## Exponential decay of reconstruction error from binary measurements of sparse signals

## Deanna Needell



Joint work with R. Baraniuk, S. Foucart, Y. Plan, and M. Wootters

## Outline

$\diamond$ Introduction
$\triangleleft$ Mathematical Formulation \& Methods
$\diamond$ Practical CS
$\diamond$ Other notions of sparsity
$\diamond$ Heavy quantization
↔ Adaptive sampling

## The mathematical problem

1. Signal of interest $f \in \mathbb{C}^{d}\left(=\mathbb{C}^{N \times N}\right)$
2. Measurement operator $\mathscr{A}: \mathbb{C}^{d} \rightarrow \mathbb{C}^{m}(m \ll d)$
3. Measurements $y=\mathscr{A} f+\xi$

$$
[y]=[\quad \mathscr{A} \quad] f]+[\xi]
$$

4. Problem: Reconstruct signal from measurements $y$

## Sparsity

Measurements $y=\mathscr{A} f+\xi$.



Assume $f$ is sparse:
$\diamond$ In the coordinate basis: $\|f\|_{0} \stackrel{\text { def }}{=}|\operatorname{supp}(f)| \leq s \ll d$
$\triangleleft$ In orthonormal basis: $f=B x$ where $\|x\|_{0} \leq s \ll d$

In practice, we encounter compressible signals.
$\checkmark \quad f_{s}$ is the best $s$-sparse approximation to $f$

## Many applications...

« Radar, Error Correction
$\diamond$ Computational Biology, Geophysical Data Analysis
$\triangleleft$ Data Mining, classification
$\diamond$ Neuroscience
$\diamond$ Imaging
$\diamond$ Sparse channel estimation, sparse initial state estimation
$\widehat{\psi}$ Topology identification of interconnected systems
ヶ ...

## Sparsity...

Sparsity in coordinate basis: $f=x$


## Reconstructing the signal $f$ from measurements $y$

$\checkmark \ell_{1}$-minimization [Candès-Romberg-Tao]
Let $A$ satisfy the Restricted Isometry Property and set:

$$
\hat{f}=\underset{g}{\operatorname{argmin}}\|g\|_{1} \quad \text { such that } \quad\|\mathscr{A} f-y\|_{2} \leq \varepsilon
$$

where $\|\xi\|_{2} \leq \varepsilon$. Then we can stably recover the signal $f$ :

$$
\|f-\hat{f}\|_{2} \lesssim \varepsilon+\frac{\left\|x-x_{s}\right\|_{1}}{\sqrt{s}}
$$

This error bound is optimal.

## Restricted Isometry Property

$\diamond \mathscr{A}$ satisfies the Restricted Isometry Property (RIP) when there is $\delta<c$ such that

$$
(1-\delta)\|f\|_{2} \leq\|\mathscr{A} f\|_{2} \leq(1+\delta)\|f\|_{2} \quad \text { whenever }\|f\|_{0} \leq s
$$

$\diamond m \times d$ Gaussian or Bernoulli measurement matrices satisfy the RIP with high probability when

$$
m \gtrsim s \log d
$$

$\diamond$ Random Fourier and others with fast multiply have similar property: $m \gtrsim s \log ^{4} d$.

## Other recovery methods

Greedy Algorithms
$\diamond$ If $A$ satisfies the RIP, then $A^{*} A$ is "close" to the identity on sparse vectors
$\stackrel{\rightharpoonup}{ }$ Use proxy $p=A^{*} y=A^{*} A x \approx x$
$\triangleleft$ Threshold to maintain sparsity: $\hat{x}=H_{s}(p)$
$\diamond$ Repeat
« (Iterative Hard Thresholding)

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\|\mathbf{x}-\Delta(\mathbf{y})\| \leq \gamma
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provided the oversampling factor satisfies

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for $f$ slowly increasing when $\gamma$ decreases to zero, equivalently

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\|\mathbf{x}-\Delta(\mathbf{y})\| \leq g(\lambda)
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for $g$ rapidly decreasing to zero when $\lambda$ increases.

A visual


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\left\|\mathbf{x}-\frac{\Delta_{\mathrm{LP}}(\mathbf{y})}{\left\|\Delta_{\mathrm{LP}}(\mathbf{y})\right\|_{2}}\right\|_{2} \lesssim \lambda^{-1 / 5} \quad \text { whenever }\|\mathbf{x}\|_{0} \leq \boldsymbol{s},\|\mathbf{x}\|_{2}=1
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- Uniform, nonadaptive, adversarial quantization error: If $\mathbf{A} \in \mathbb{R}^{m \times n}$ is a Gaussian matrix, then w/hp

$$
\left\|\mathbf{x}-\Delta_{\mathrm{SOCP}}(\mathbf{y})\right\|_{2} \lesssim \lambda^{-1 / 12} \quad \text { whenever }\|\mathbf{x}\|_{0} \leq \boldsymbol{s},\|\mathbf{x}\|_{2}=1
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- Power decay is optimal since

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\left\|\mathbf{x}-\Delta_{\mathrm{opt}}(\mathbf{y})\right\|_{2} \gtrsim \lambda^{-1}
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- Remedy: adaptive choice of dithers $\tau_{1}, \ldots, \tau_{m}$ in

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\mathbf{x}^{t+1}=\quad \mathbf{x}^{t}+\widehat{\mathbf{x}-\mathbf{x}^{t}}, \quad \text { so that } \quad\left\|\mathbf{x}-\mathbf{x}^{t+1}\right\|_{2} \leq R / 4^{t+1}
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- After $T$ iterations, number of measurements is $m=q T$, and

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- Software step needed to compute the thresholds $\tau_{i}=\left\langle\mathbf{a}_{i}, \mathbf{x}^{t}\right\rangle$.


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- Cons: slow, post-quantization error not handled.


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- Plane geometry to estimate direction and magnitude of $\mathbf{x}$.
- Cons: dithers $\left\langle\mathbf{a}_{i}, \mathbf{w}\right\rangle$ are adaptive.
- Pros: deterministic, fast, handles pre/post-quantization errors.


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- If $\|\mathbf{e}\|_{\infty} \leq \varepsilon R 2^{-T}$ (or $\left\|\mathbf{e}^{t}\right\|_{2} \leq \varepsilon \sqrt{q}\left\|\mathbf{x}-\mathbf{x}^{t}\right\|_{2}$ throughout), then

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for the convex-optimization and hard-thresholding schemes.

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- Post-quantization error $\mathbf{f} \in\{ \pm 1\}^{m}$ in

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- Pre-quantization error $\mathbf{e} \in \mathbb{R}^{m}$ in

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y_{i}=\operatorname{sign}\left(\left\langle\mathbf{a}_{i}, \mathbf{x}\right\rangle-\tau_{i}+e_{i}\right) .
$$

- If $\|\mathbf{e}\|_{\infty} \leq \varepsilon R 2^{-T}$ (or $\left\|\mathbf{e}^{t}\right\|_{2} \leq \varepsilon \sqrt{q}\left\|\mathbf{x}-\mathbf{x}^{t}\right\|_{2}$ throughout), then

$$
\left\|\mathbf{x}-\mathbf{x}^{T}\right\|_{2} \leq R 2^{-T}=R \exp (-c \lambda)
$$

for the convex-optimization and hard-thresholding schemes.

- Post-quantization error $\mathbf{f} \in\{ \pm 1\}^{m}$ in

$$
y_{i}=f_{i} \operatorname{sign}\left(\left\langle\mathbf{a}_{i}, \mathbf{x}\right\rangle-\tau_{i}\right) .
$$

- If $\operatorname{card}\left(\left\{i: f_{i}^{t}=-1\right\}\right) \leq \eta q$ throughout, then

$$
\left\|\mathbf{x}-\mathbf{x}^{T}\right\|_{2} \leq R 2^{-T}=R \exp (-c \lambda)
$$

for the hard-thresholding scheme.

## Numerical Illustration




## Numerical Illustration, ctd

500 hard-thresholding-based tests


100 second-order-cone-programming-based tests $\mathrm{n}=100, \mathrm{~m} / \mathrm{s}=20: 20: 200, \mathrm{~T}=5$ (11 hours)


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\left|\frac{\sqrt{\pi / 2}}{q}\langle\mathbf{A} \mathbf{w}, \operatorname{sign}(\mathbf{A} \mathbf{x})\rangle-\langle\mathbf{w}, \mathbf{x}\rangle\right| \leq \delta
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$$
\mathbf{e}=\mathbf{A u} \quad \text { with } \quad \begin{cases}\|\mathbf{u}\|_{2} & \leq d\|\mathbf{e}\|_{2} / \sqrt{q} \\ \|\mathbf{u}\|_{1} & \leq d^{\prime} \sqrt{s_{*}}\|\mathbf{e}\|_{2} / \sqrt{q}\end{cases}
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- Restricted Isometry Property: if $q \geq C \delta^{-2} s \ln (n / s)$, then with w/hp

$$
\left|\frac{1}{q}\|\mathbf{A} \mathbf{x}\|_{2}^{2}-\|\mathbf{x}\|_{2}^{2}\right| \leq \delta\|\mathbf{x}\|_{2}^{2}
$$

for all $\mathbf{x} \in \mathbb{R}^{n}$ with $\|\mathbf{x}\|_{0} \leq s$.

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## Thank you!

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