iterative Greedy Algorithms with Sparse Constraints by N. Nguyen, D. Needell, and T. Woolf. Submitted (Arxiv arxiv.org/abs/1407.0088)